

Scheduling in Hard Real Time Systems



Known chores and durations

— find schedule satisfying constraints

- each job has a *completion deadline*

— uniprocessor

- rate-monotonic scheduling of *cyclic chores*, i.e., *periodic tasks*

- ◇ $t_i \rightarrow T_i, P_i$ where T is *execution time* and P is *period*

- ◇ e.g., a job is required to execute 2 seconds every 6 seconds

- ◇ deadline is implied

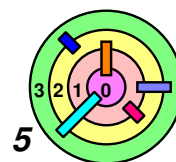
- Earliest Deadline First (EDF)

- ◇ $t_i \rightarrow d_i$ where d is a *deadline*

- ◇ optimality proof exists under certain conditions

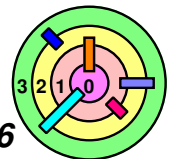
— multiprocessor

- often NP-complete ...



Assumptions

- ➡ Interrupts don't interfere (too much) with schedule
 - ▬ bounded interrupt delays
- ➡ Execution time really is predictable
 - ▬ what about effects of caching and paging?



Rate-Monotonic Scheduling

- ➡ Periodic chores
- ➡ period P_i
 - ➡ *per-cycle* processing time $T_i (\leq P_i)$
 - T_i / P_i is the *utilization* (or "*duty cycle*") of chore i

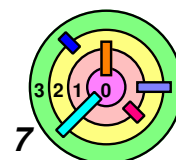
➡ necessary condition for feasibility: $\sum_{i=0}^{n-1} (T_i / P_i) \leq 1$

- ➡ Rate-monotonic scheduling
- ➡ each chore i is handled by a thread with *priority* $1/P_i$
 - ➡ *preemptive, priority scheduling*

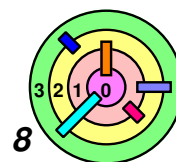
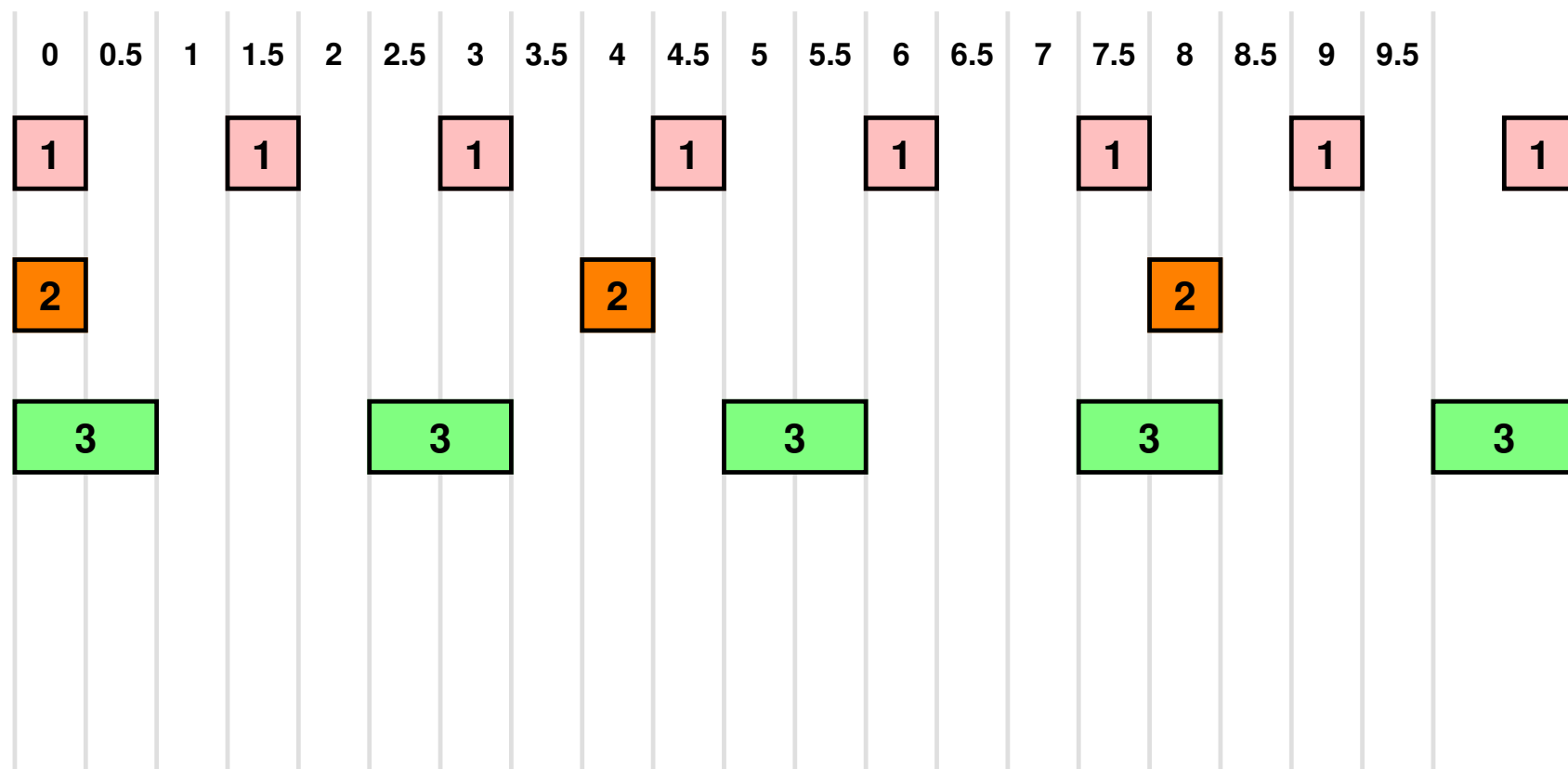
➡ works when $\sum_{i=0}^{n-1} (T_i / P_i) \leq n(2^{1/n} - 1)$

➡ $= \ln 2 \approx 0.6931$ in the limit

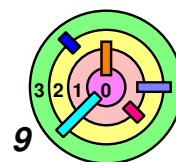
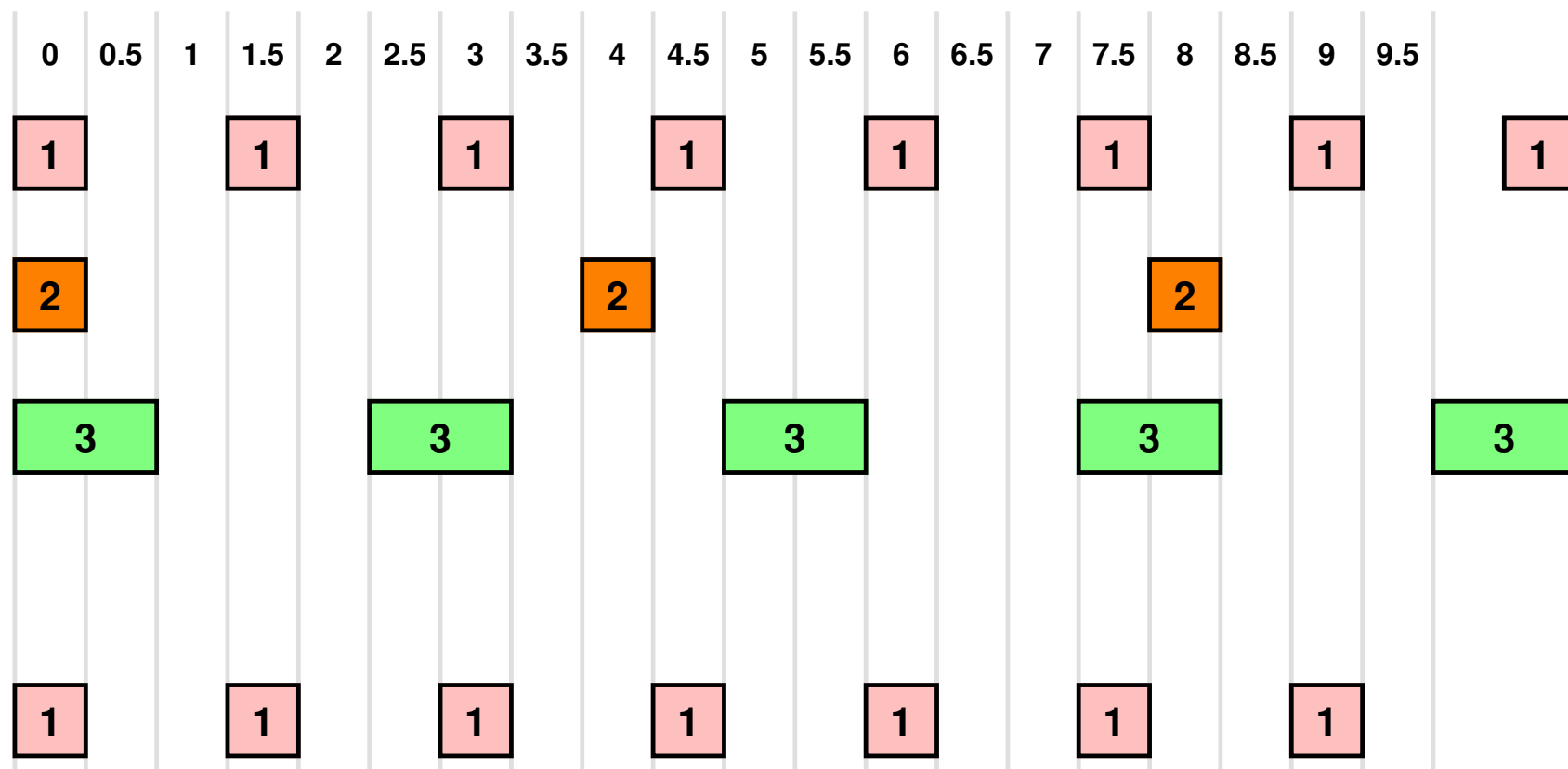
n	$n(2^{1/n} - 1)$
1	1
2	0.8284
3	0.7798
4	0.7568
5	0.7435
...	...



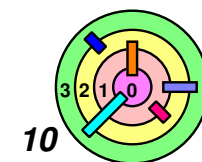
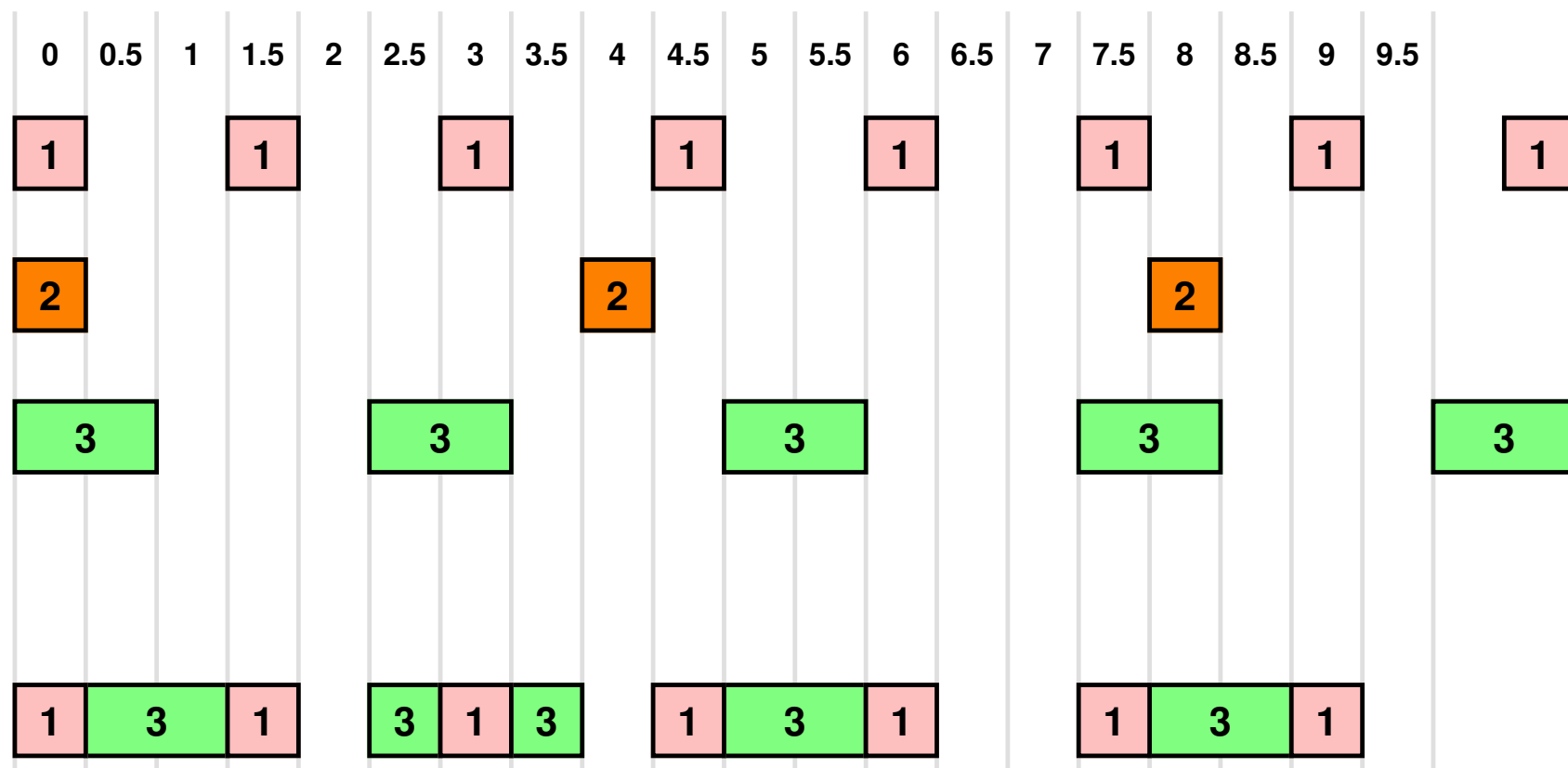
Scenario 1



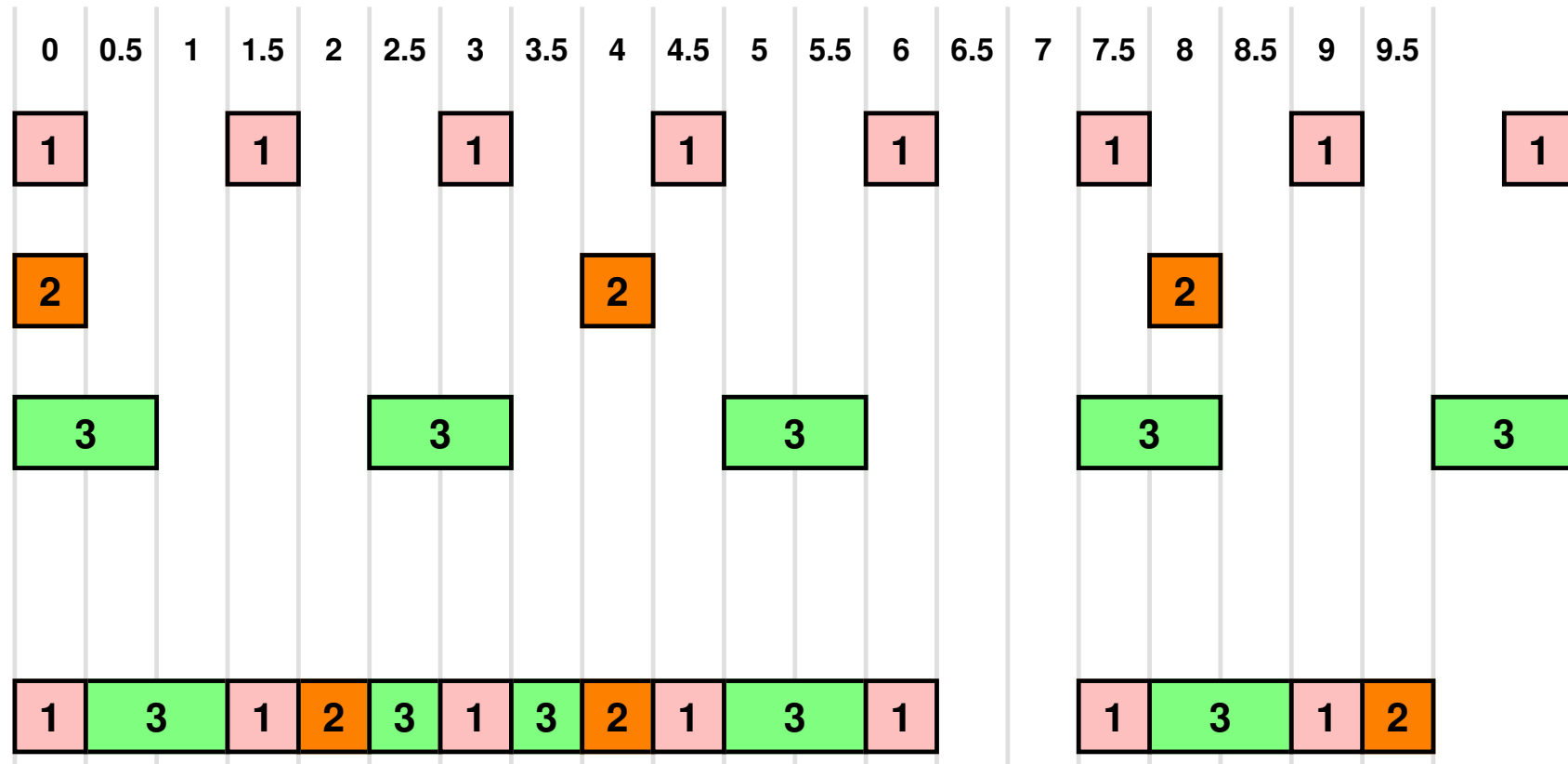
Scenario 1



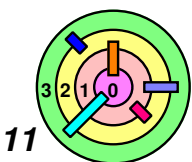
Scenario 1



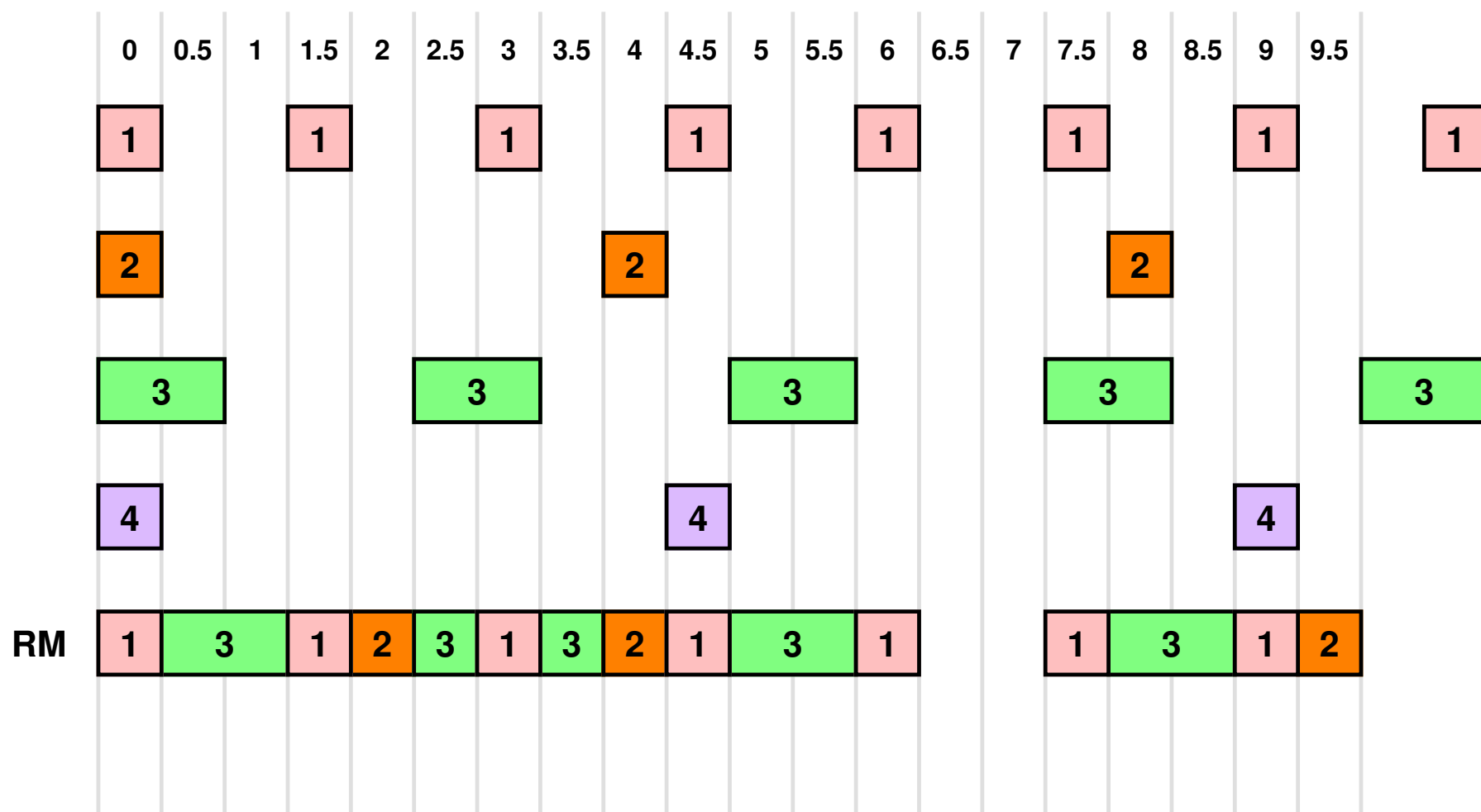
Scenario 1



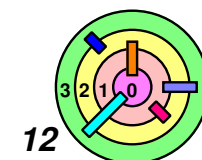
- rate-monotonic scheduler succeeded
- need to simulate till time = 60 to be certain



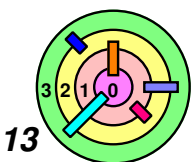
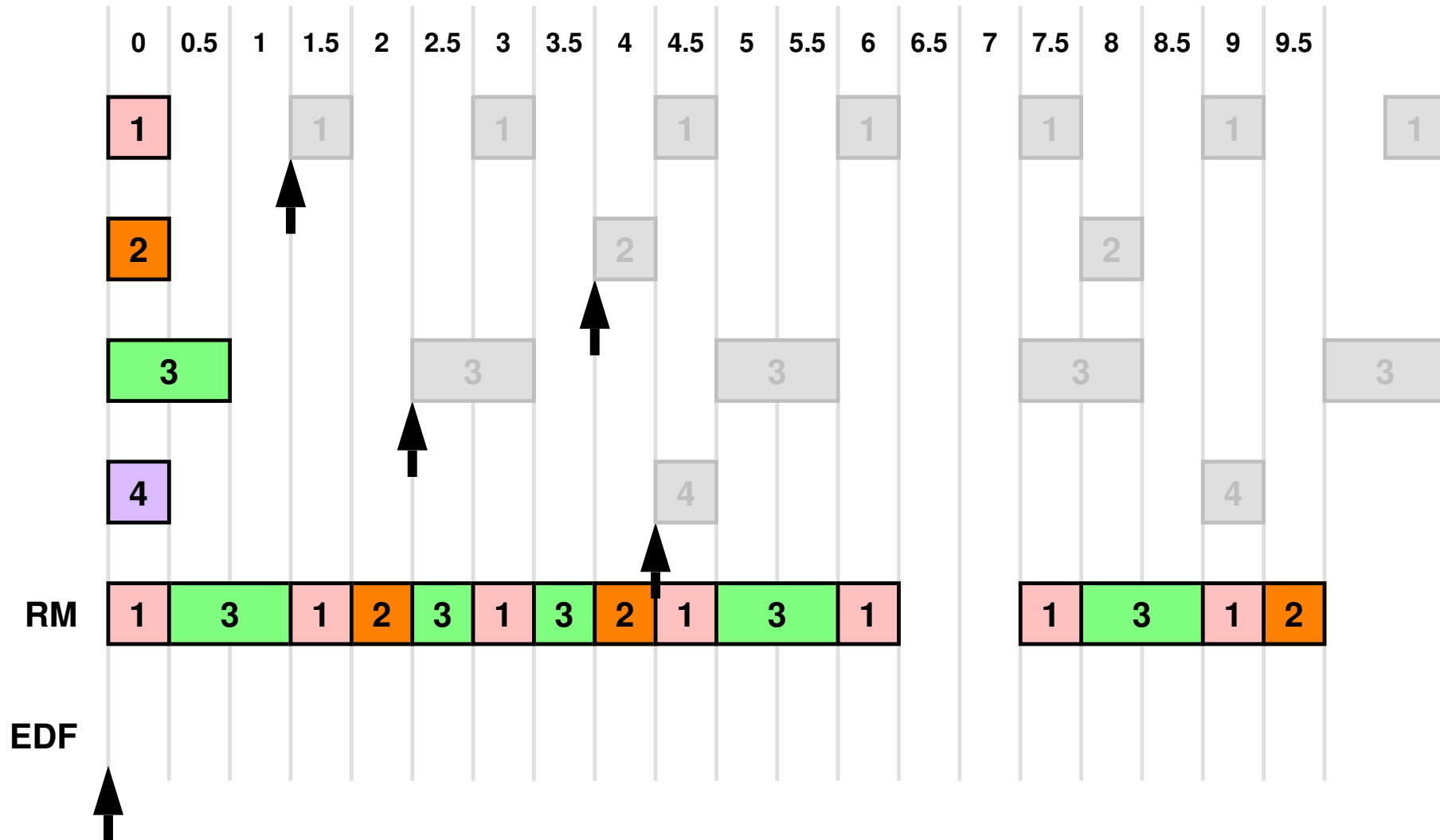
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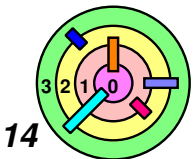
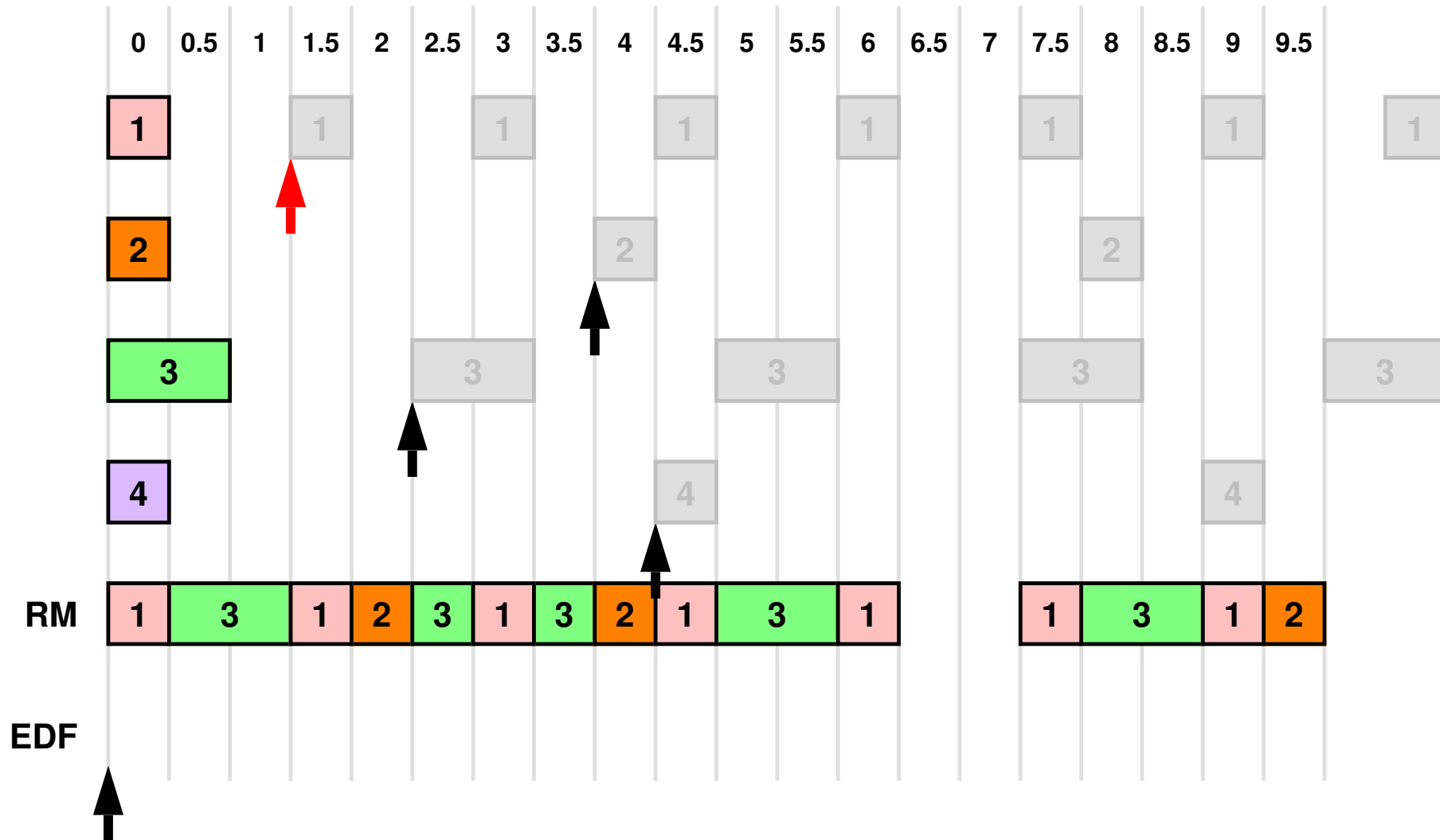
- rate-monotonic scheduler would fail in this case
- try EDF



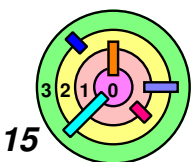
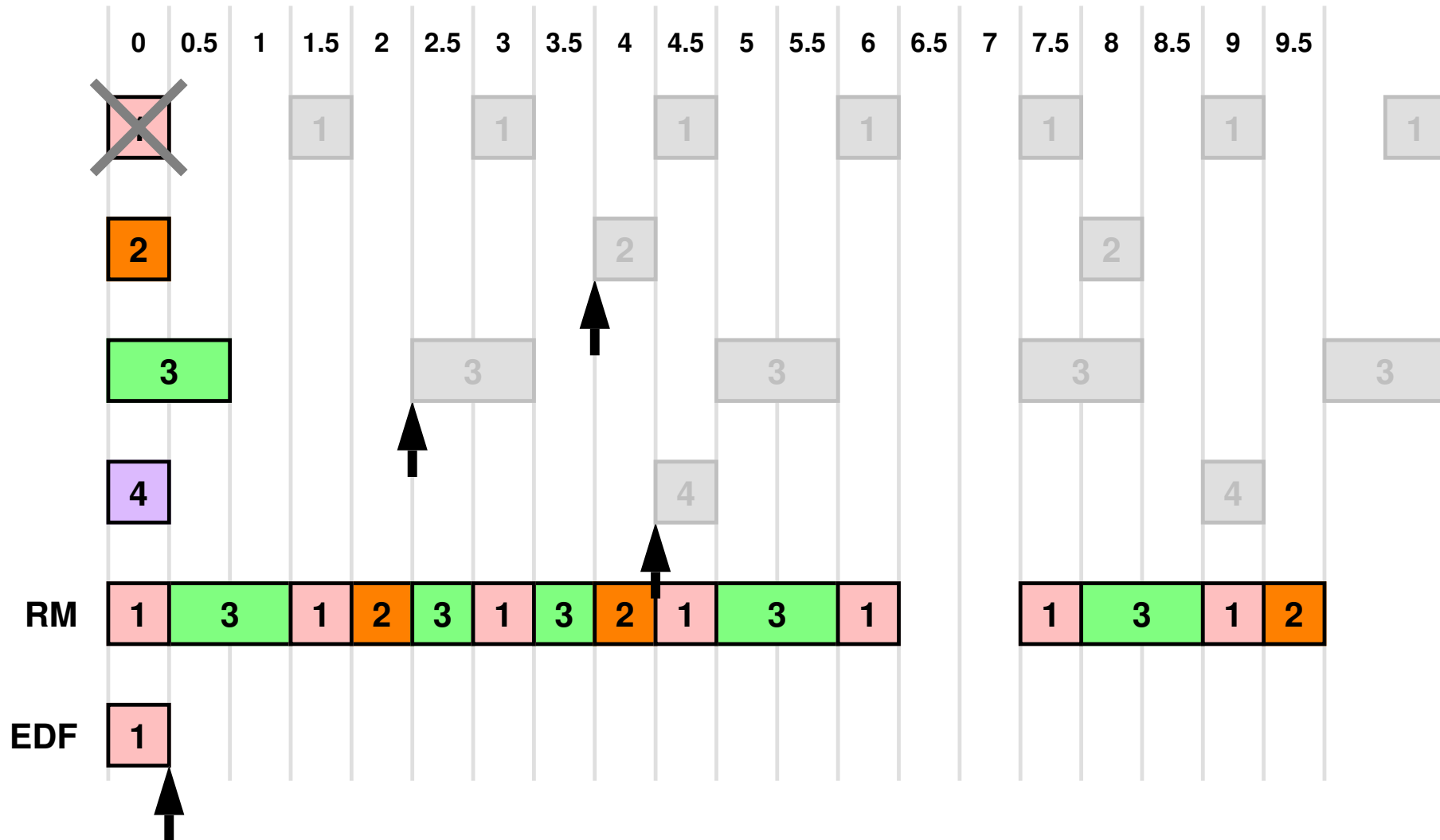
Earliest Deadline First



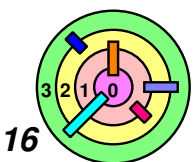
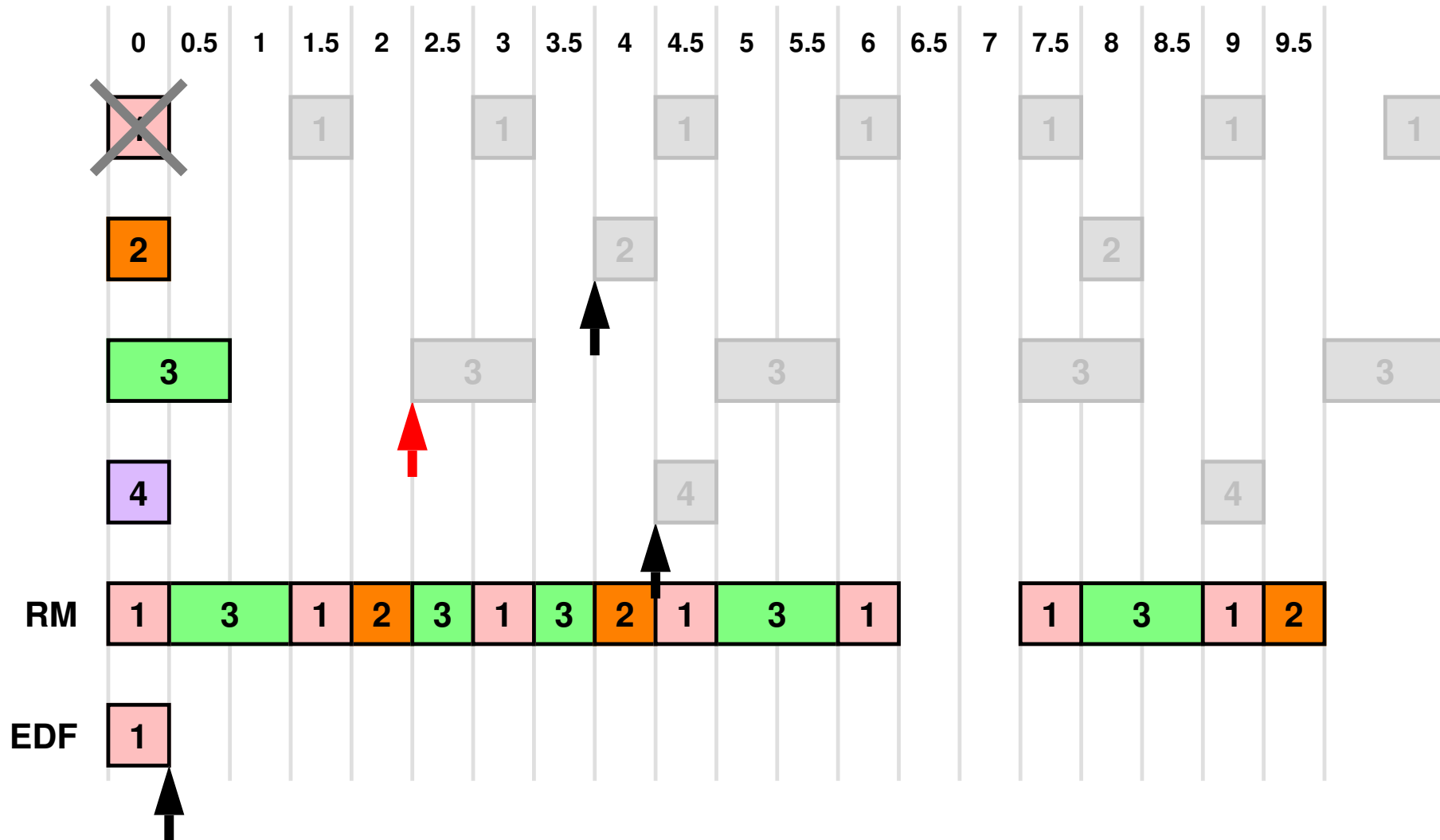
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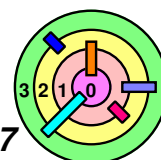


Earliest Deadline First

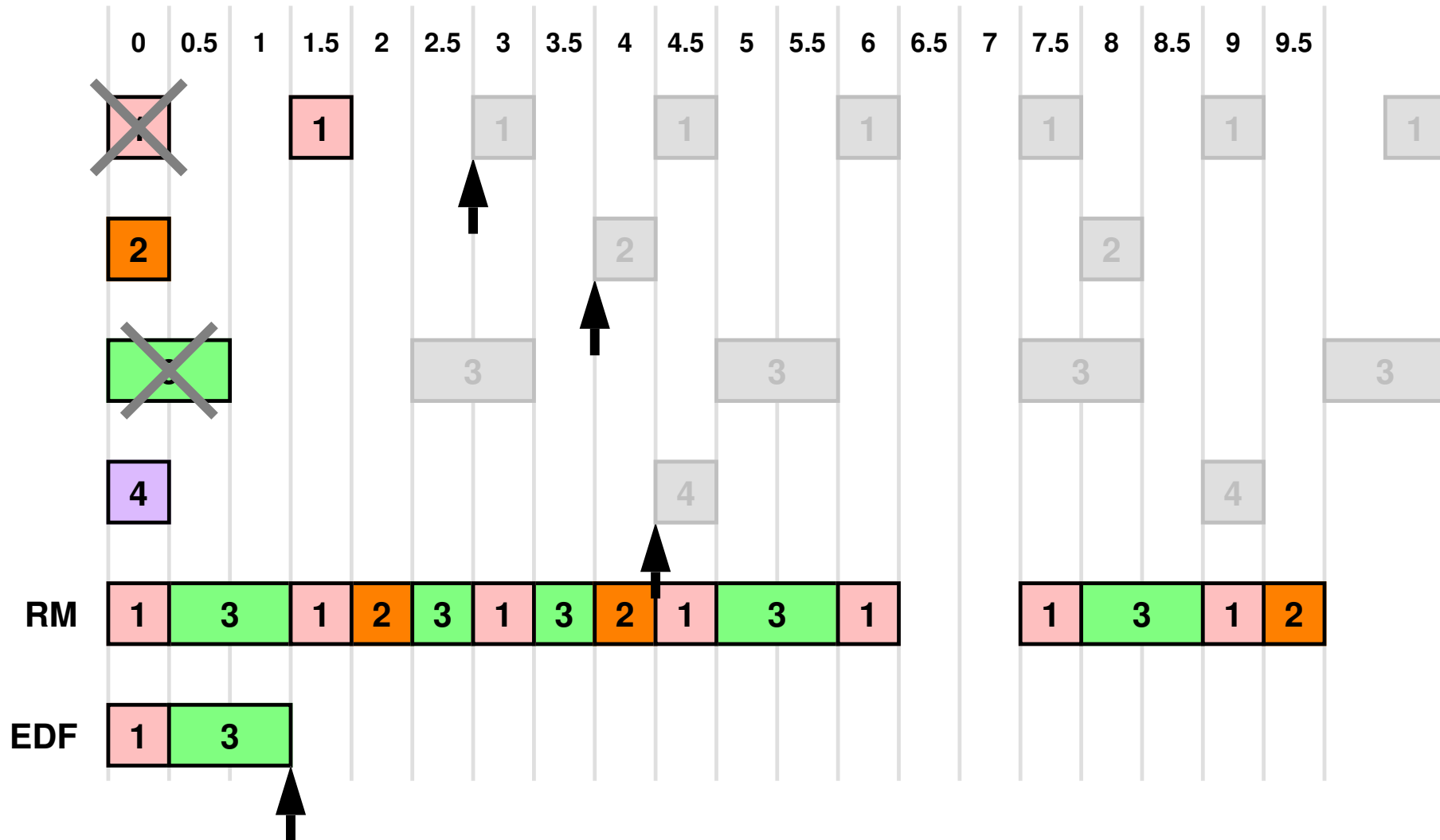


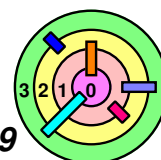
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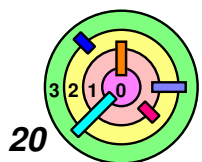
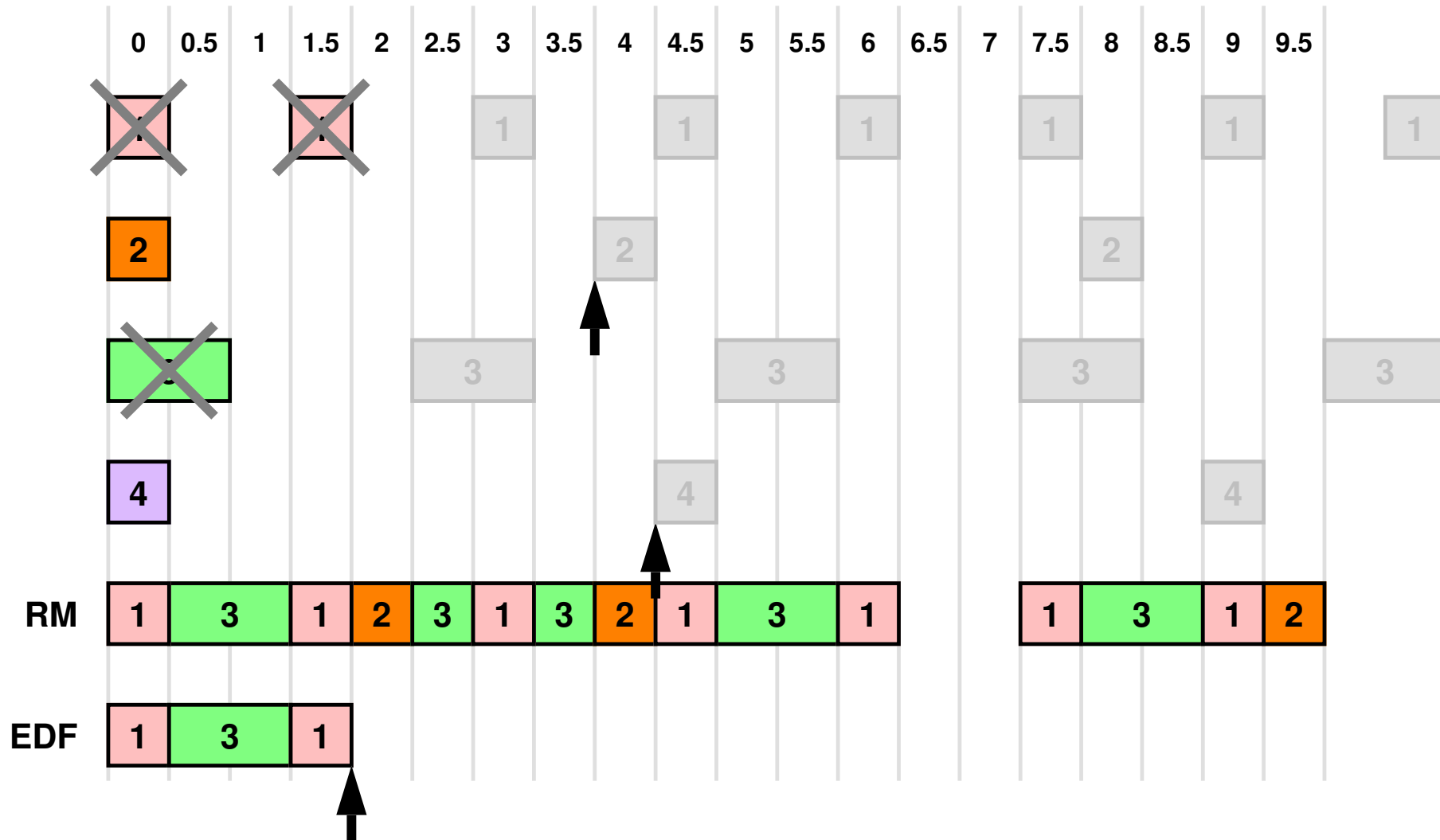


Earliest Deadline First

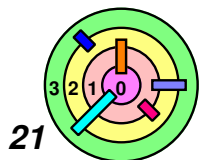
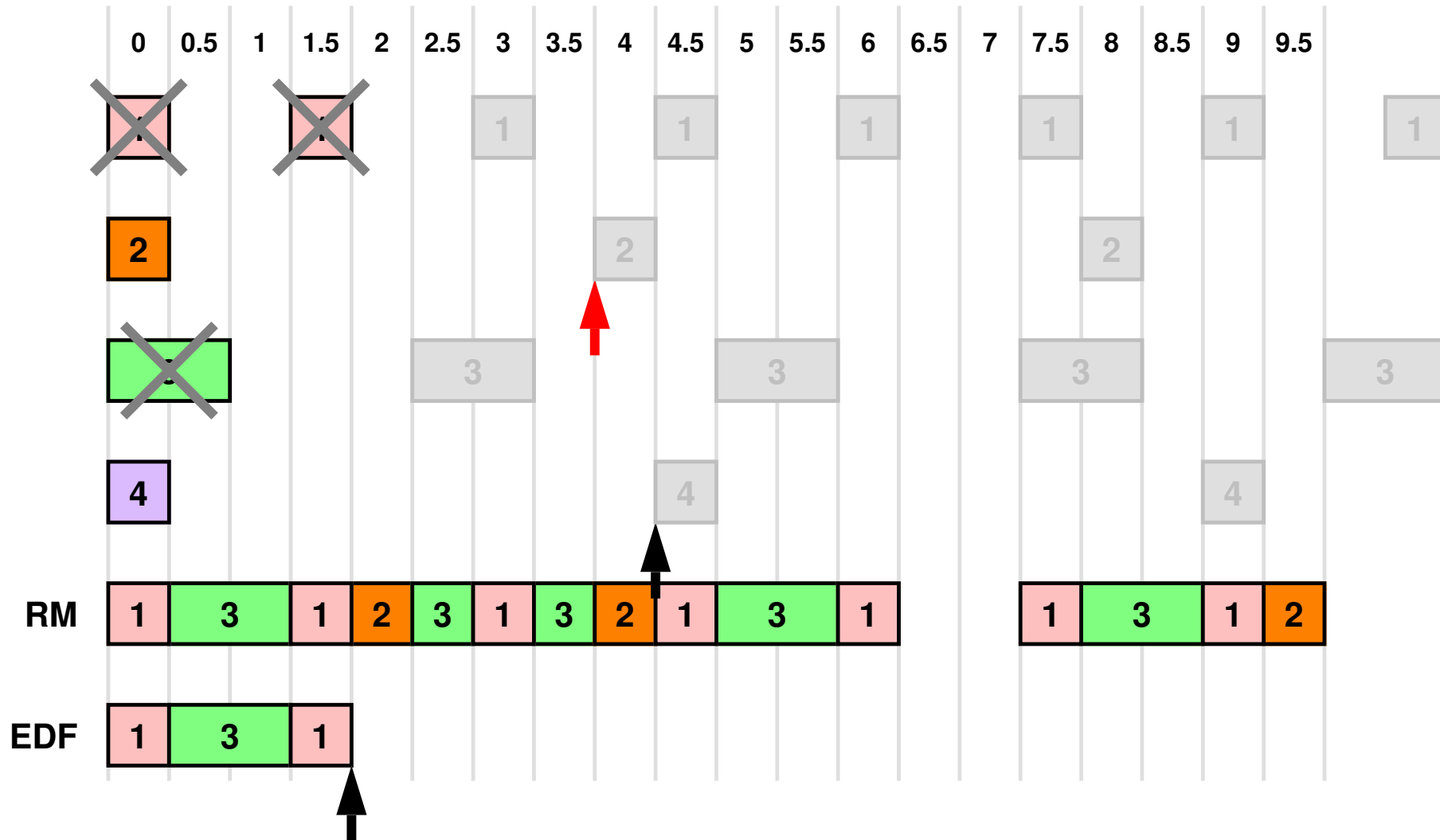




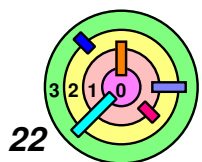
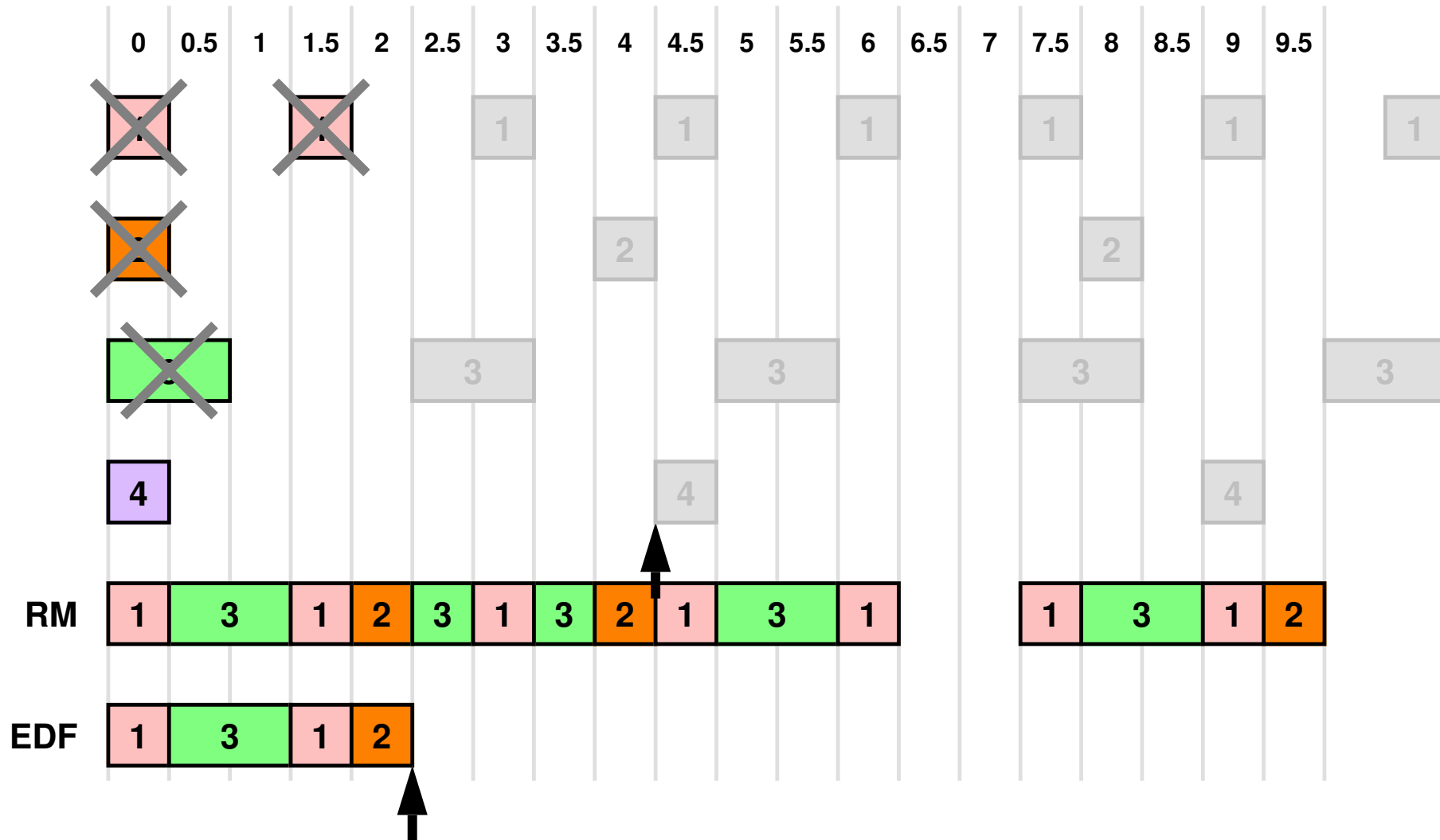
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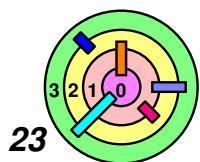
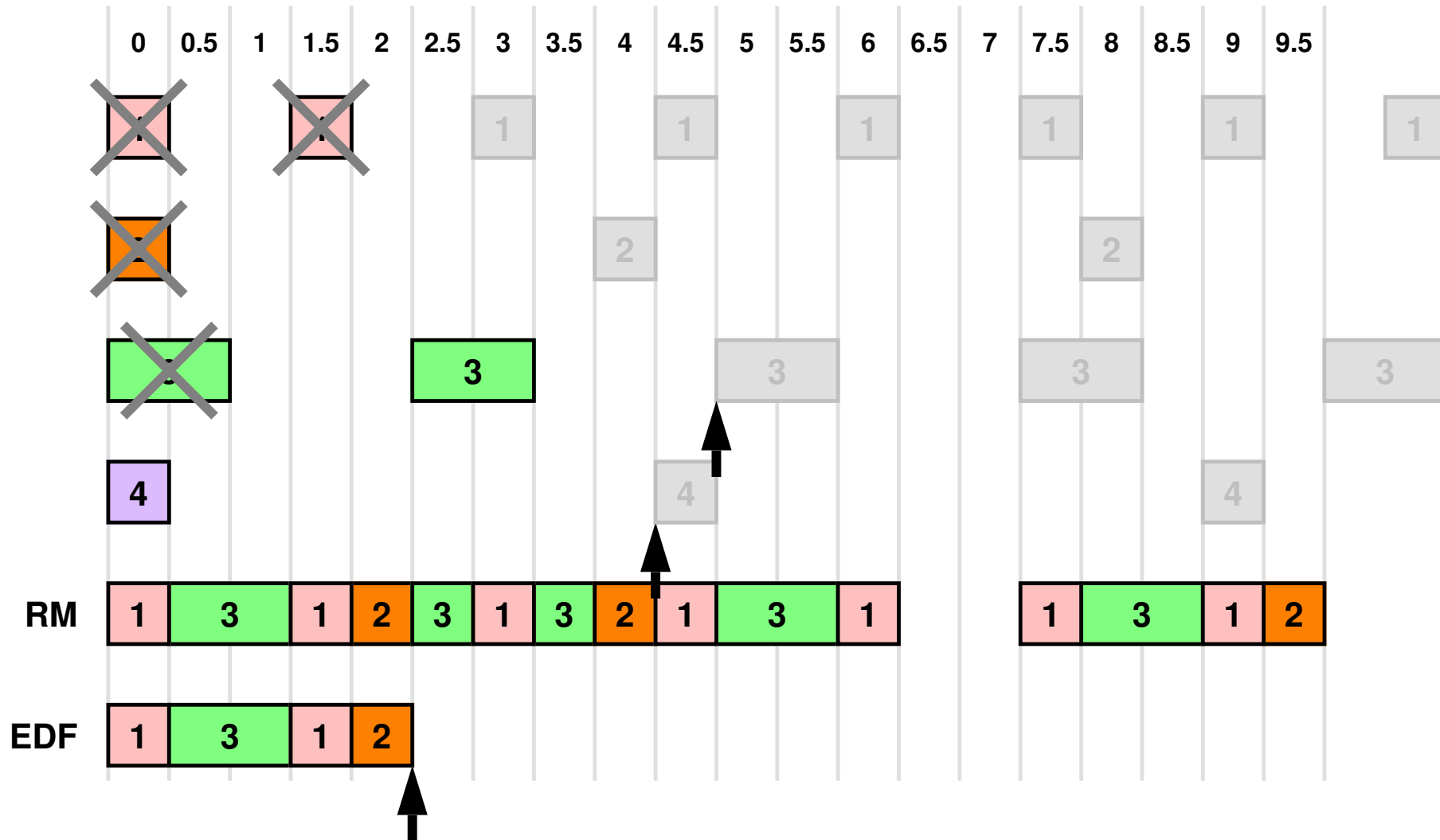
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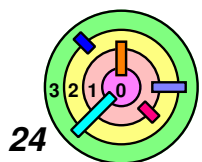
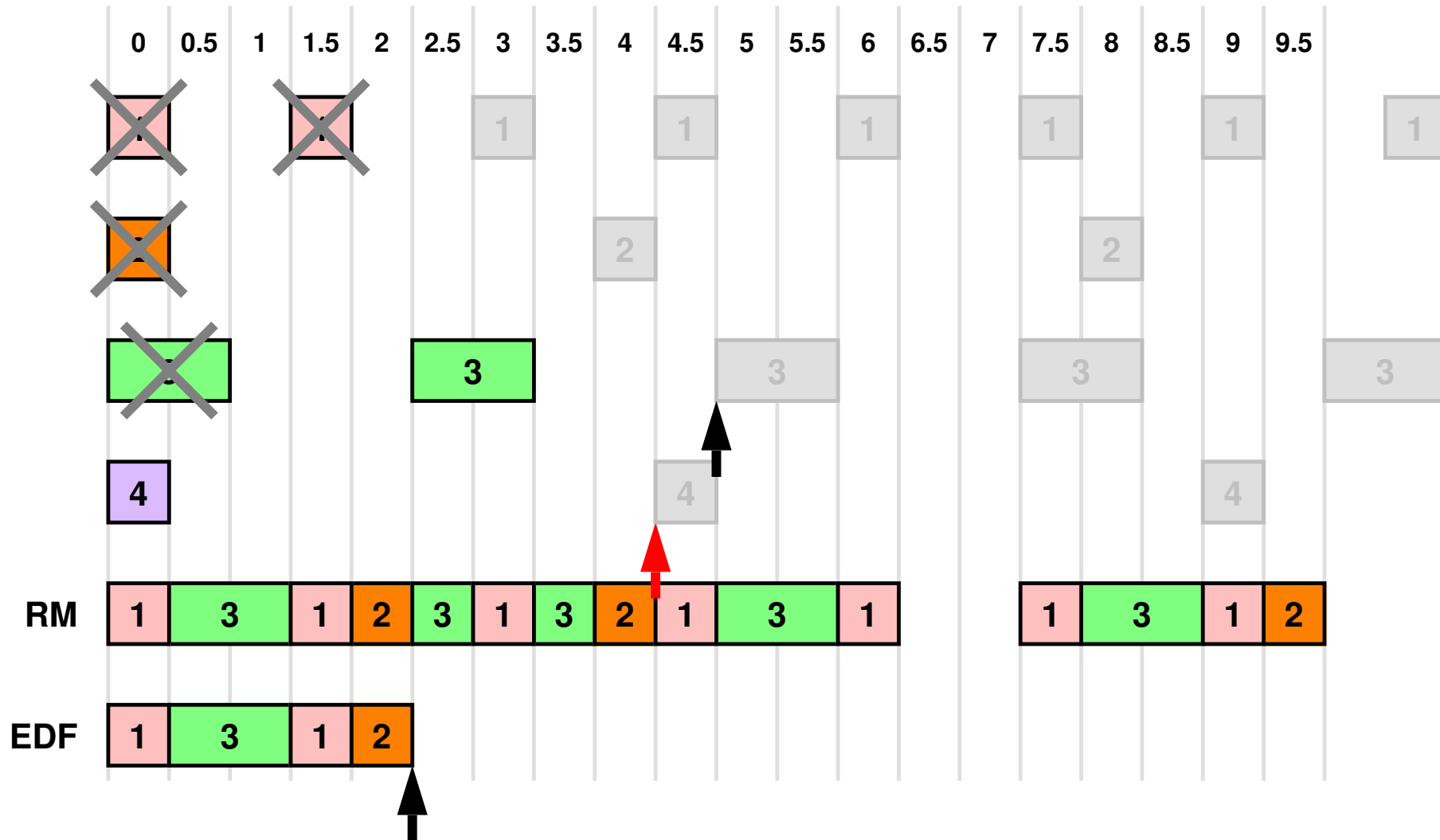
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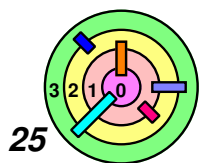
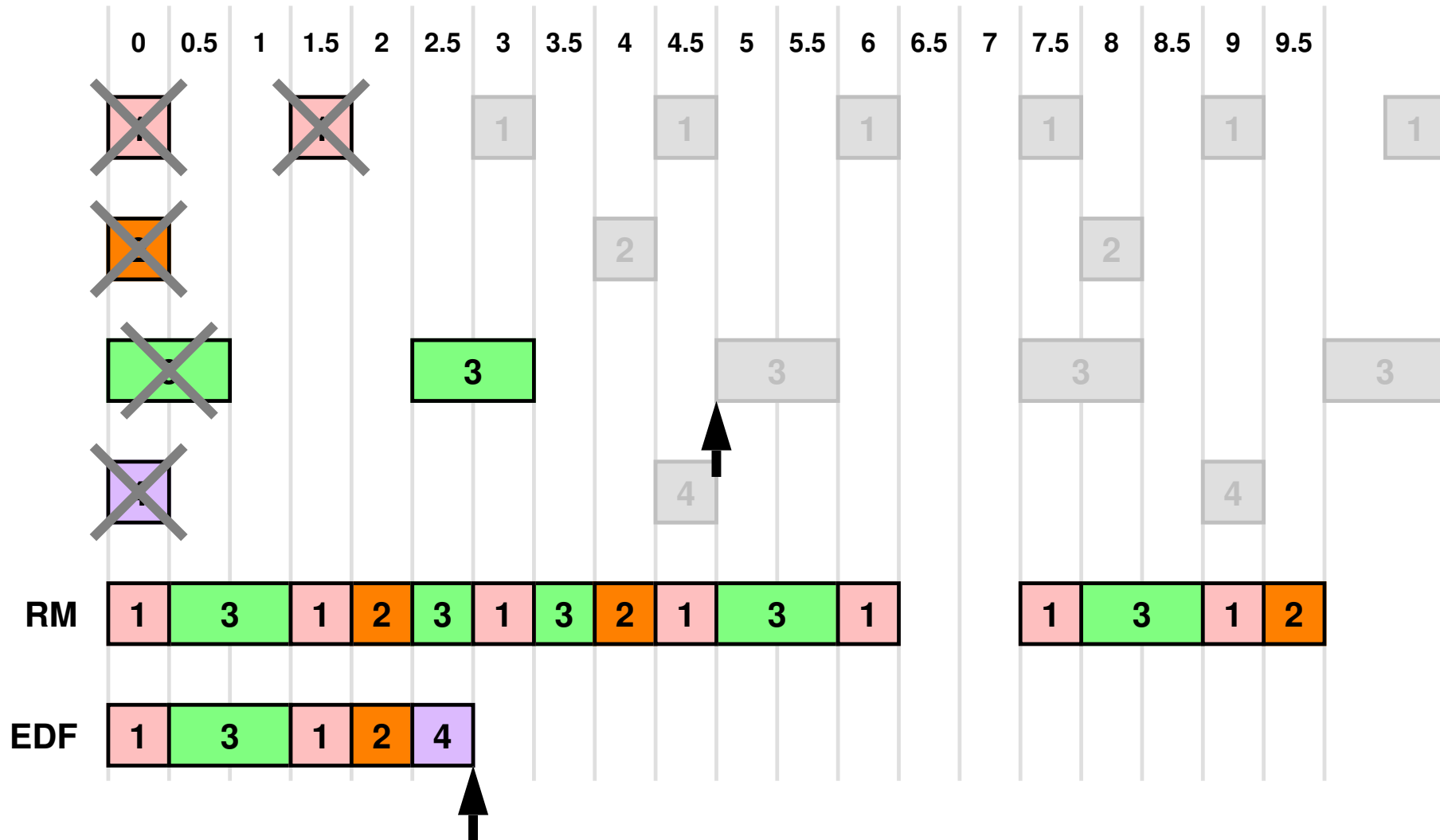
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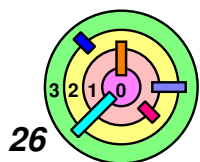
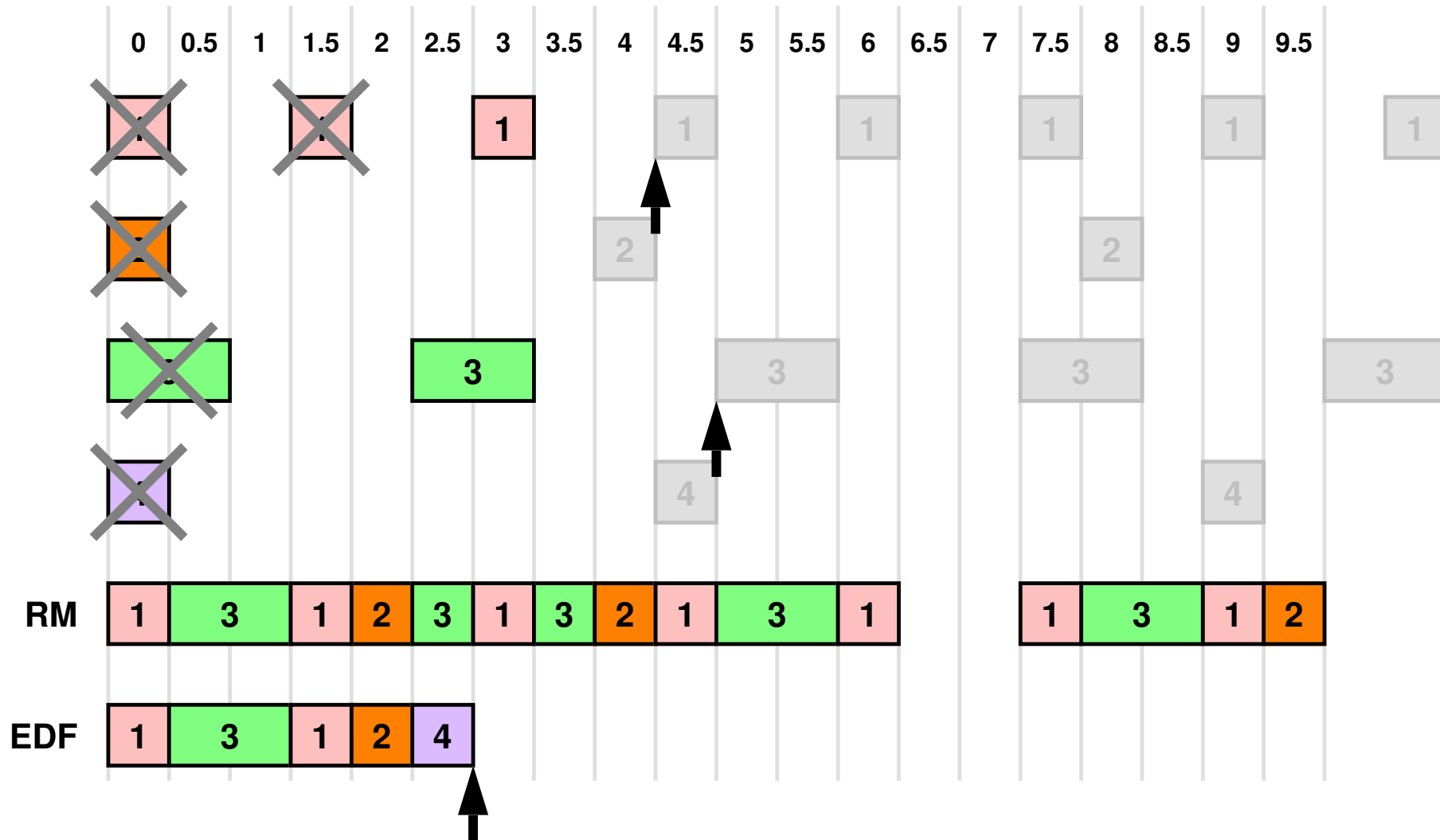
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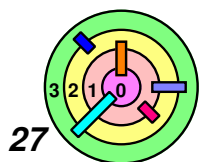


Earliest Deadline First

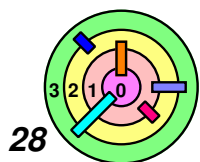
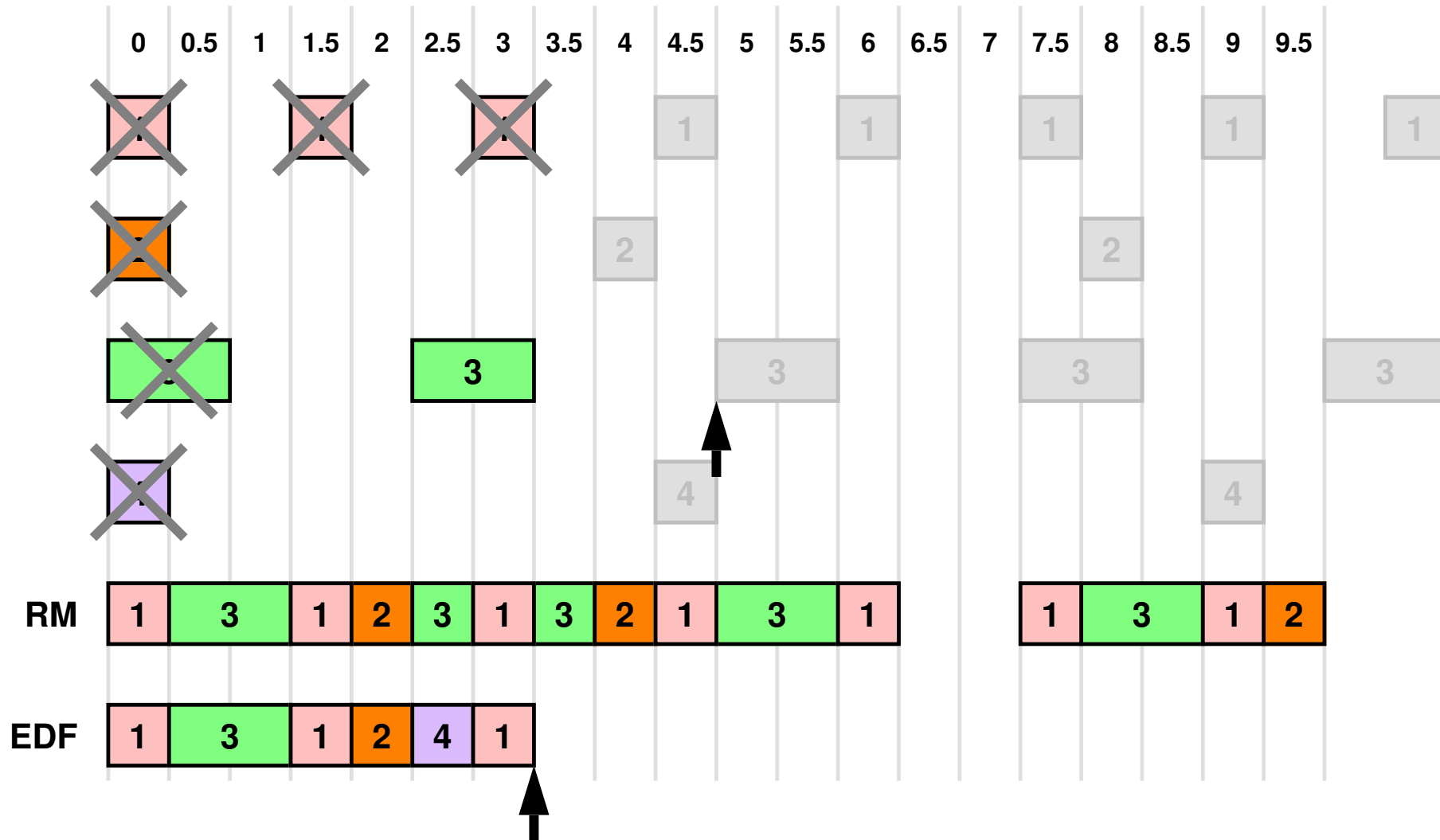


Earliest Deadline First

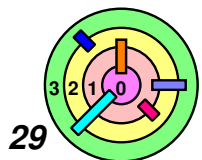
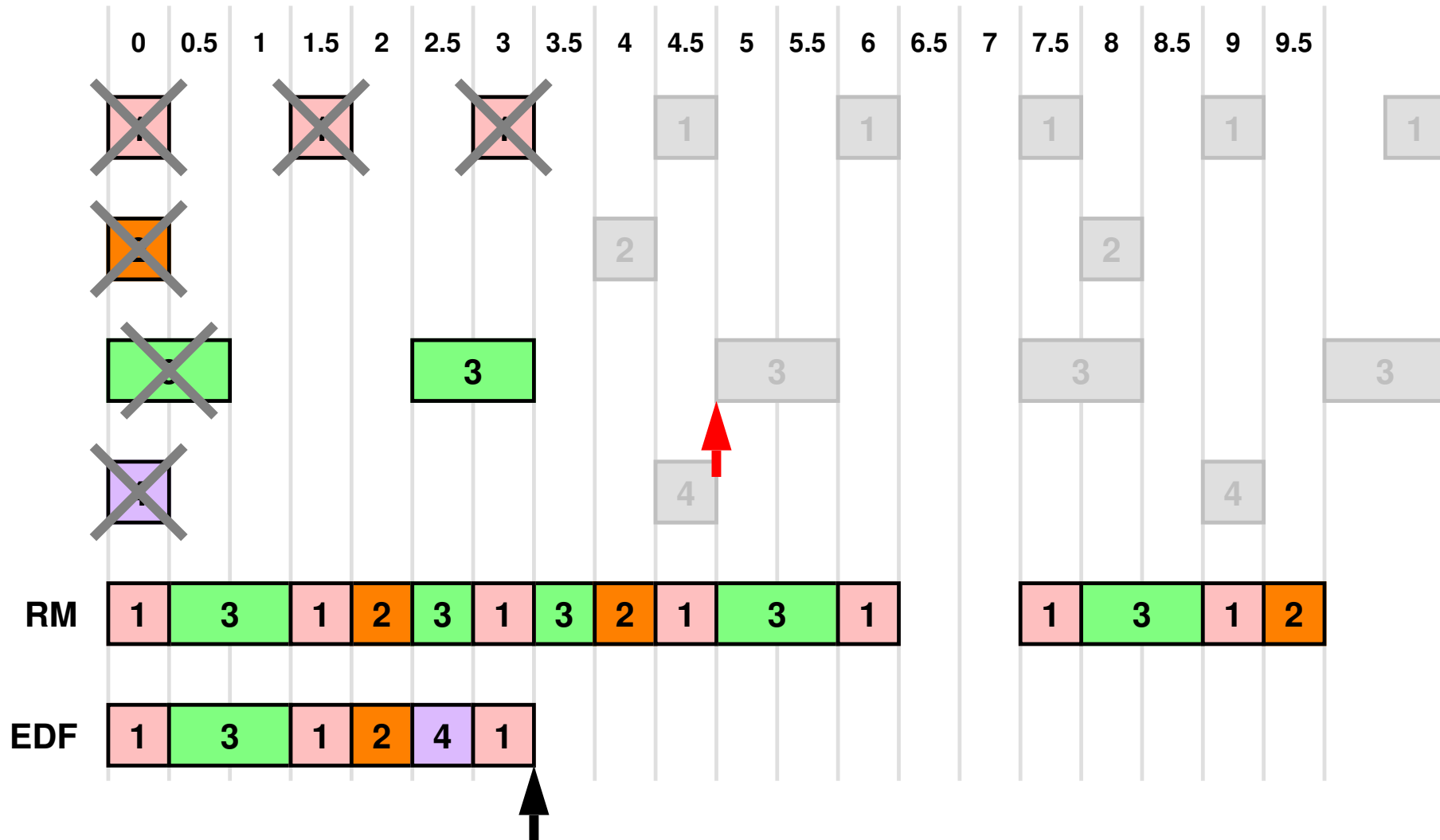




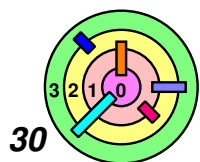
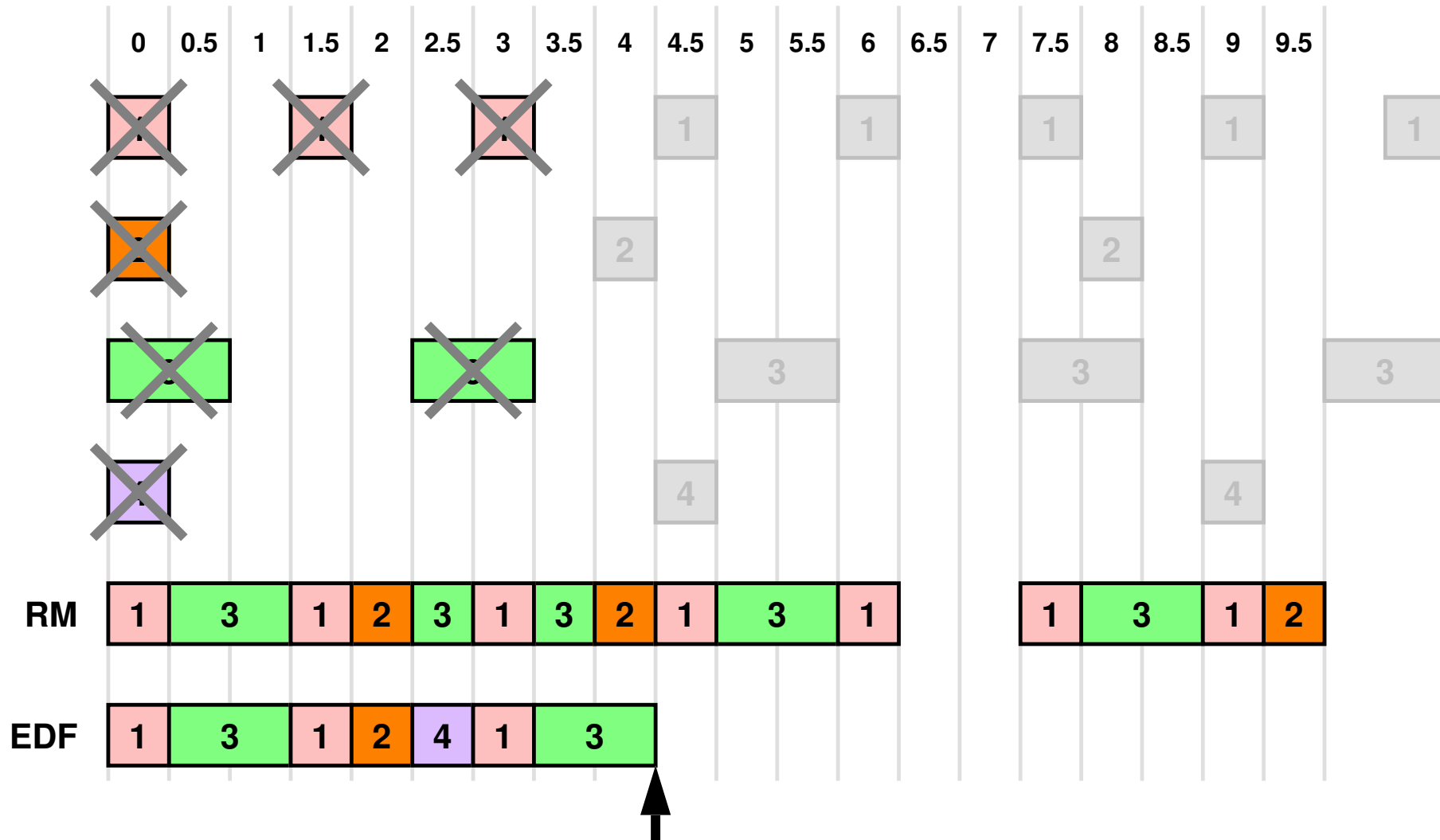
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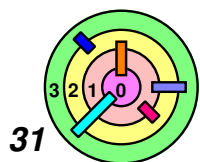
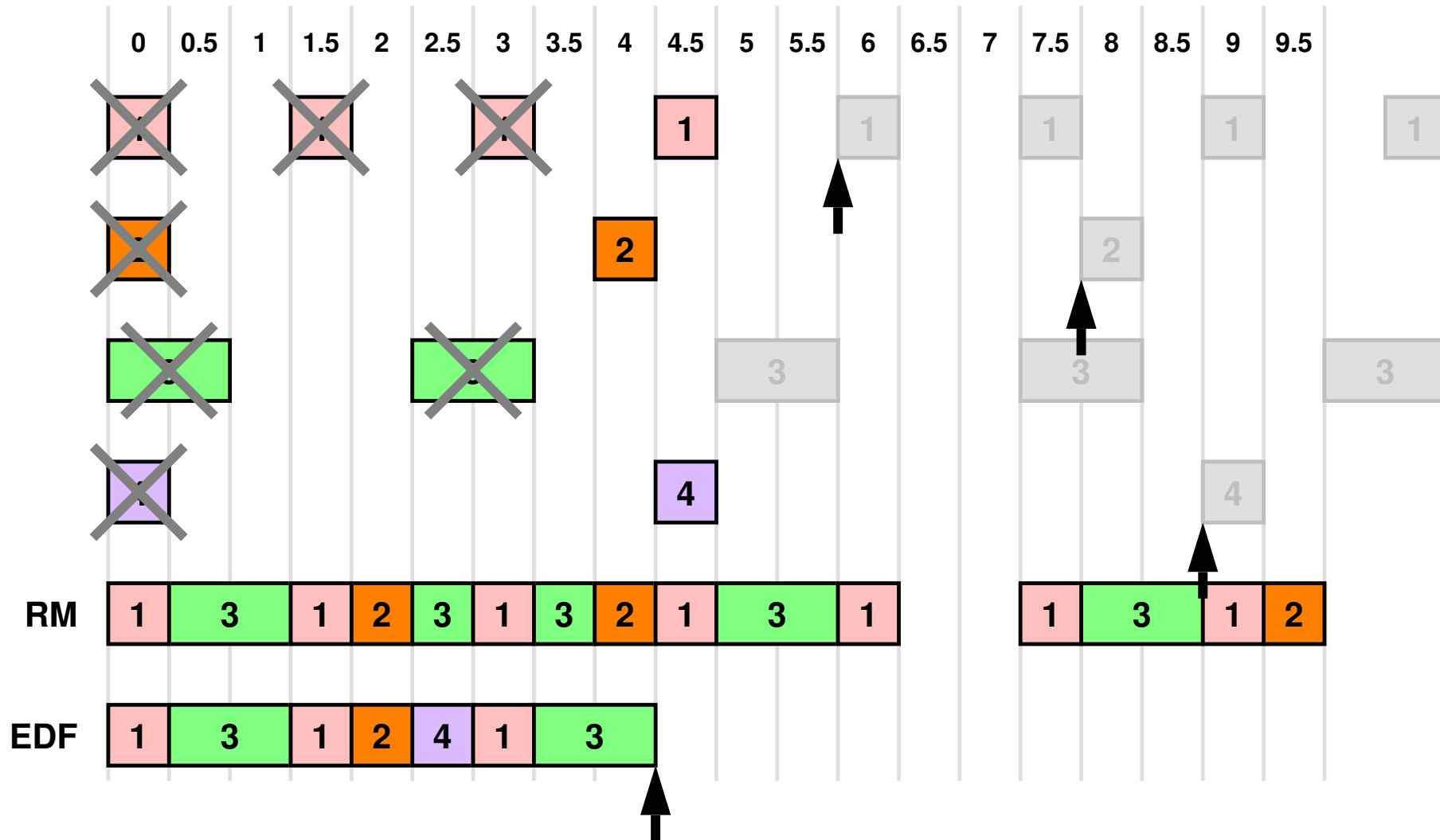
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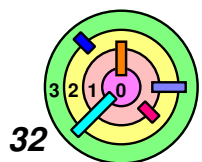
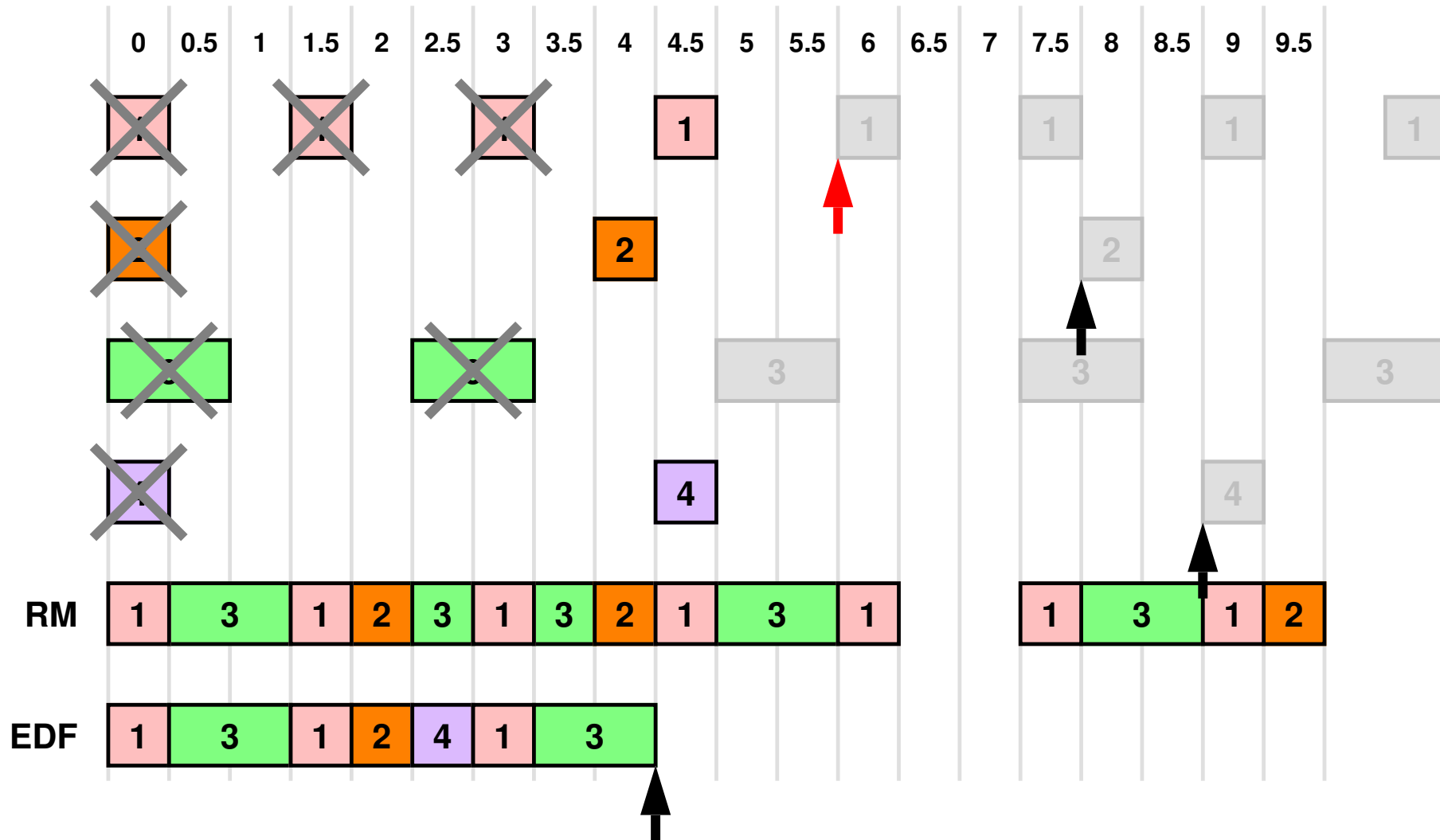
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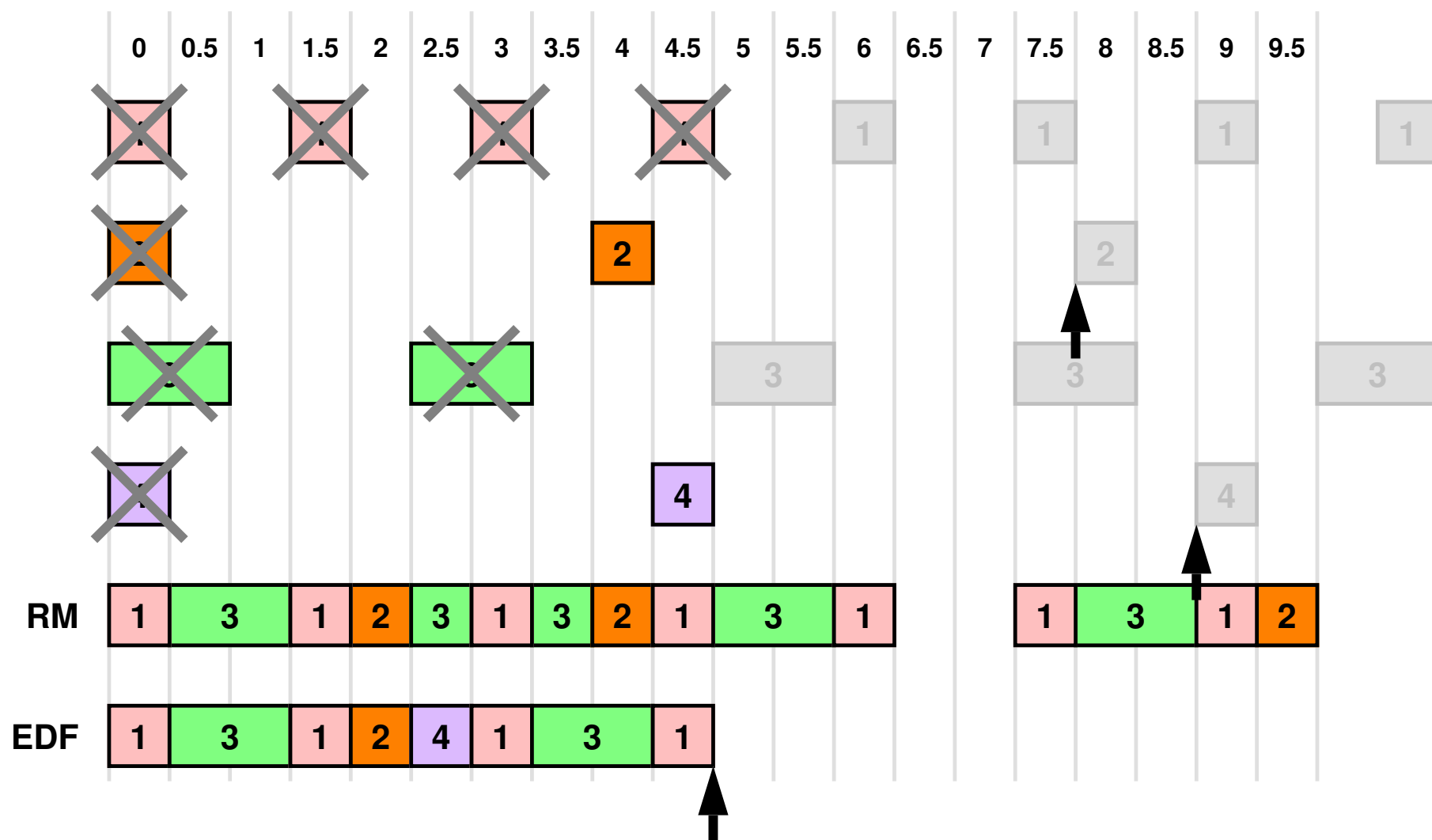
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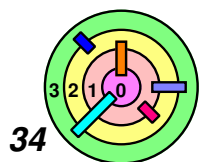
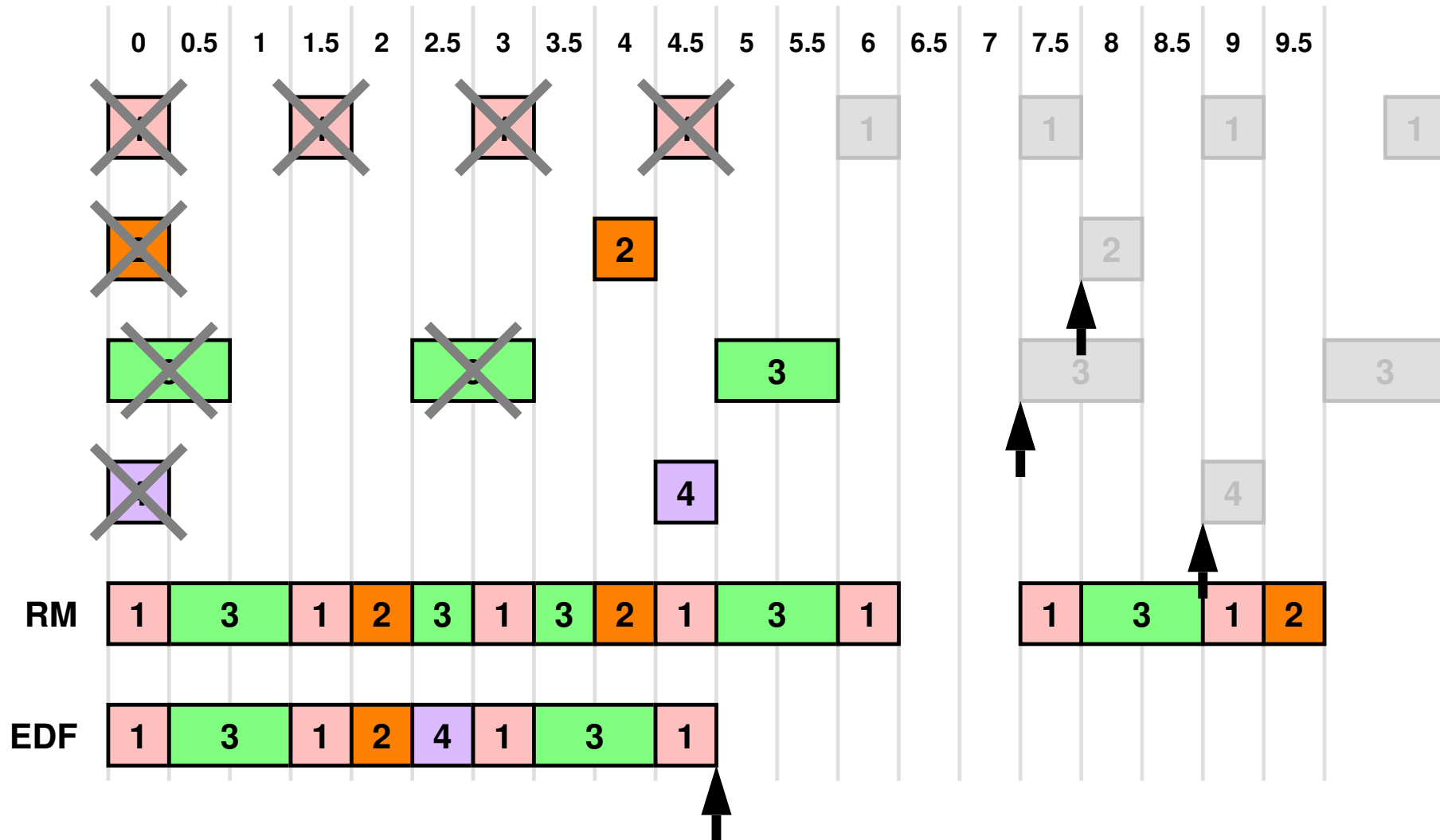
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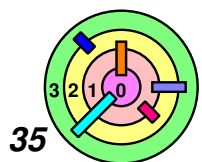
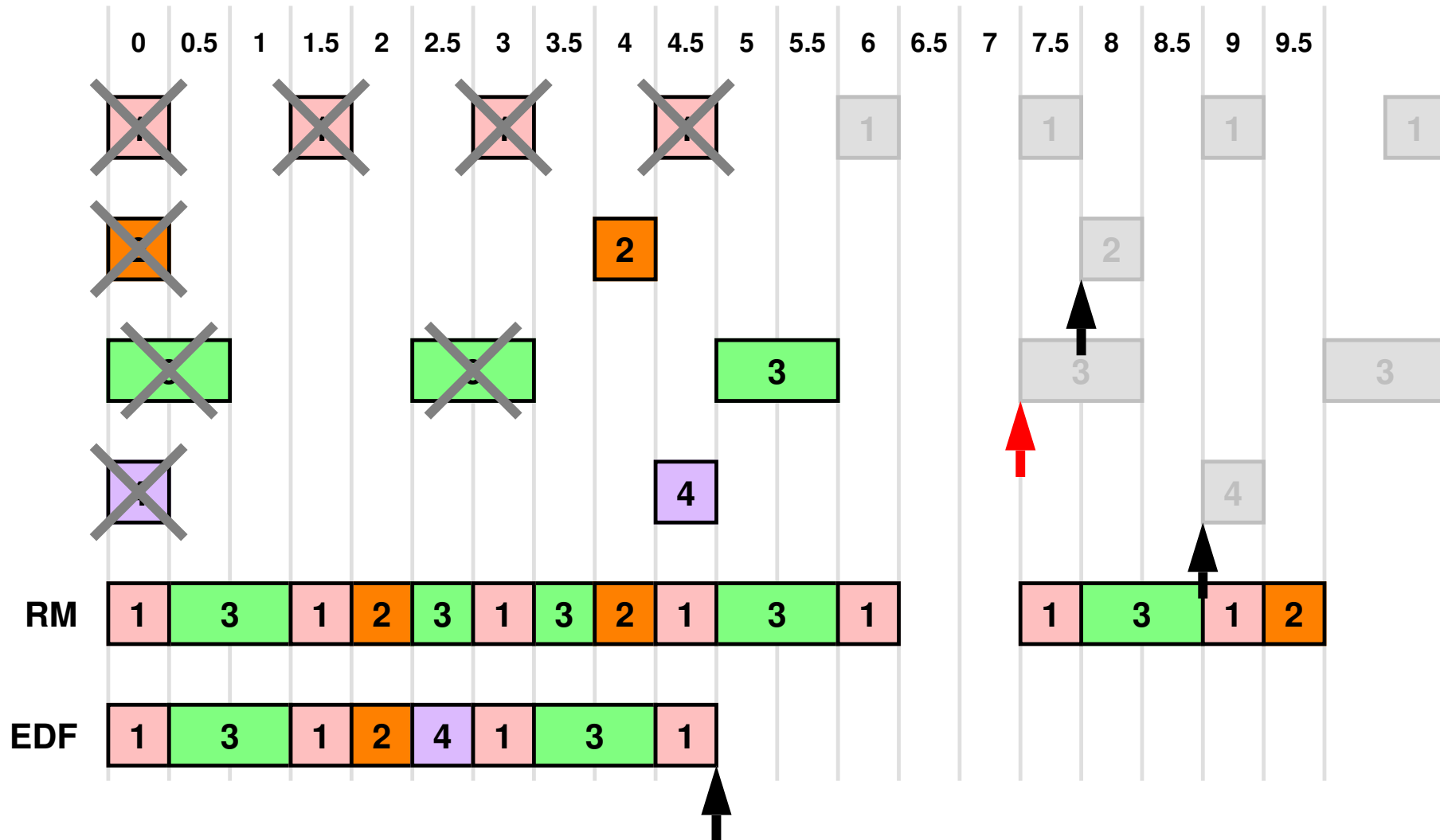
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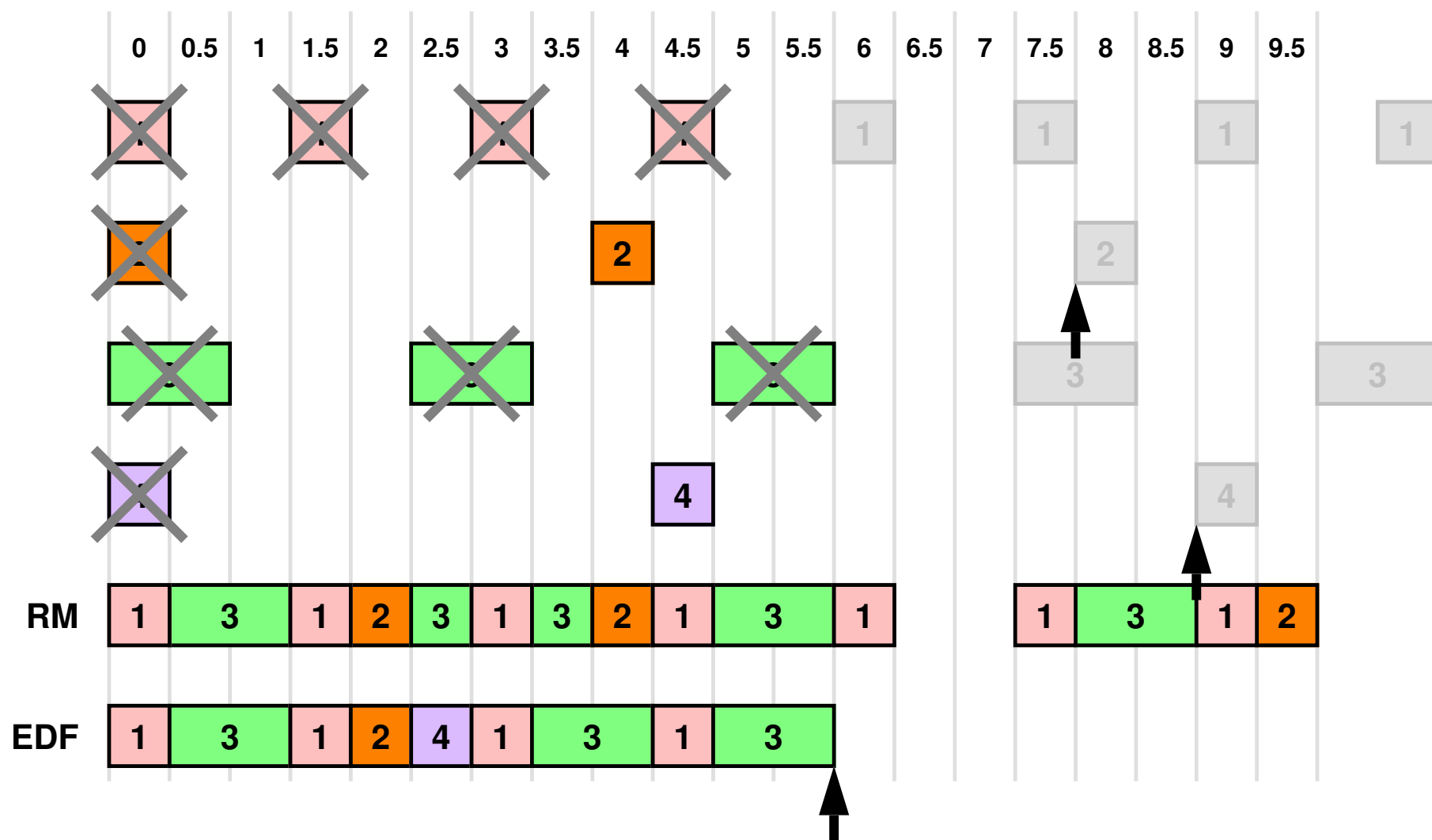
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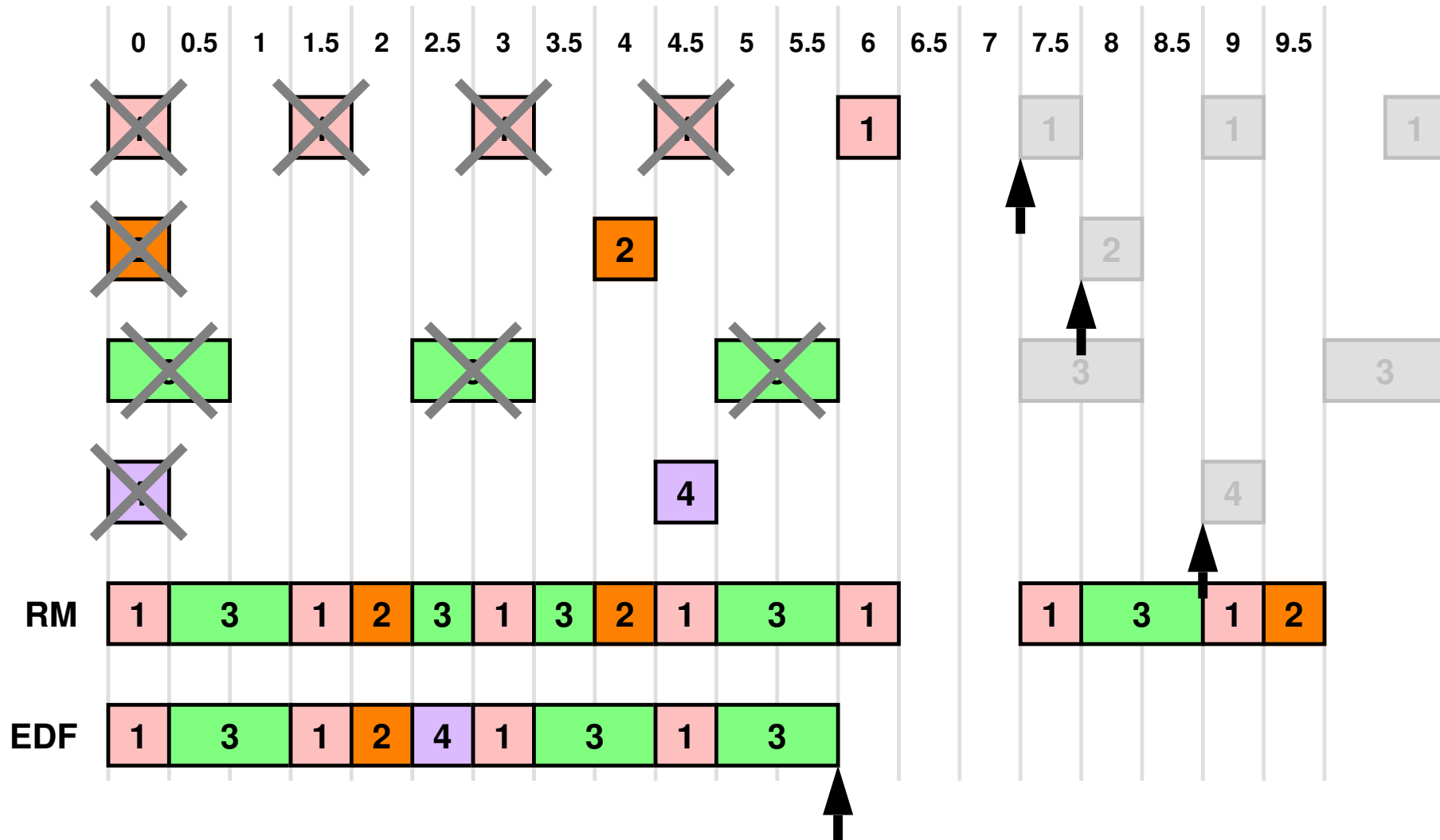
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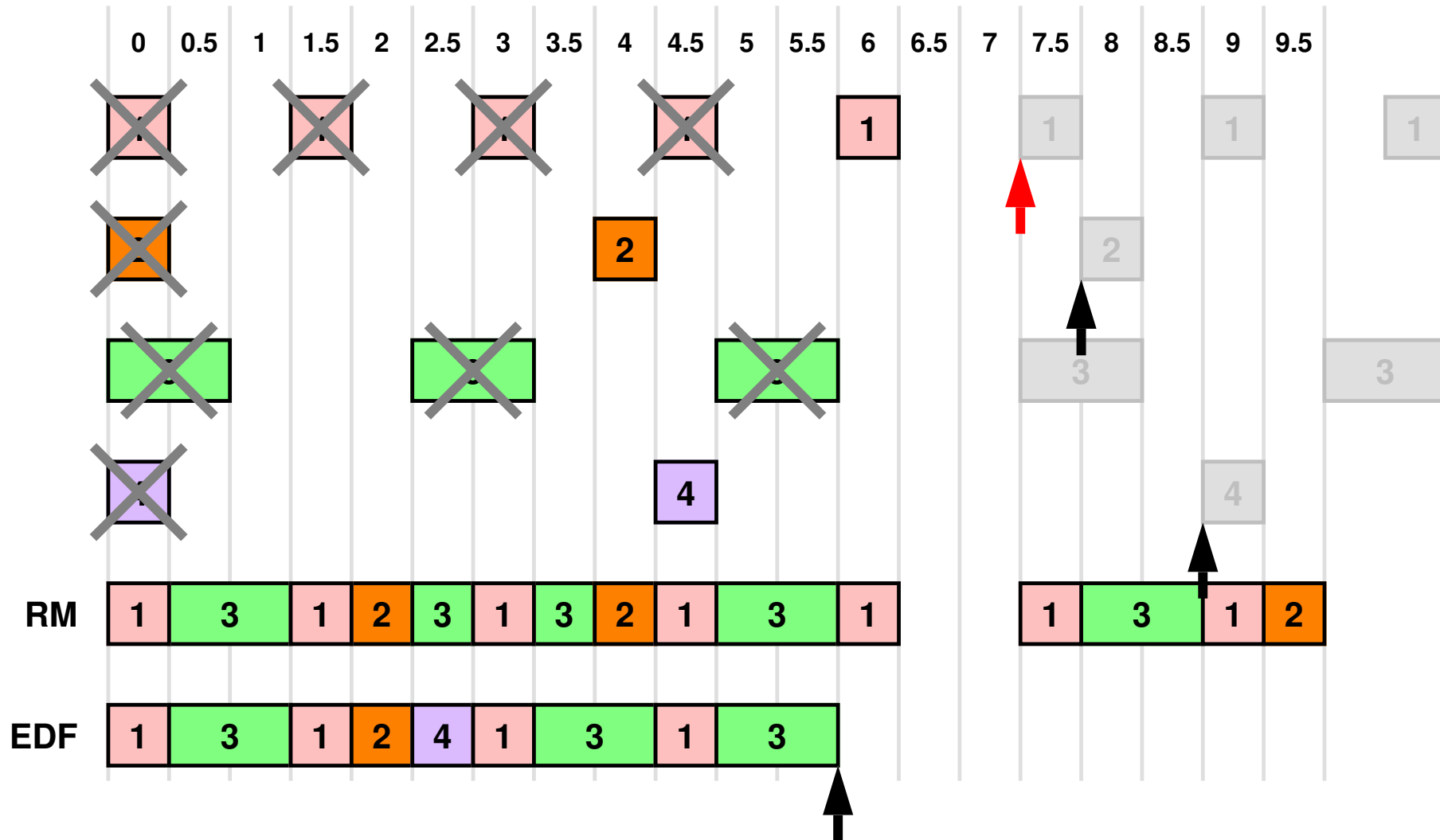
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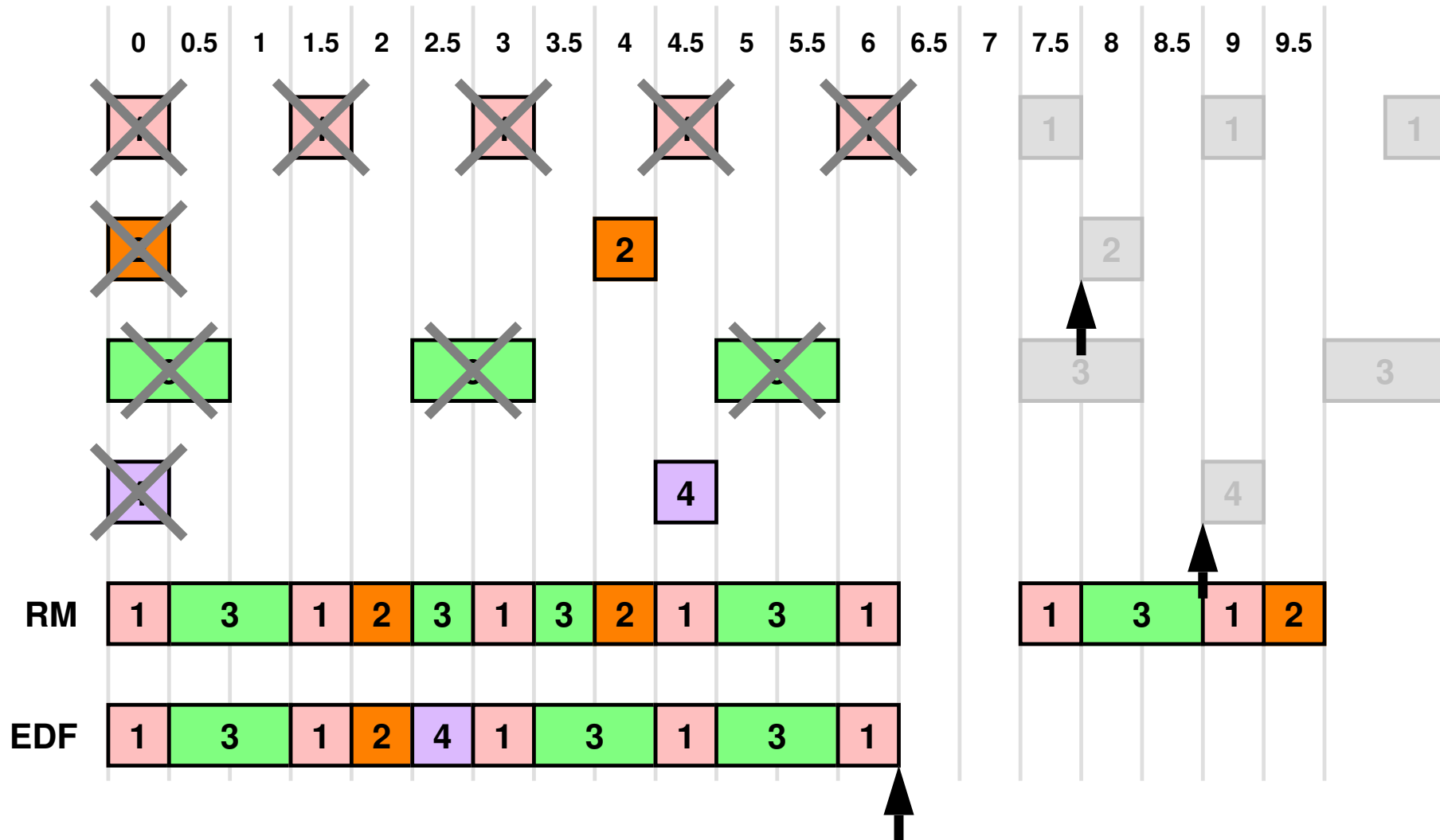
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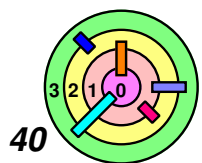
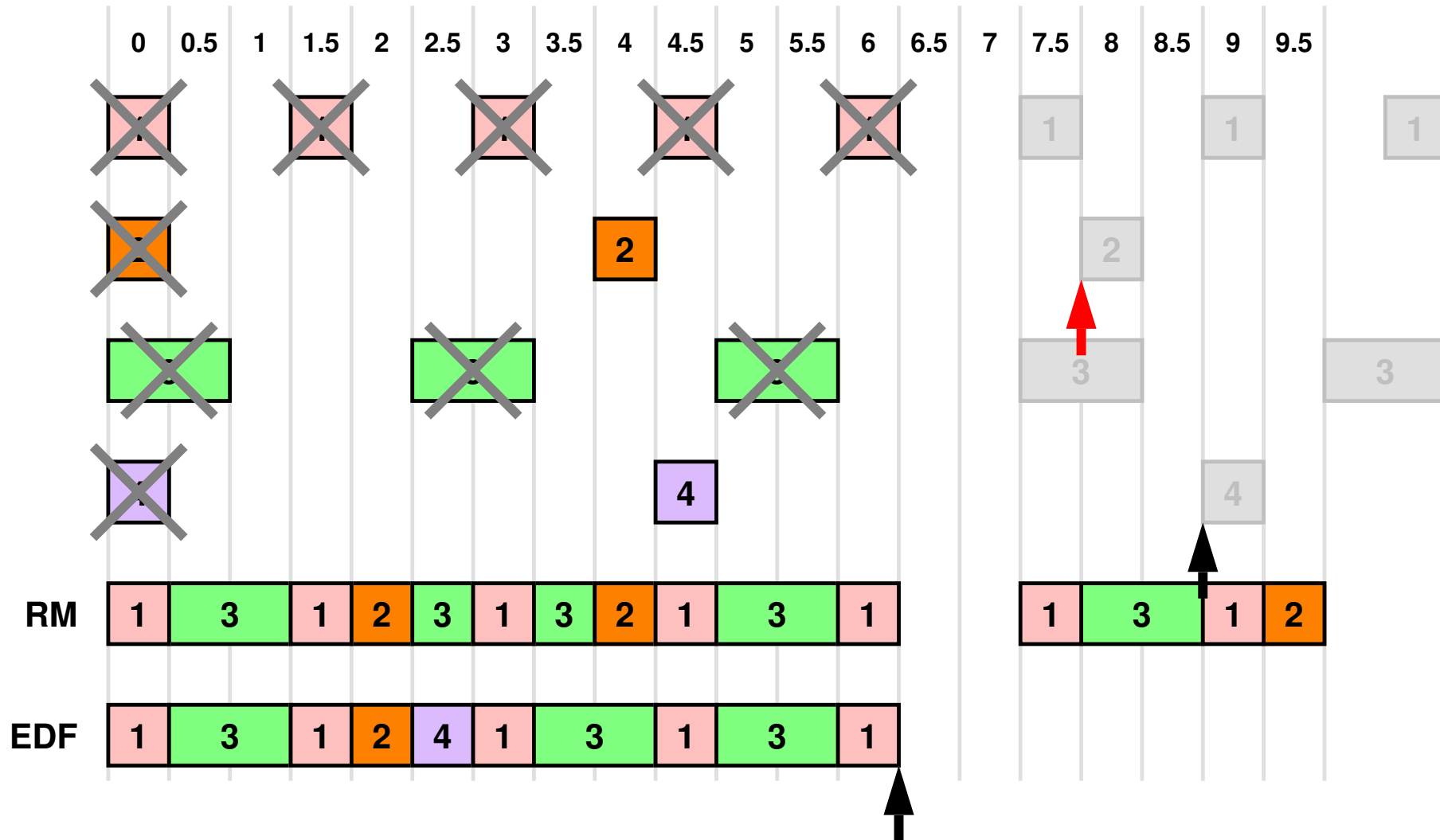
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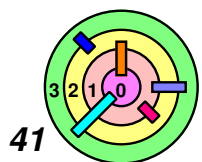
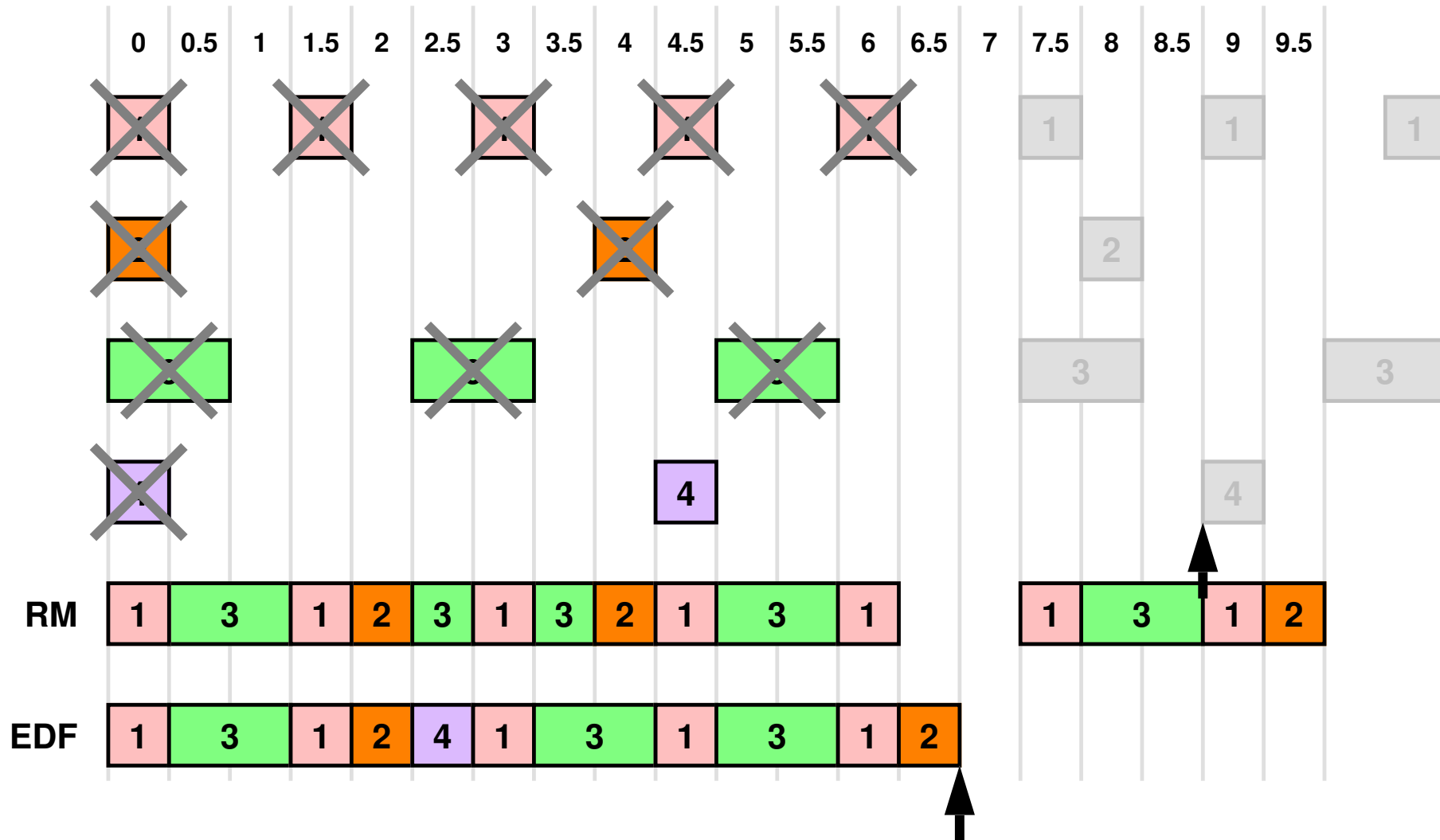
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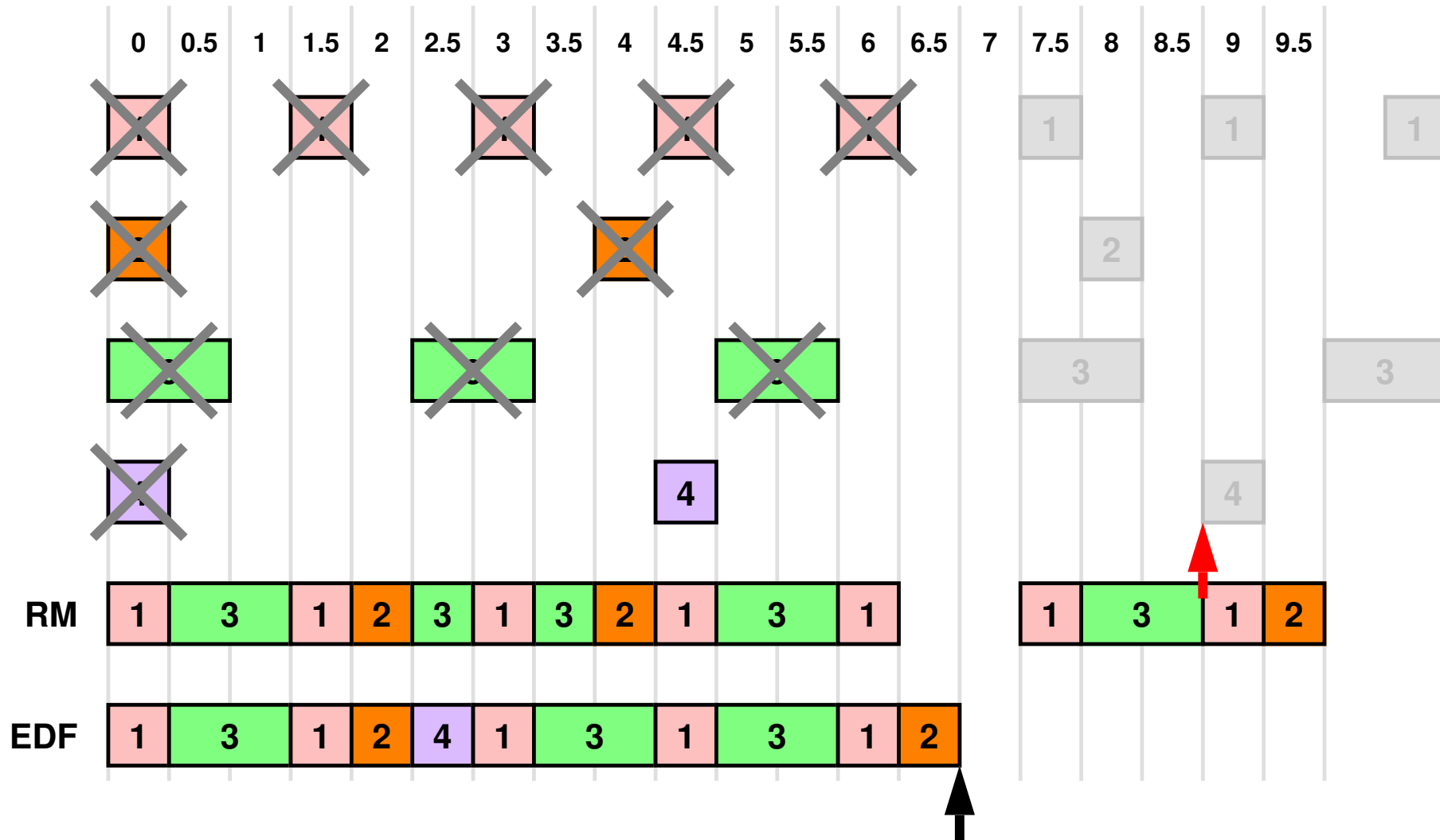
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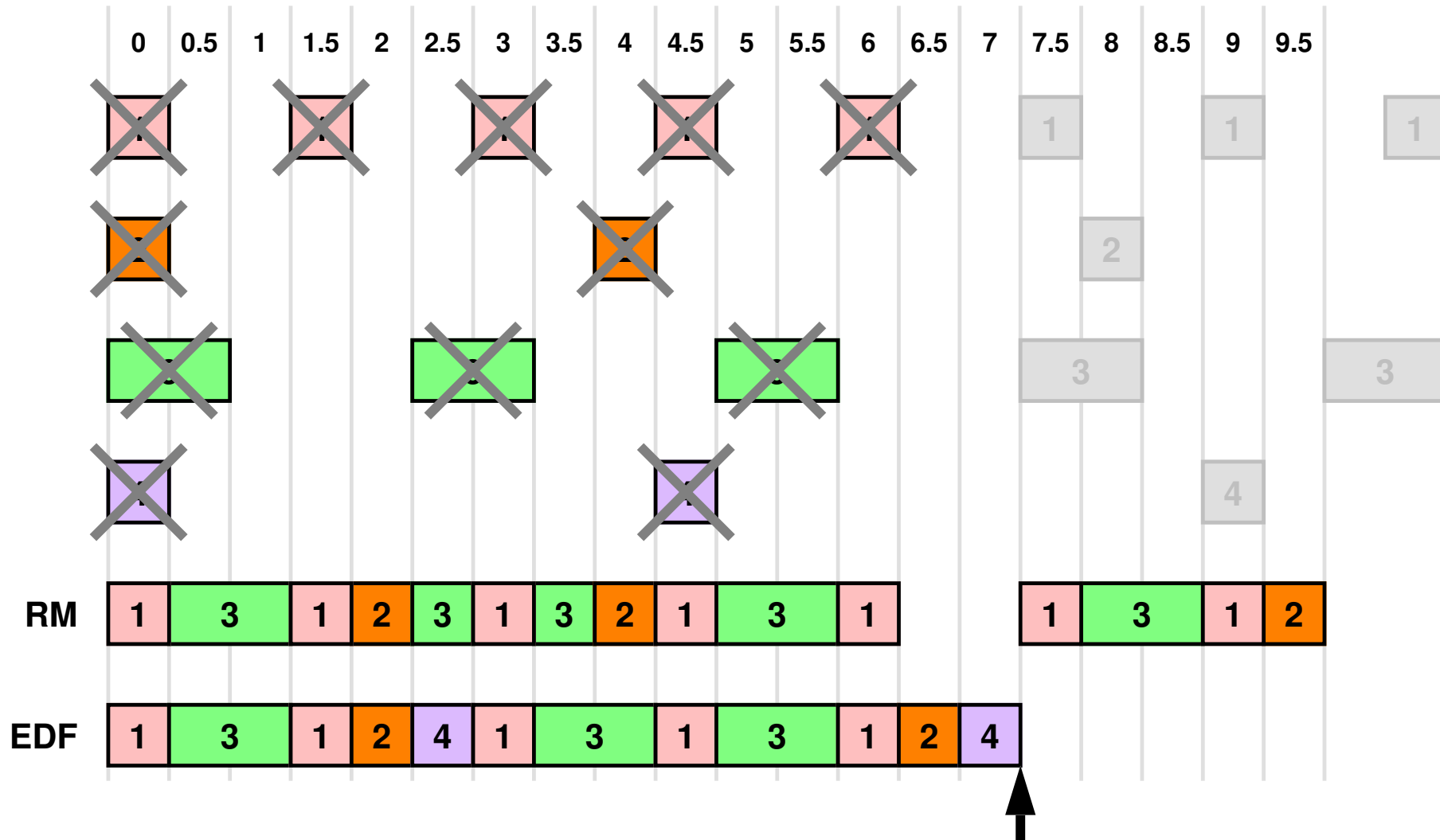
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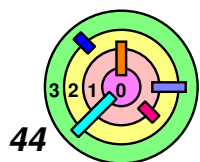
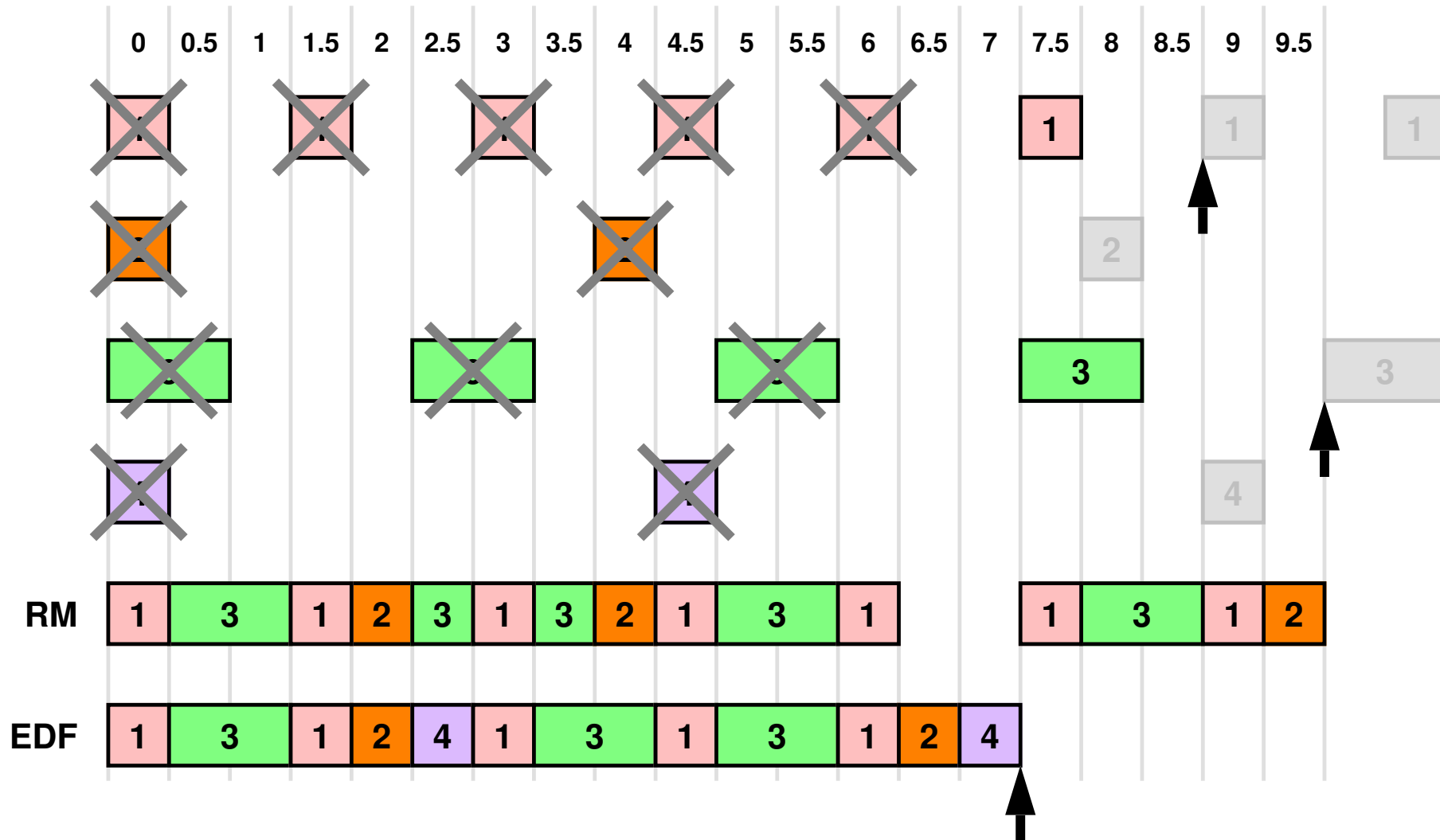
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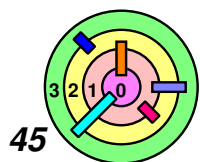
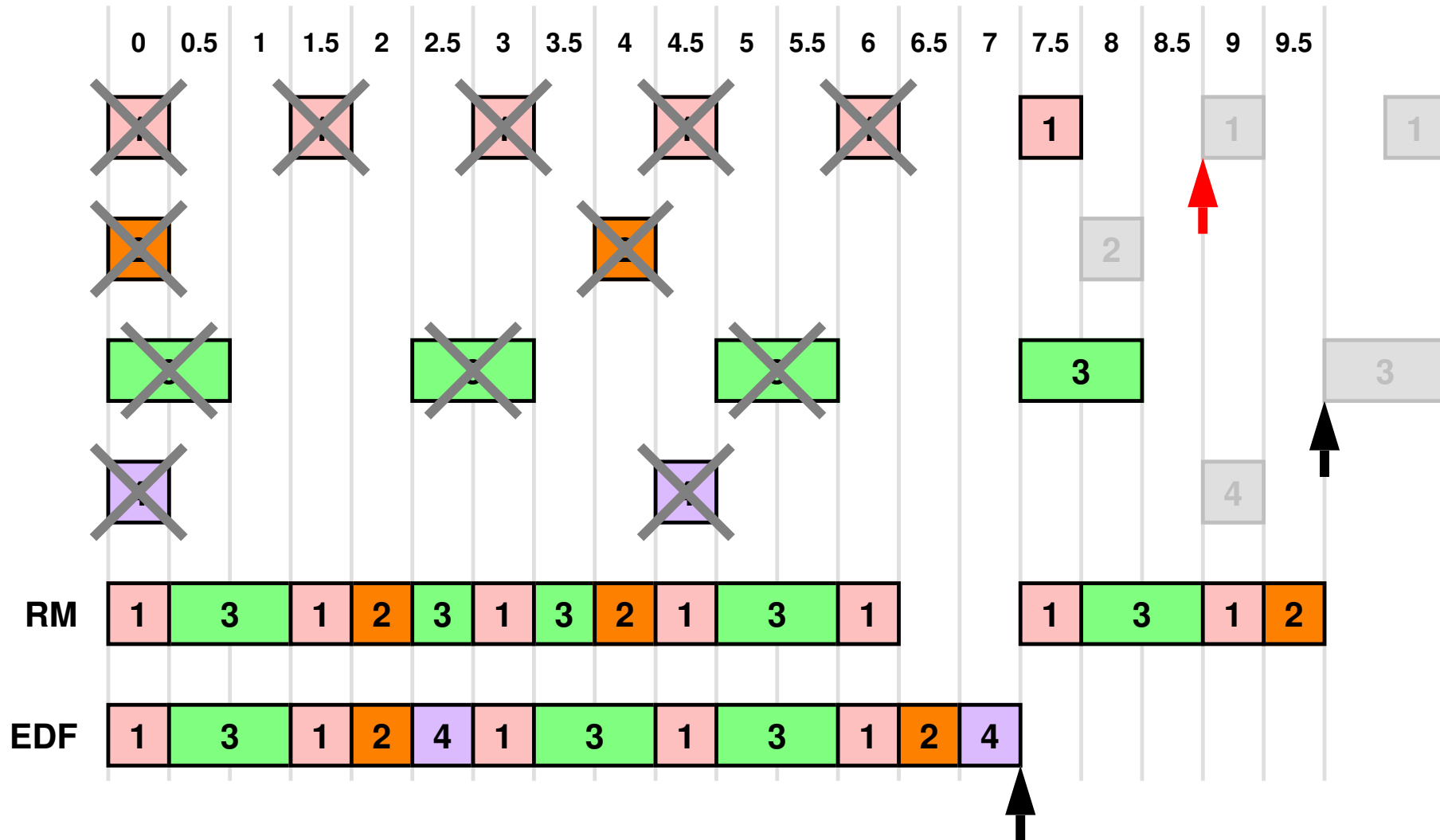
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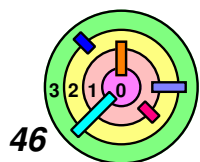
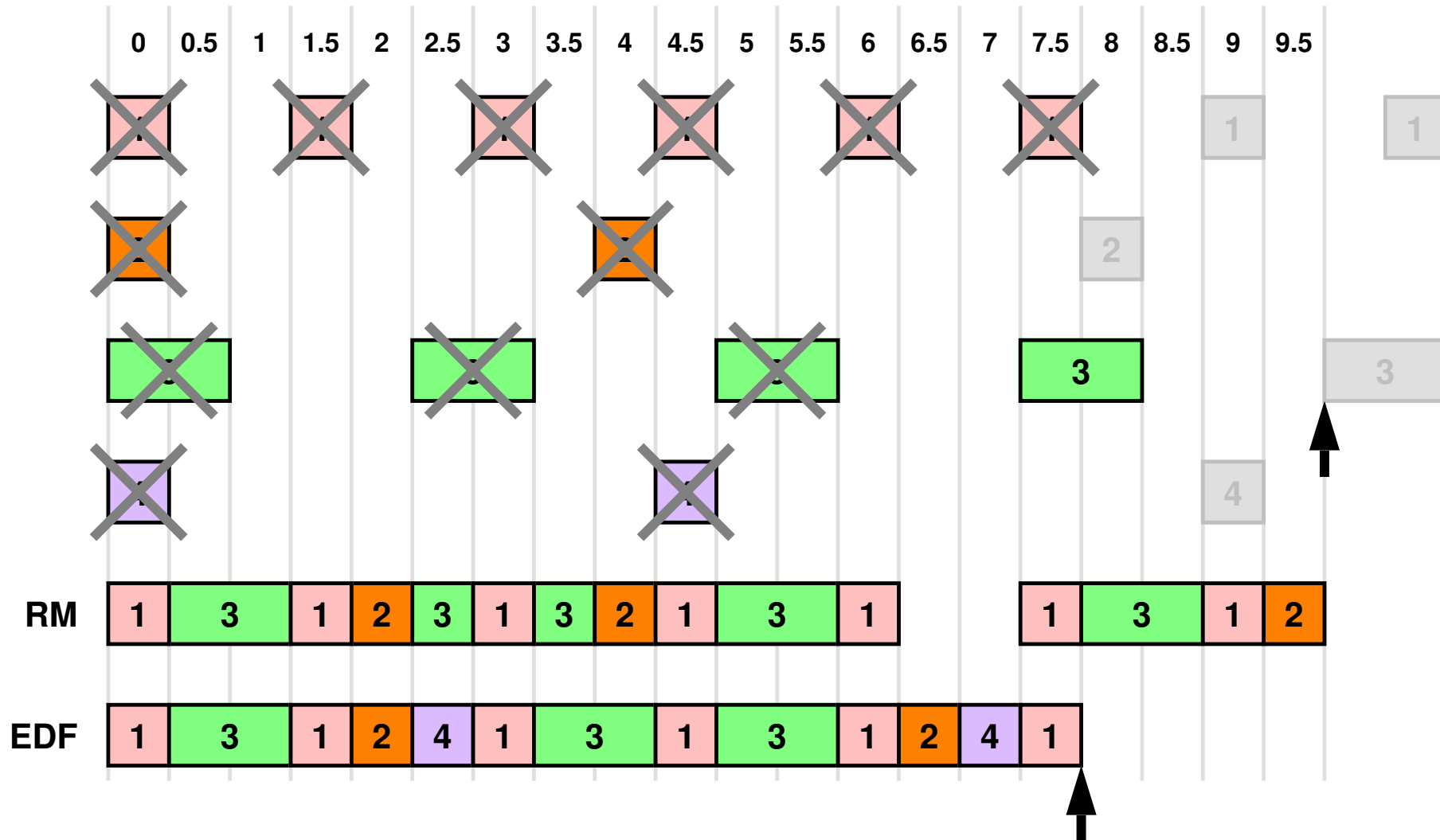
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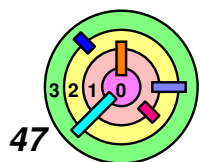
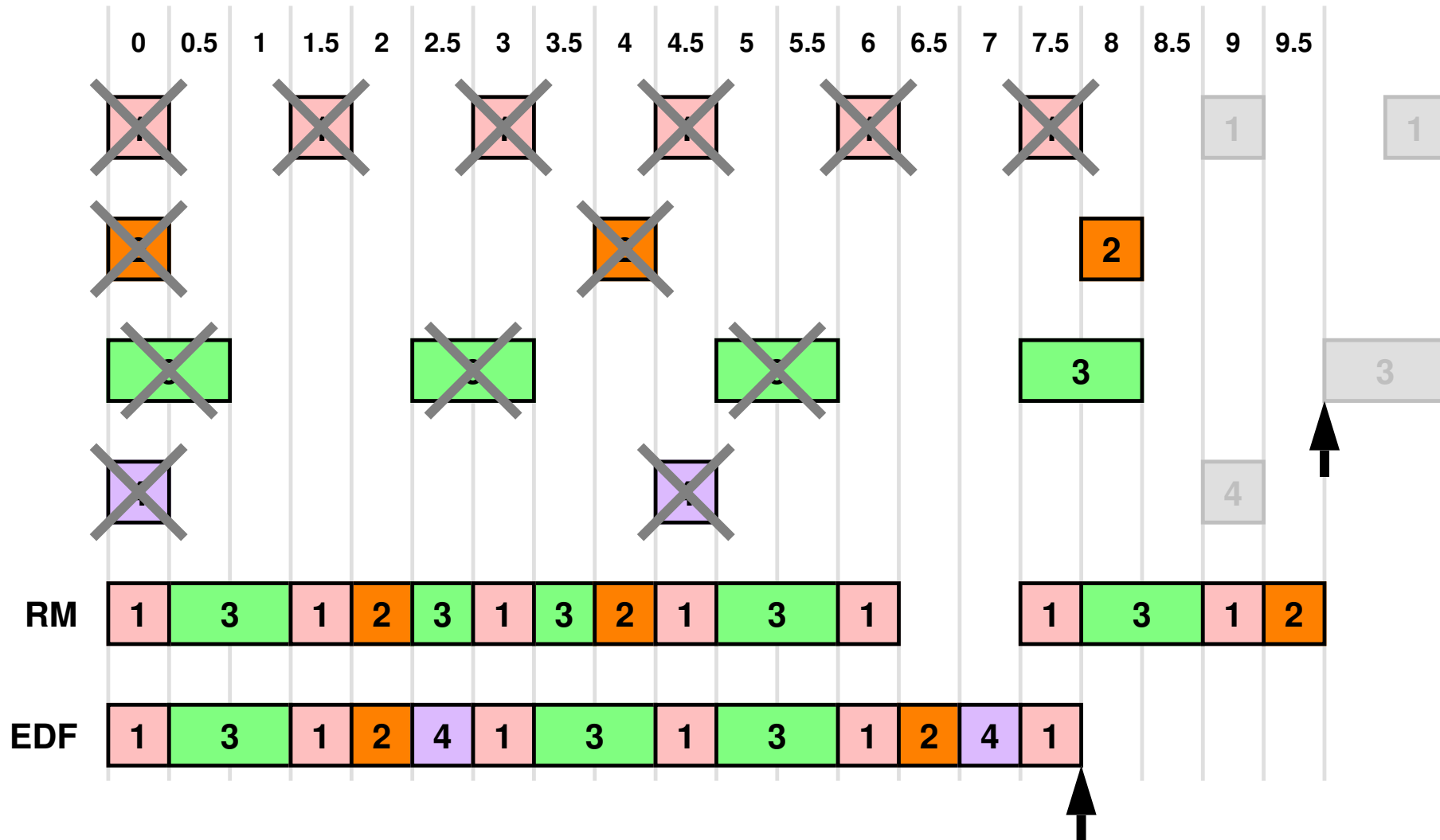
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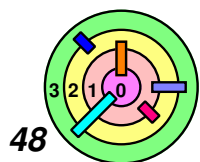
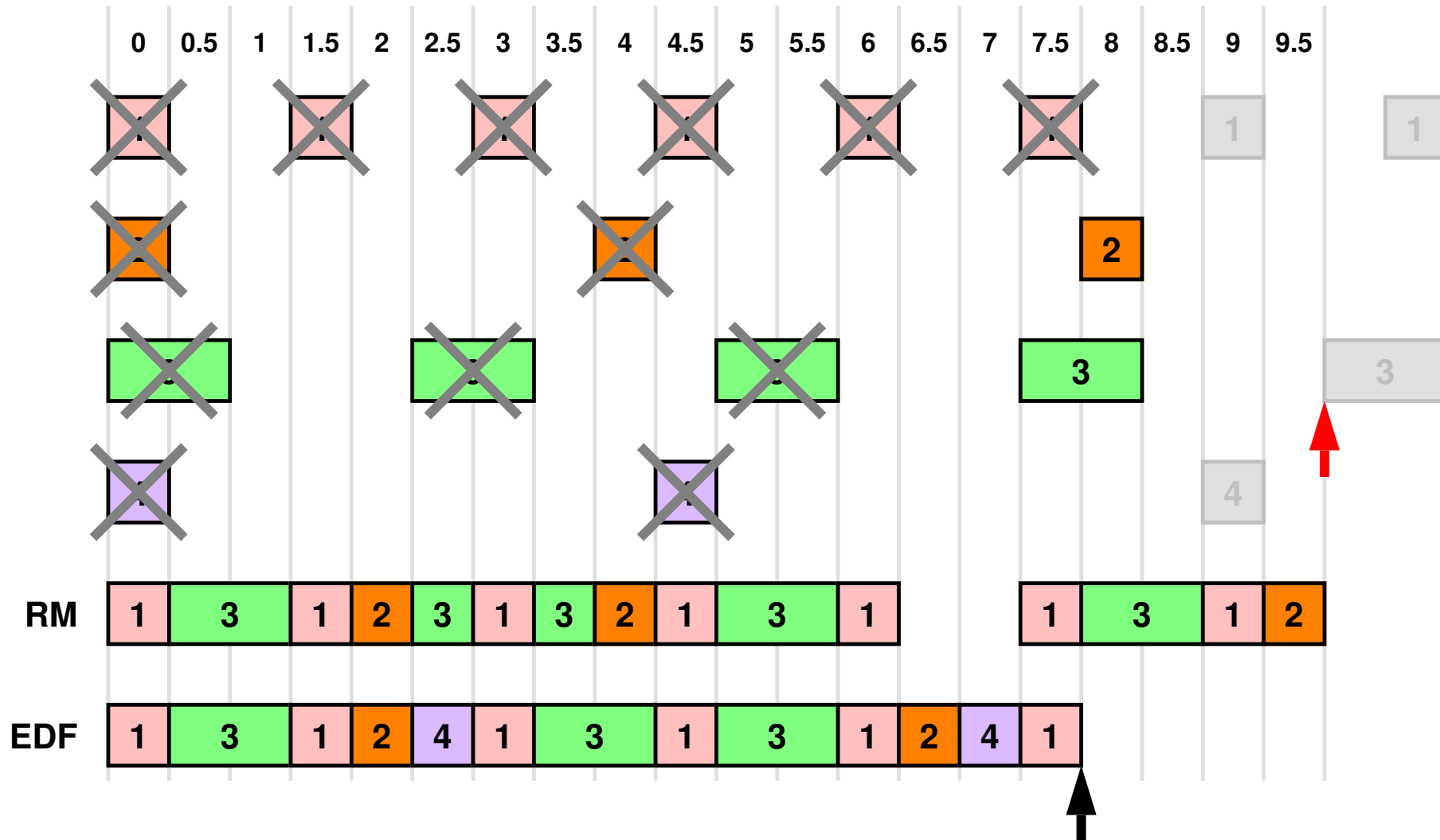
Earliest Deadline First



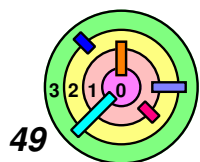
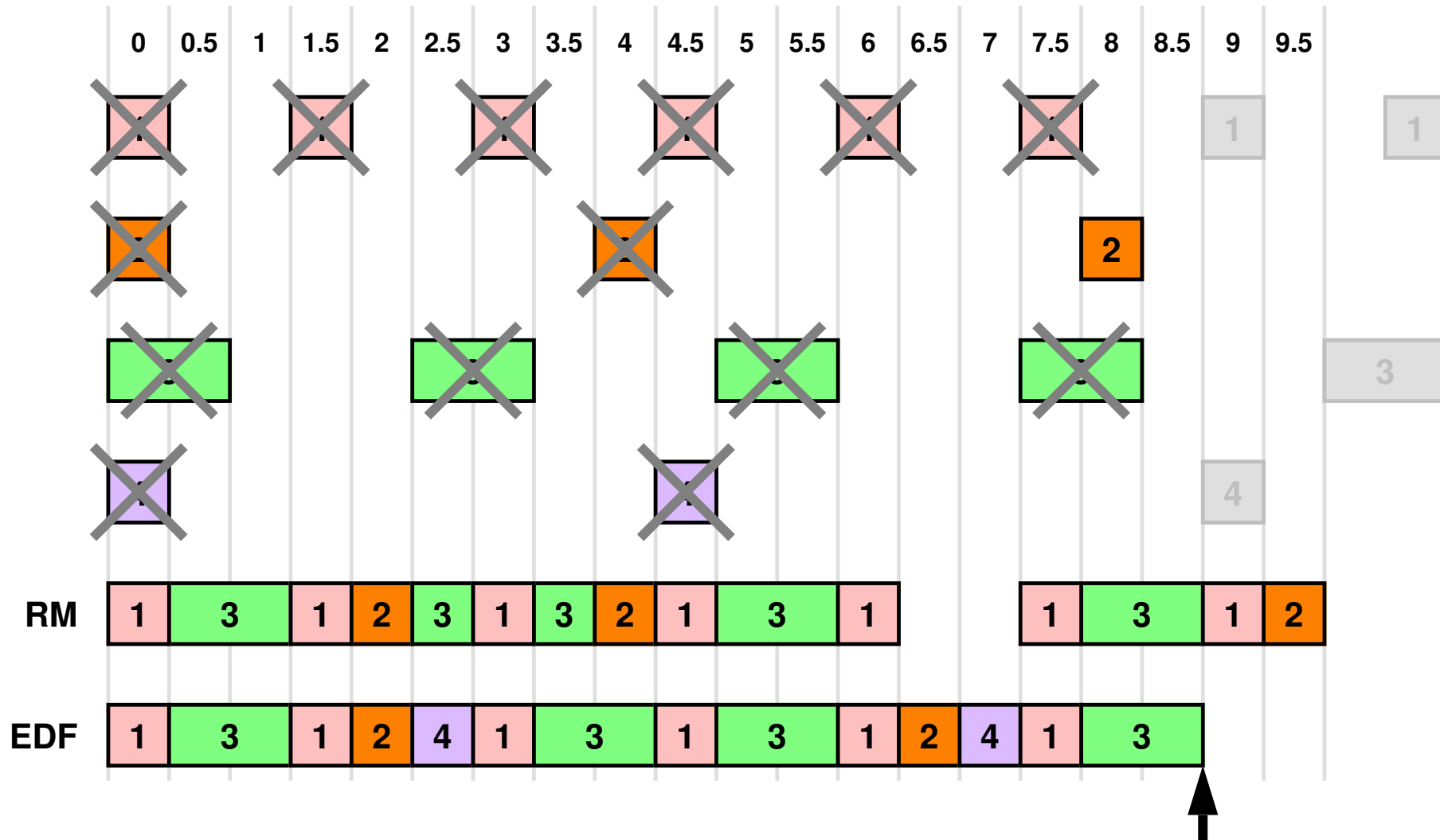
Earliest Deadline First

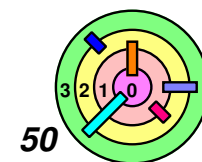


Earliest Deadline First

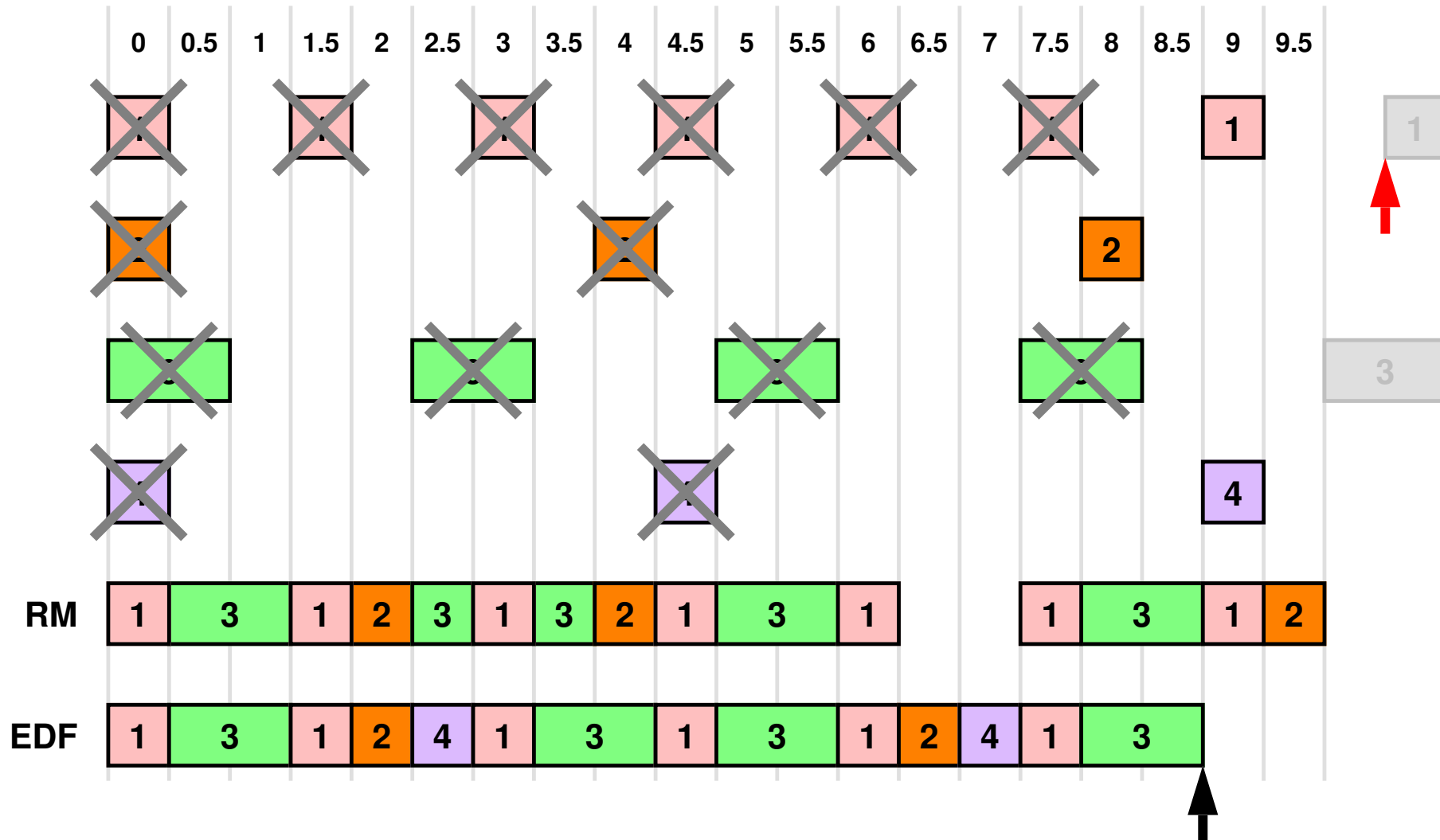


Earliest Deadline First

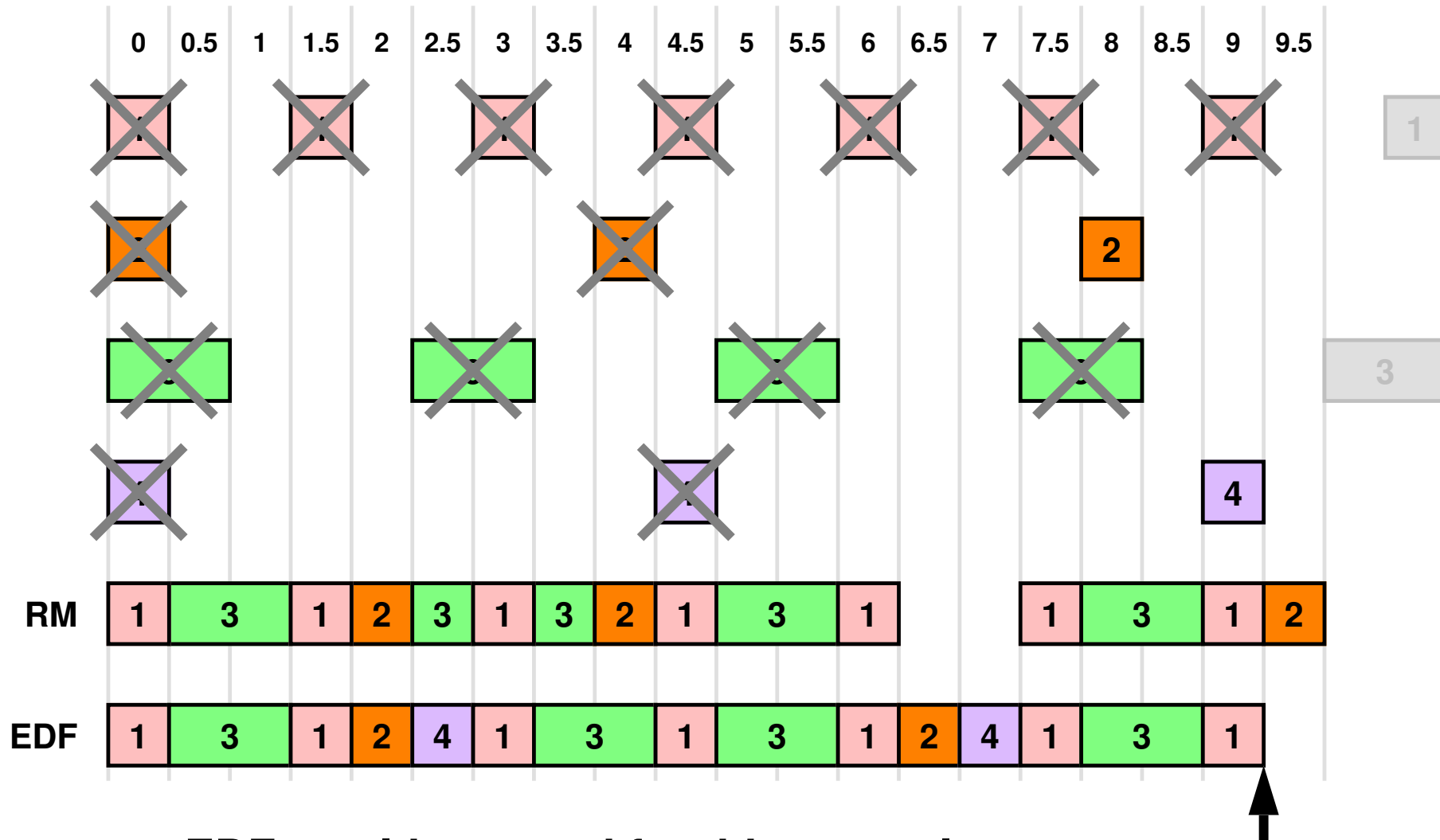




Earliest Deadline First

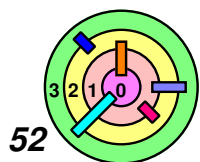


Earliest Deadline First

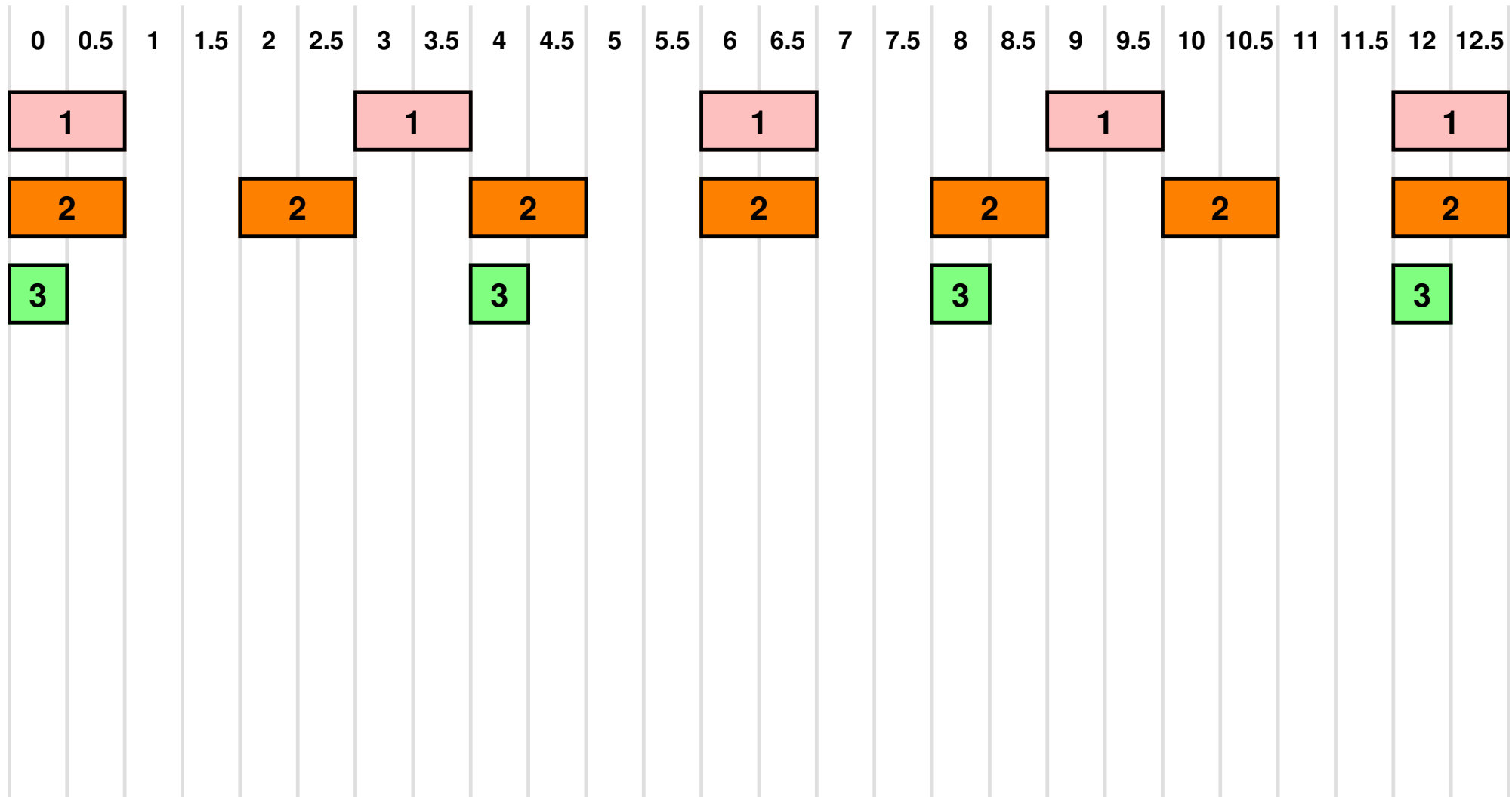


EDF would succeed for this example

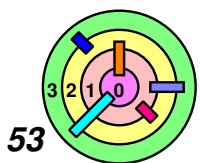
need to simulate till time = 180 to be certain



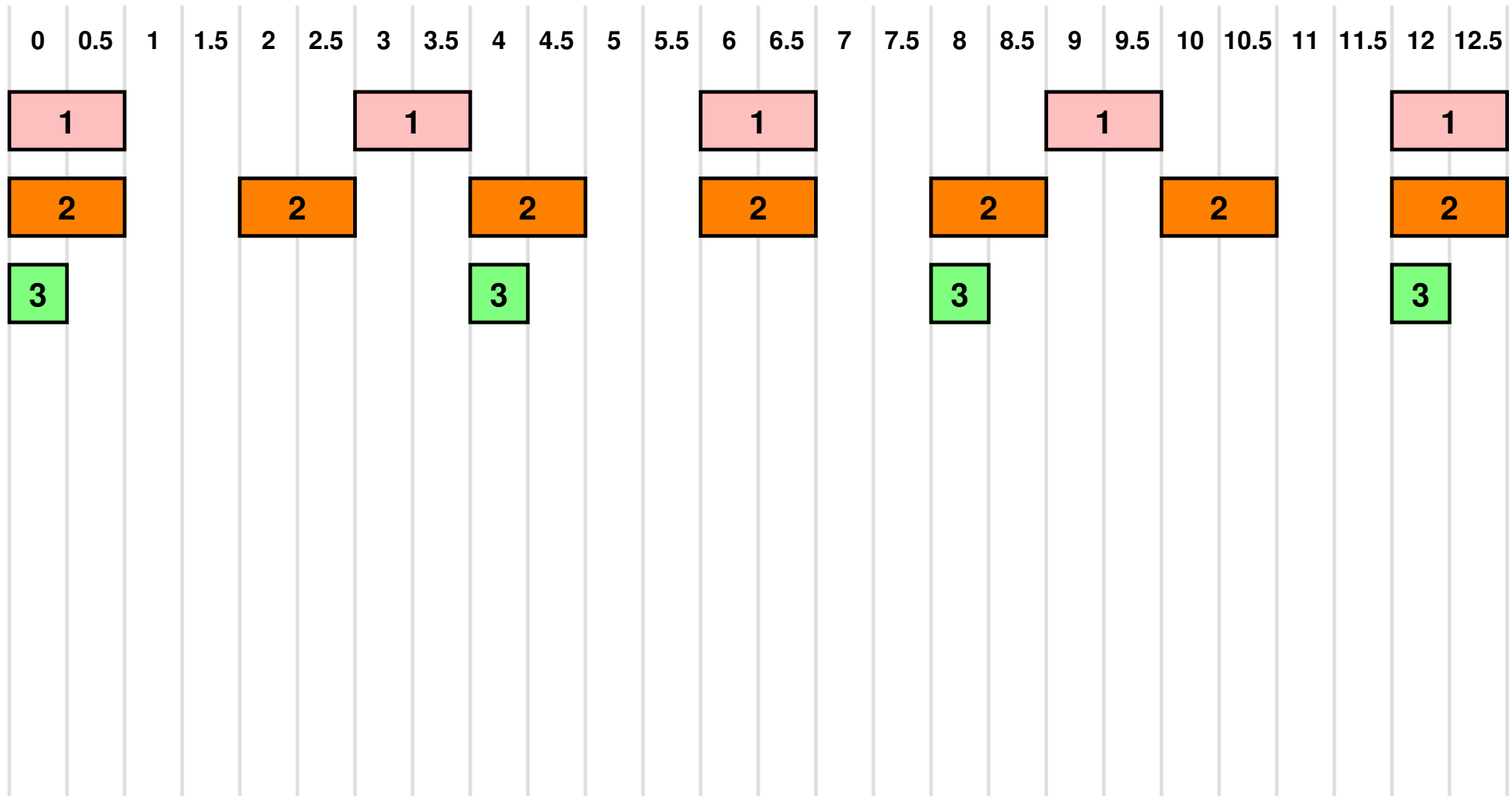
Phase Problems



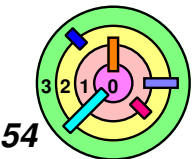
- if jobs do not start in-phase, rate-monotonic scheduling still may work



Phase Problems

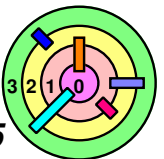


see "extra slides" at the end



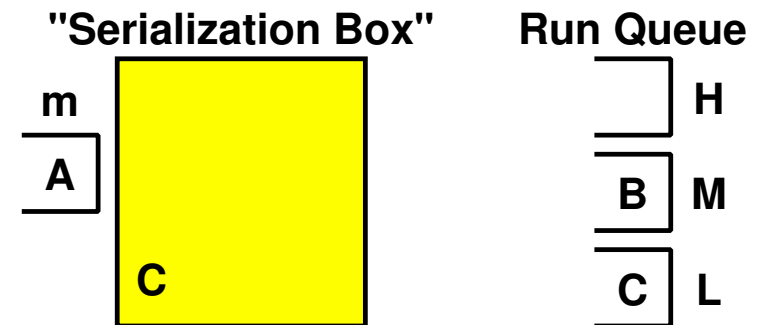
5.3 Scheduling

- ➡ **Goals**
- ➡ **Scheduling Algorithms**
- ➡ ***Implementation Issues***
- ➡ **Case Studies**



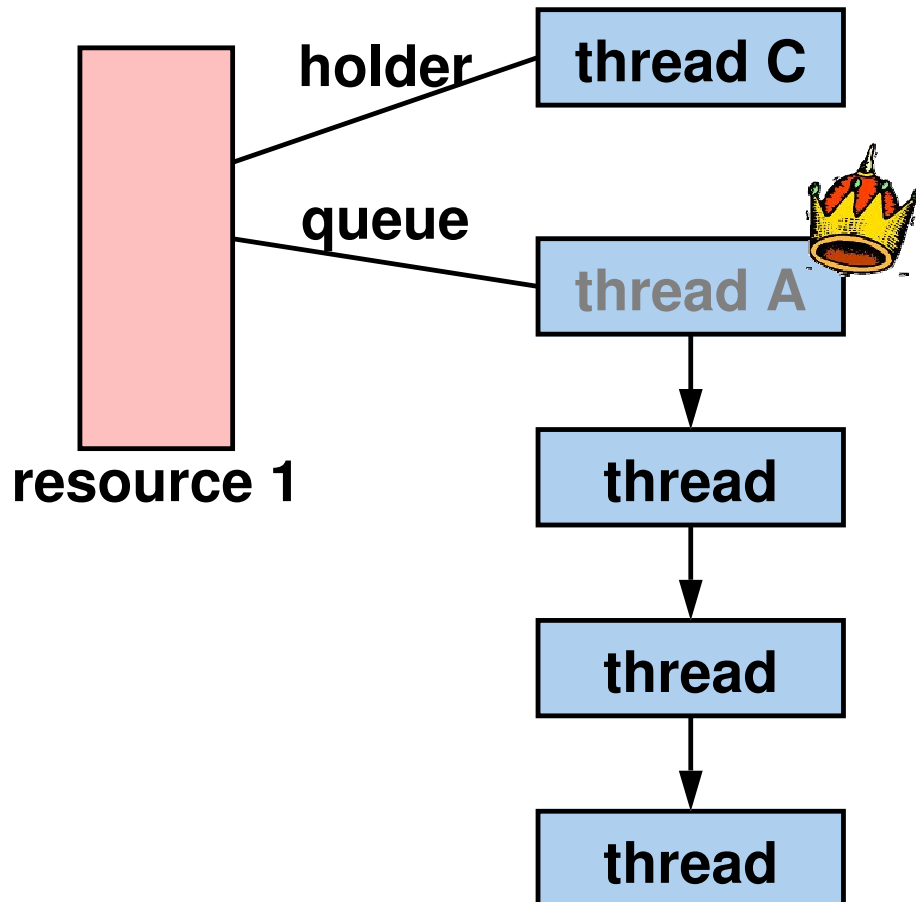
Priority Problem

- ➡ High-priority thread A blocks on mutex 1
- ➡ Low-priority thread C holds mutex 1
- ➡ Thread C cannot run because medium-priority thread B is running
- ➡ A is effectively waiting at C's priority
= *priority inversion*



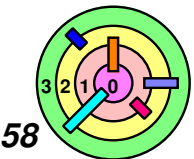
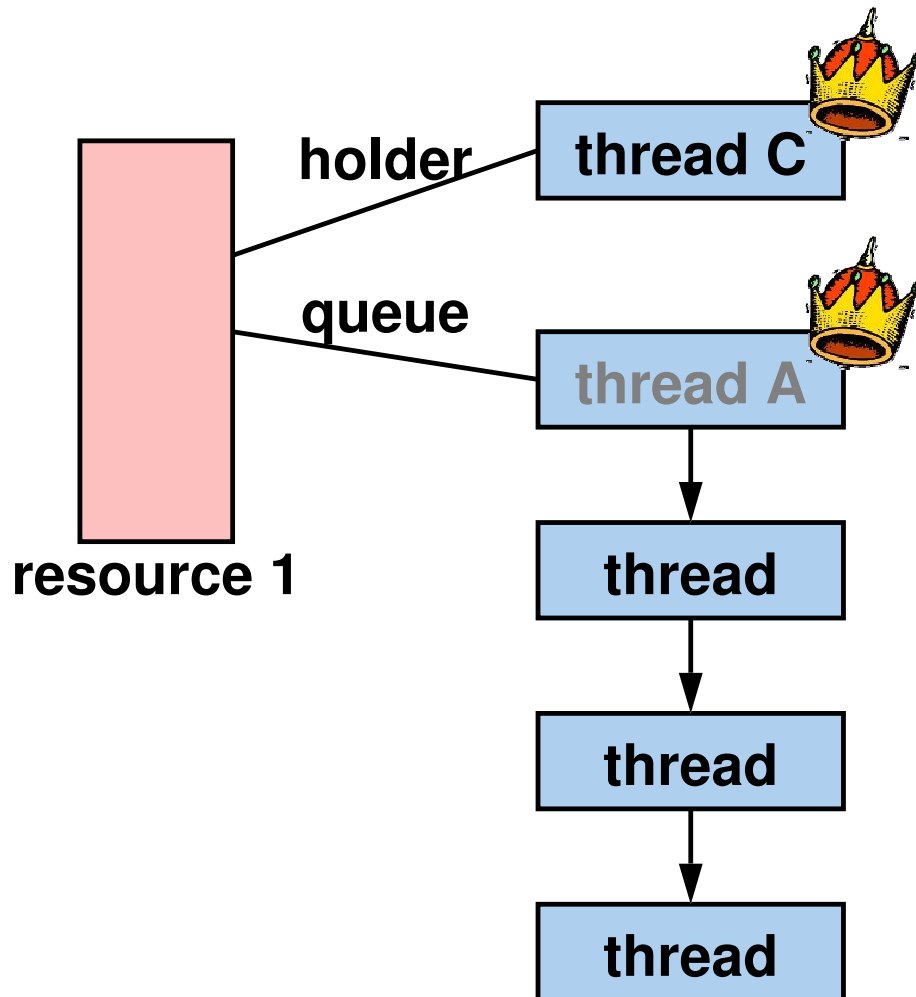
Priority Inheritance

➡ While A is waiting for resource held by C, it gives C its priority

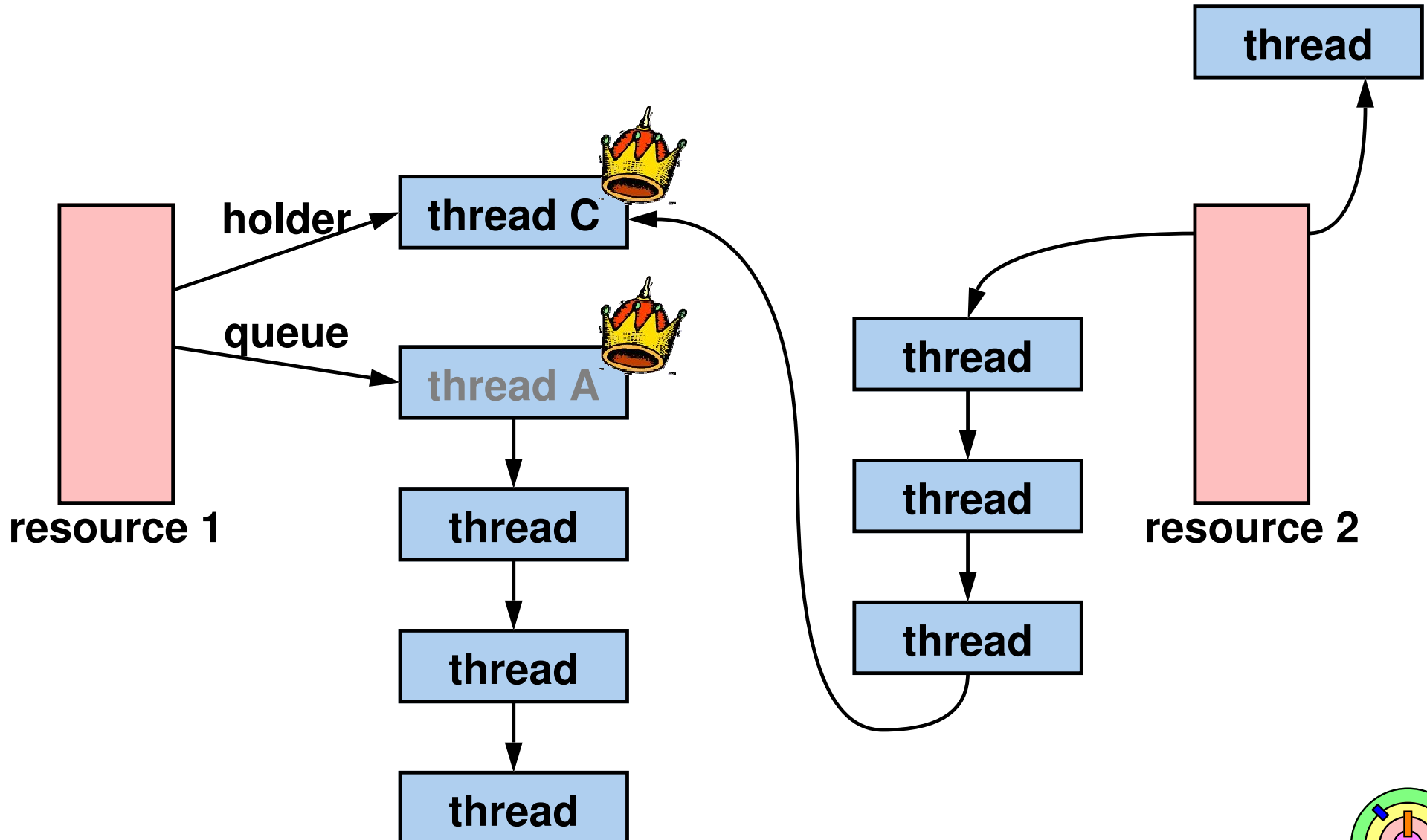


Priority Inheritance

➡ While A is waiting for resource held by C, it gives C its priority



Cacading Inheritance



Cacading Inheritance

