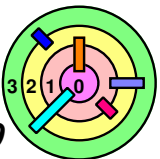


3.4 Linking & Loading

➡ *Static Linking & Loading*

➡ **Shared Libraries**



Remember This?

```

main:
→ pushl %ebp
  movl %esp, %ebp
  pushl %esi
  pushl %edi
  subl $8, %esp
  ...
  pushl $1
  movl -12(%ebp), %eax
  pushl %eax
  call sub
  addl $8, %esp
  movl %eax, -16(%ebp)
  ...
  addl $8, %esp
  movl $0, %eax
  popl %edi
  popl %esi
  movl %ebp, %esp
  popl %ebp
  ret

```

set up stack frame

push args

pop args; get result

set return value and restore frame



```

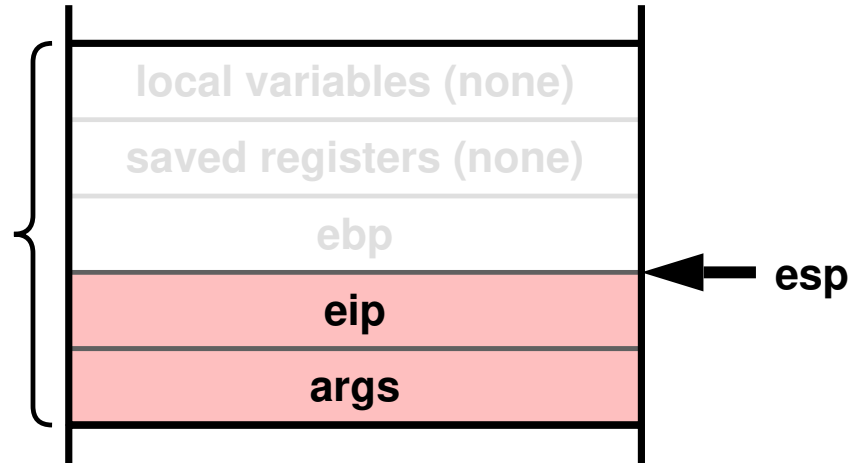
int main() {
    int i;
    int a;
    ...
    i = sub(a, 1);
    ...
    return(0);
}

```

Something Simpler

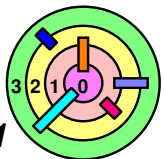
```
int main(int argc,
         char *[]) {
    return(argc);
}
```

stack frame
of main()



```
main:
→ pushl %ebp           } set up
  movl %esp, %ebp      } stack frame
  movl 8(%ebp), %eax
  movl %ebp, %esp      } set return
  popl %ebp            } value and
  ret                  } restore frame
```

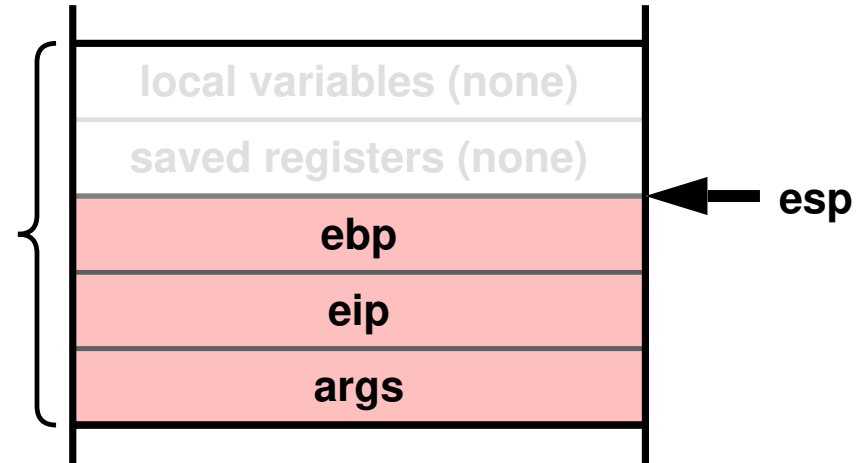
➡ Does location matter?



Something Simpler

```
int main(int argc,
        char *[]) {
    return(argc);
}
```

stack frame
of main()



main:

```

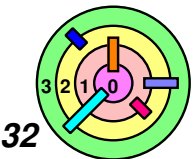
    pushl %ebp
    → movl %esp, %ebp
    movl 8(%ebp), %eax
    movl %ebp, %esp
    popl %ebp
    ret

```

set up
stack frame

set return
value and
restore frame

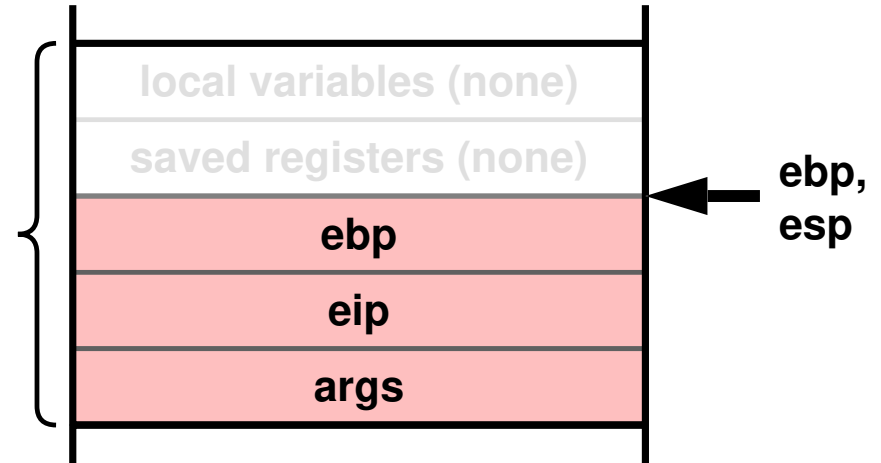
➡ Does location matter?



Something Simpler

```
int main(int argc,
         char *[]) {
    return(argc);
}
```

stack frame
of main()

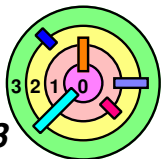


```
main:
    pushl %ebp
    movl %esp, %ebp
    → movl 8(%ebp), %eax
    movl %ebp, %esp
    popl %ebp
    ret
```

set up
stack frame

set return
value and
restore frame

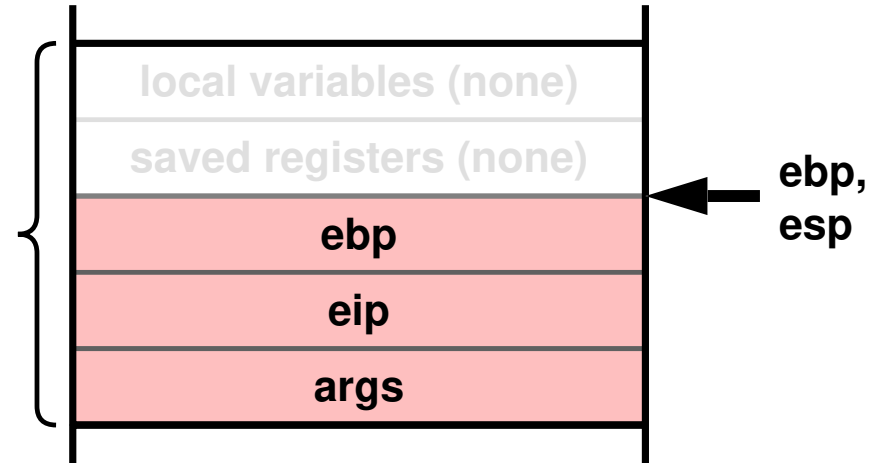
➡ Does location matter?



Something Simpler

```
int main(int argc,
         char *[]) {
    return(argc);
}
```

stack frame
of main()

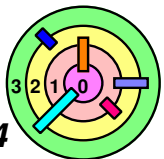


```
main:
    pushl %ebp
    movl %esp, %ebp
    movl 8(%ebp), %eax
    → movl %ebp, %esp
    popl %ebp
    ret
```

set up
stack frame

set return
value and
restore frame

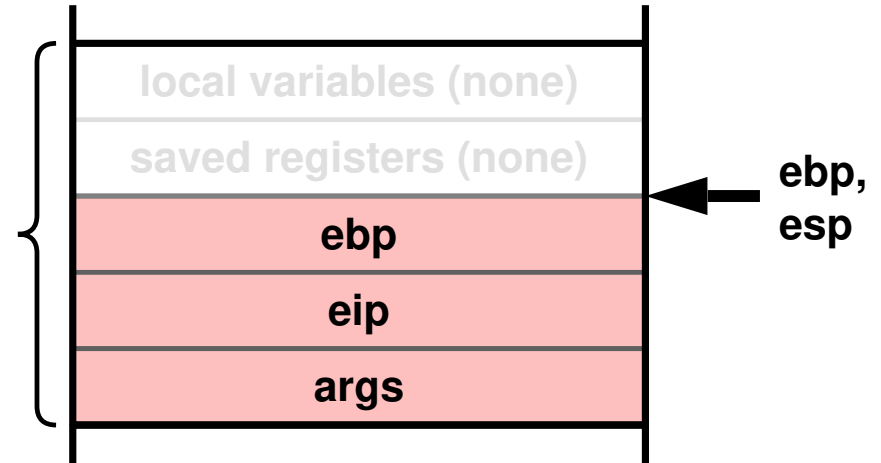
➡ Does location matter?



Something Simpler

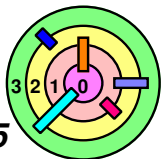
```
int main(int argc,
        char *[]) {
    return(argc);
}
```

stack frame
of main()



```
main:
    pushl %ebp          } set up
    movl %esp, %ebp     } stack frame
    movl 8(%ebp), %eax
    movl %ebp, %esp
    → popl %ebp          } set return
    ret                 } value and
                        } restore frame
```

➡ Does location matter?

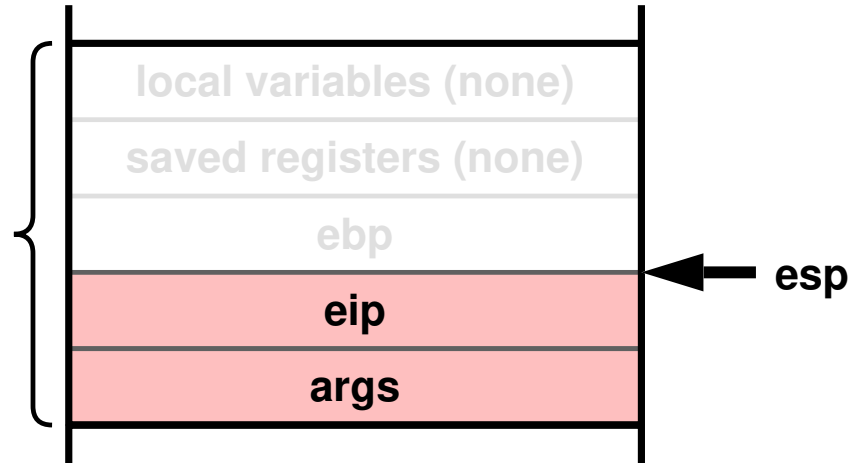


Something Simpler

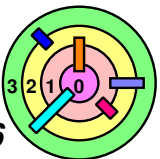
```
int main(int argc,
        char *[]) {
    return(argc);
}
```

```
main:
    pushl %ebp          } set up
    movl %esp, %ebp     } stack frame
    movl 8(%ebp), %eax
    movl %ebp, %esp      } set return
    popl %ebp            } value and
    ret                  } restore frame
```

stack frame
of main()



➡ Does location matter?

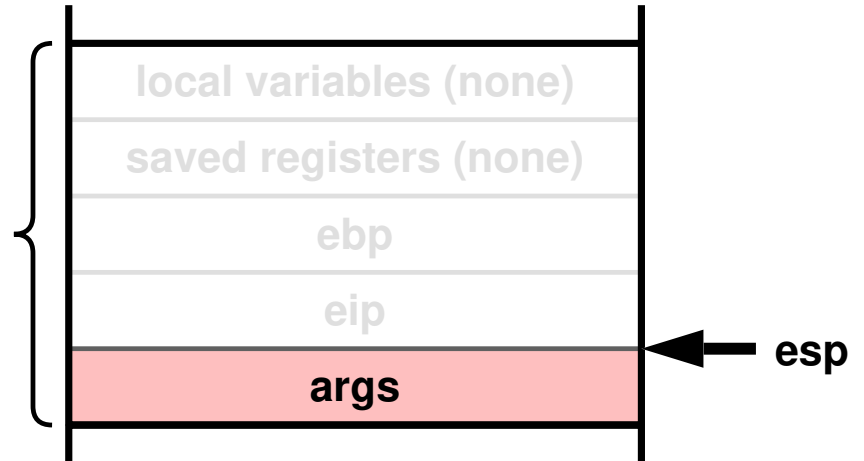


Something Simpler

```
int main(int argc,
        char *[]) {
    return(argc);
}
```

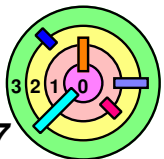
```
main:
    pushl %ebp          } set up
    movl %esp, %ebp     } stack frame
    movl 8(%ebp), %eax
    movl %ebp, %esp     } set return
    popl %ebp           } value and
    ret                 } restore frame
```

stack frame
of main()



Does location matter?

- if everything can be accessed relative to the *frame pointer*, then you don't need to know the actual address of an object
 - just use *relative-addresses* to *access variables*
 - the *code* is also *location-independent*



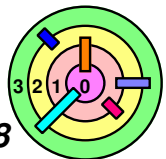
Location Matters ...

```
int X = 6;  
int *aX = &X;
```

```
void subr(int i) {  
    printf("i = %d\n", i);  
}
```

```
int main( ) {  
    void subr(int);  
    int y = X;  
    subr(y);  
    return(0);  
}
```

- ➡ Why does it matter here?
- ➡ need to put the *address* of x into aX
 - *what* is the address of x?
 - ◆ *when* do you know?
 - remember, both x and aX are in the *data segment*
 - *who* would put the actual value into aX?
 - ➡ also need to put the *address* of subr into main()



Coping



Done in two steps

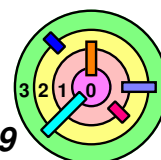
- **compiler** figure out *how many bytes* each object is
 - assign them temporary locations
- **linker** figures out *where* each object is and lays out the entire address space
 - functions, global variables, and more



Relocation

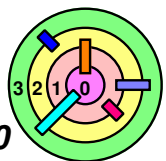
- modify internal references in memory depending on where module is expected to be *loaded*
 - one of the **exec** system calls loads a program into memory
 - ◆ everything is laid out carefully in memory
- modules requiring relocation are said to be *relocatable*
- the act of modifying such a module to *resolve these references* is called *relocation*
- the program that performs relocation is called a *linker*

textbook
is wrong



Linker

- ➡ Two main functions of a *linker*
 - 1) *relocation*
 - 2) *symbol resolution*
 - ◆ find unresolved symbols and figure out how to resolve them
- ➡ A *loader* loads a program into memory
 - ▬ "unfolds" a program from disk into memory and "transfer control" to it (i.e., "executes" it)
 - ▬ a "relocating loader" may perform additional relocation
 - but that's *dynamic* linking



A Slight Revision

```
extern int X;
int *aX = &X;

int main( ) {
    void subr(int);
    int y = *aX;
    subr(y);
    return(0);
}
```

main.c

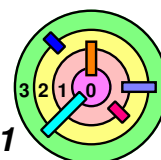
```
#include <stdio.h>
int X;

void subr(int i) {
    printf("i = %d\n", i);
}
```

subr.c

```
% gcc -o prog main.c subr.c
```

- main.c is compiled into main.o
 - subr.c is compiled into subr.o
- } relocatable modules
- ld is then invoked to combine them into prog
 - ld knows where to find printf()
 - prog can be loaded into memory through one of the **exec** system calls



A Slight Revision

```
extern int X;
int *aX = &X;

int main( ) {
    void subr(int);
    int y = *aX;
    subr(y);
    return(0);
}
```

main.c

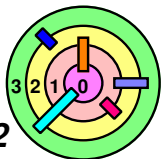
```
#include <stdio.h>
int X;

void subr(int i) {
    printf("i = %d\n", i);
}
```

subr.c

```
% gcc -o prog main.c subr.c
```

- how does ld decides what needs to be done?
- main.c contains undefined references to x and subr()
 - instructions for doing this are provided in main.o
- later on, when the actual locations for these are determined, ld will modify them when main.o is *copied* into prog



A Slight Revision

```
extern int X;  
int *aX = &X;  
  
int main( ) {  
    void subr(int);  
    int y = *aX;  
    subr(y);  
    return(0);  
}
```

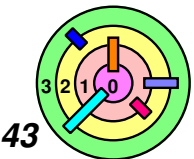
main.c

```
#include <stdio.h>  
int X;  
  
void subr(int i) {  
    printf("i = %d\n", i);  
}
```

subr.c

```
% gcc -o prog main.c subr.c
```

- main.o must contain a list of external symbols, along with their types, and instructions for updating this code



Can Also Use A Header File

```
#include "subr.c";
int *aX = &X;
```

```
int main( ) {
    int y = *aX;
    subr(y);
    return(0);
}
```

main.c

```
extern int X;
void subr(int);
```

subr.h

```
#include <stdio.h>
int X;

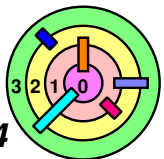
void subr(int i) {
    printf("i = %d\n", i);
}
```

subr.c



A **header** file typically contains

- declaration/definition of **data structures**
- declaration of **exported symbols**

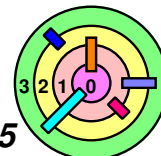


main.s

Offset	Op	Arg
0:	.data	; what follows is initialized data
0:	.globl	aX ; aX is global: it may be used ; by others
0:	aX:	
0:	.long	X
4:		
0:	.text	; offset restarts; what follows is ; text (read-only code)
0:	.globl	main
0:	main:	
0:	pushl	%ebp ; save the frame pointer
1:	movl	%esp,%ebp ; point to
3:	subl	\$4,%esp ; make space
6:	movl	aX,%eax ; put conten
11:	movl	(%eax),%eax ; put *X
13:	movl	%eax,-4(%ebp) ; stor
16:	pushl	-4(%ebp) ; push y on
19:	call	subr
24:	addl	\$4,%esp ; remove y f
27:	movl	\$0,%eax ; set return
31:	movl	%ebp, %esp ; restore
33:	popl	%ebp ; pop frame pointer
35:	ret	

```
extern int X;
int *aX = &X;

int main( ) {
    void subr(int);
    int y = *aX;
    subr(y);
    return(0);
}
```



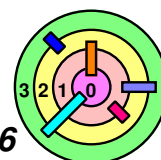
main.s

Offset	Op	Arg
0:	.data	; what follows is initialized data
0:	.globl	aX ; aX is global: it may be used ; by others
0:	aX:	
0:	.long	X
4:		
0:	.text	; offset restarts; w ; text (read-only co
0:	.globl	main
0:	main:	
0:	pushl	%ebp ; save the frame pointer
1:	movl	%esp,%ebp ; point to
3:	subl	\$4,%esp ; make space
6:	movl	aX,%eax ; put conten
11:	movl	(%eax),%eax ; put *X
13:	movl	%eax,-4(%ebp) ; stor
16:	pushl	-4(%ebp) ; push y on
19:	call	subr
24:	addl	\$4,%esp ; remove y f
27:	movl	\$0,%eax ; set return
31:	movl	%ebp, %esp ; restore
33:	popl	%ebp ; pop frame pointer
35:	ret	

What follows goes into the *data* segment

```
extern int X;
int *aX = &X;

int main( ) {
    void subr(int);
    int y = *aX;
    subr(y);
    return(0);
}
```



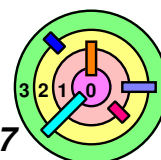
main.s

Offset	Op	Arg
0:	.data	; what follows is initialized data
0:	.globl	aX ; aX is global: it may be used ; by others
0:	aX:	
0:	.long	X
4:		
→ 0:	.text	; offset restarts; w ; text (read-only co
0:	.globl	main
0:	main:	
0:	pushl	%ebp ; save the frame pointer
1:	movl	%esp,%ebp ; point to
3:	subl	\$4,%esp ; make space
6:	movl	aX,%eax ; put conten
11:	movl	(%eax),%eax ; put *X
13:	movl	%eax,-4(%ebp) ; stor
16:	pushl	-4(%ebp) ; push y on
19:	call	subr
24:	addl	\$4,%esp ; remove y f
27:	movl	\$0,%eax ; set return
31:	movl	%ebp, %esp ; restore
33:	popl	%ebp ; pop frame pointer
35:	ret	

What follows goes into the **text** segment

```
extern int X;
int *aX = &X;

int main( ) {
    void subr(int);
    int y = *aX;
    subr(y);
    return(0);
}
```



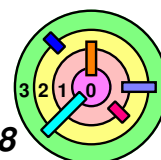
main.s

Offset	Op	Arg
→ 0:	.data	; what follows is initialized data
0:	.globl	aX ; aX is global: it may be used ; by others
0:	aX:	
0:	.long	X
4:		
→ 0:	.text	; offset restarts; w ; text (read-only co
0:	.globl	main
0:	main:	
0:	pushl	%ebp ; save the frame pointer
1:	movl	%esp,%ebp ; point to
3:	subl	\$4,%esp ; make space
6:	movl	aX,%eax ; put conten
11:	movl	(%eax),%eax ; put *X
13:	movl	%eax,-4(%ebp) ; stor
16:	pushl	-4(%ebp) ; push y on
19:	call	subr
24:	addl	\$4,%esp ; remove y f
27:	movl	\$0,%eax ; set return
31:	movl	%ebp, %esp ; restore
33:	popl	%ebp ; pop frame pointer
35:	ret	

offset got restarted because
segments are relocatable

```
extern int X;
int *aX = &X;

int main( ) {
    void subr(int);
    int y = *aX;
    subr(y);
    return(0);
}
```



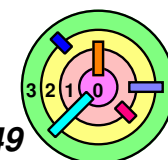
main.s

Offset	Op	Arg
0:	.data	; what follows is initialized data
→ 0:	.globl	aX ; aX is global: it may be used ; by others
0:	aX:	
0:	.long	X
4:		
0:	.text	; offset restarts; w ; text (read-only co
→ 0:	.globl	main
0:	main:	
0:	pushl	%ebp ; save the frame pointer
1:	movl	%esp,%ebp ; point to
3:	subl	\$4,%esp ; make space
6:	movl	aX,%eax ; put conten
11:	movl	(%eax),%eax ; put *X
13:	movl	%eax,-4(%ebp) ; stor
16:	pushl	-4(%ebp) ; push y on
19:	call	subr
24:	addl	\$4,%esp ; remove y f
27:	movl	\$0,%eax ; set return
31:	movl	%ebp, %esp ; restore
33:	popl	%ebp ; pop frame pointer
35:	ret	

.global directive means that the symbol mentioned is defined *here* and is *exported*
 ➡ i.e., can be referenced by other modules
 ➡ aX and main are global

```
extern int X;
int *aX = &X;

int main( ) {
    void subr(int);
    int y = *aX;
    subr(y);
    return(0);
}
```



main.s

Offset	Op	Arg
0:	.data	; what follows is initialized data
0:	.globl	aX ; aX is global: it may be used ; by others
0:	aX:	
0:	.long	X
4:		
0:	.text	; offset restarts; w ; text (read-only co
0:	.globl	main
0:	main:	
0:	pushl	%ebp ; save the frame pointer
1:	movl	%esp,%ebp ; point to
3:	subl	\$4,%esp ; make space
6:	movl	aX,%eax ; put conten
11:	movl	(%eax),%eax ; put *X
13:	movl	%eax,-4(%ebp) ; stor
16:	pushl	-4(%ebp) ; push y on
19:	call	subr
24:	addl	\$4,%esp ; remove y f
27:	movl	\$0,%eax ; set return
31:	movl	%ebp, %esp ; restore
33:	popl	%ebp ; pop frame pointer
35:	ret	

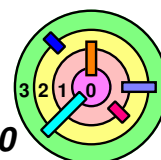
aX is 4 bytes long and put
the value of x here

☞ x here is an *address*

☞ x will remain *unresolved*

```
extern int X;
int *aX = &X;
```

```
int main( ) {
    void subr(int);
    int y = *aX;
    subr(y);
    return(0);
}
```



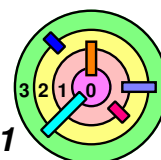
main.s

Offset	Op	Arg
0:	.data	; what follows is initialized data
0:	.globl	aX ; aX is global: it may be used ; by others
0:	aX:	
0:	.long	X
4:		
0:	.text	; offset restarts; w ; text (read-only co
0:	.globl	main
0:	main:	
0:	pushl	%ebp ; save the frame pointer
1:	movl	%esp,%ebp ; point to
3:	subl	\$4,%esp ; make space
6:	movl	aX,%eax ; put conten
11:	movl	(%eax),%eax ; put *X
13:	movl	%eax,-4(%ebp) ; stor
16:	pushl	-4(%ebp) ; push y on
19:	call	subr
24:	addl	\$4,%esp ; remove y f
27:	movl	\$0,%eax ; set return
31:	movl	%ebp, %esp ; restore
33:	popl	%ebp ; pop frame pointer
35:	ret	

these 3 places require
relocation

```
extern int X;
int *aX = &X;

int main( ) {
    void subr(int);
    int y = *aX;
    subr(y);
    return(0);
}
```



main.s

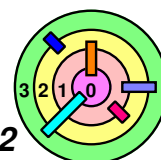
Offset	Op	Arg
0:	.data	; what follows is initialized data
0:	.globl	aX ; aX is global: it may be used ; by others
0:	aX:	
0:	.long	X
4:		
0:	.text	; offset restarts; w ; text (read-only co
0:	.globl	main
0:	main:	
0:	pushl	%ebp ; save the frame pointer
1:	movl	%esp,%ebp ; point to
3:	subl	\$4,%esp ; make space
6:	movl	aX,%eax ; put conten
11:	movl	(%eax),%eax ; put *X
13:	movl	%eax,-4(%ebp) ; stor
16:	pushl	-4(%ebp) ; push y on
19:	call	subr
24:	addl	\$4,%esp ; remove y f
27:	movl	\$0,%eax ; set return
31:	movl	%ebp, %esp ; restore
33:	popl	%ebp ; pop frame pointer
35:	ret	

this call is a PC-relative call

- what's stored at offset 20 is not the absolute address of subr, but a relative address
- this is just how x86 works

```
extern int X;
int *aX = &X;

int main( ) {
    void subr(int);
    int y = *aX;
    subr(y);
    return(0);
}
```



subr.s

Offset	Op	Arg
0:	.data	; what follows is initialized data
0:	printfarg:	
0:	.string	"i = %d\n"
8:		
0:	.comm	X, 4 ; 4 bytes in BSS is required ; for global X
4:		
0:	.text	; offset restarts; what follows is ; text (read-only code)
0:	.globl	subr
0:	subr:	
0:	pushl	%ebp ; save the frame pointer
1:	movl	%esp, %ebp ; point to new frame
3:	pushl	8(%ebp) ; push i onto stack
6:	pushl	\$printfarg ; push address of printf arg ; onto stack
11:	call	printf
16:	addl	\$8, %esp ; pop arguments
19:	movl	%ebp, %esp ; restore stack pointer
21:	popl	%ebp ; pop frame pointer
23:	ret	

```
#include <stdio.h>
int X;

void subr(int i) {
    printf("i = %d\n", i);
}
```

subr.s

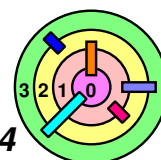
Offset	Op	Arg
0:	.data	; what follows is initialized data
→ 0:	printfarg:	
→ 0:	.string	"i = %d\n"
8:		
0:	.comm	X, 4 ; 4 bytes in BSS ; for global X
4:		
0:	.text	; offset restarts; w ; text (read-only co
0:	.globl	subr
0:	subr:	
0:	pushl	%ebp ; save the fra
1:	movl	%esp, %ebp ; point t
3:	pushl	8(%ebp) ; push i on
6:	pushl	\$printfarg ; push ad ; onto st
11:	call	printf
16:	addl	\$8, %esp ; pop argum
19:	movl	%ebp, %esp ; restore stack pointer
21:	popl	%ebp ; pop frame pointer
23:	ret	

this is how you create a string constant

- ▢ this one is 8 bytes long
- ▢ and local to this module (since it's not global)
- ▢ can "live" somewhere else

```
#include <stdio.h>
int X;
```

```
void subr(int i) {
    printf("i = %d\n", i);
}
```



subr.s

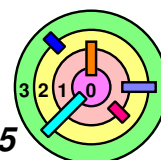
Offset	Op	Arg
0:	.data	; what follows is initialized data
0:	printfarg:	
0:	.string	"i = %d\n"
8:		
0:	.comm	X, 4 ; 4 bytes in BSS ; for global X
4:		
0:	.text	; offset restarts; w ; text (read-only co
0:	.globl	subr
0:	subr:	
0:	pushl	%ebp ; save the fra
1:	movl	%esp, %ebp ; point t
3:	pushl	8(%ebp) ; push i on
6:	pushl	\$printfarg ; push ad ; onto st
11:	call	printf
16:	addl	\$8, %esp ; pop argum
19:	movl	%ebp, %esp ; restore stack pointer
21:	popl	%ebp ; pop frame pointer
23:	ret	

this is how you create a string constant

- ▢ this one is 8 bytes long
- ▢ and local to this module (since it's not global)
- ▢ can "live" somewhere else
- ▢ it is used here

```
#include <stdio.h>
int X;
```

```
void subr(int i) {
    printf("i = %d\n", i);
}
```



subr.s

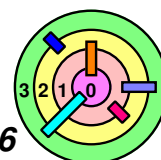
Offset	Op	Arg
0:	.data	; what follows is initialized data
0:	printfarg:	
0:	.string	"i = %d\n"
8:		
0:	.comm	X, 4 ; 4 bytes in BSS ; for global X
4:		
0:	.text	; offset restarts; w ; text (read-only co
0:	.globl	subr
0:	subr:	
0:	pushl	%ebp ; save the fra
1:	movl	%esp, %ebp ; point t
3:	pushl	8(%ebp) ; push i on
6:	pushl	\$printfarg ; push ad ; onto st
11:	call	printf
16:	addl	\$8, %esp ; pop argum
19:	movl	%ebp, %esp ; restore stack pointer
21:	popl	%ebp ; pop frame pointer
23:	ret	

4 bytes is required in the **bss** segment for this global variable

☞ .comm directive also means that the symbol mentioned is defined **here** and is **exported**

```
#include <stdio.h>
int X;

void subr(int i) {
    printf("i = %d\n", i);
}
```



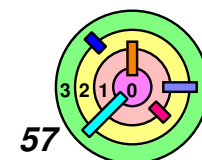
subr.s

Offset	Op	Arg
0:	.data	; what follows is initialized data
0:	printfarg:	
0:	.string	"i = %d\n"
8:		
0:	.comm	X, 4 ; 4 bytes in BSS ; for global X
4:		
0:	.text	; offset restarts; w ; text (read-only co
→ 0:	.globl	subr
0:	subr:	
0:	pushl	%ebp ; save the fra
1:	movl	%esp, %ebp ; point t
3:	pushl	8(%ebp) ; push i on
6:	pushl	\$printfarg ; push ad ; onto st
11:	call	printf
16:	addl	\$8, %esp ; pop argum
19:	movl	%ebp, %esp ; restore stack pointer
21:	popl	%ebp ; pop frame pointer
23:	ret	

subr is a global symbol
exported from here

```
#include <stdio.h>
int X;

void subr(int i) {
    printf("i = %d\n", i);
}
```



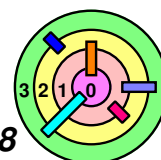
subr.s

Offset	Op	Arg
0:	.data	; what follows is initialized data
0:	printfarg:	
0:	.string	"i = %d\n"
8:		
0:	.comm	X, 4 ; 4 bytes in BSS ; for global X
4:		
0:	.text	; offset restarts; w ; text (read-only co
0:	.globl	subr
0:	subr:	
0:	pushl	%ebp ; save the fra
1:	movl	%esp, %ebp ; point t
3:	pushl	8(%ebp) ; push i on
→ 6:	pushl	\$printfarg ; push ad ; onto st
→ 11:	call	printf
16:	addl	\$8, %esp ; pop argum
19:	movl	%ebp, %esp ; restore stack pointer
21:	popl	%ebp ; pop frame pointer
23:	ret	

relocation is required for
printf and printfarg

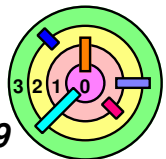
```
#include <stdio.h>
int X;

void subr(int i) {
    printf("i = %d\n", i);
}
```



Object Files

- ➡ An object file describes what's in the data, bss, and text segments in separate sections
- ➡ Along with each section is a list of:
 - ▬ global symbols
 - ▬ undefined symbols
 - ▬ instructions for relocation
 - these instructions indicate
 - ◆ which locations within the section must be modified
 - ◆ which symbol's value is used to modify the location
 - a symbol's value is the address that is ultimately determined for it
 - typically, this address is added to the location being modified
- ➡ To inspect an object file on Unix
 - ▬ `nm` - list symbols from object files
 - ▬ `objdump` - display information from object files



subr.o

Data:

Size: 8

Contents: "i = %d\n"

bss:

Size: 4

Global: X, offset 0

Text:

Size: 24

Global: subr, offset 0

Undefined: printf

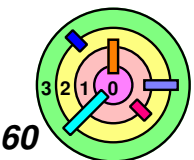
Relocation:

offset 7, size 4, value: addr of printfarg

offset 12, size 4, value: PC-relative addr of
printf

Contents: [machine instructions]

➡ No tool can generate exactly the above printout



subr.s

Offset	Op	Arg
0:	.data	; what follows is initialized data
0:	printfarg:	
0:	.string	"i = %d\n"
8:		
0:	.comm	X, 4 ; 4 bytes in BSS ; for global X
4:		
0:	.text	; offset restarts; w ; text (read-only co
0:	.globl	subr
0:	subr:	
0:	pushl	%ebp ; save the fra
1:	movl	%esp, %ebp ; point t
3:	pushl	8(%ebp) ; push i onto stack
→ 6:	pushl	\$printfarg ; push address of string ; onto stack
→ 11:	call	printf
16:	addl	\$8, %esp ; pop arguments from stack
19:	movl	%ebp, %esp ; restore stack pointer
21:	popl	%ebp ; pop frame pointer
23:	ret	

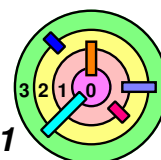
relocation is required for:

▢ printfarg at offset 7

▢ printf at offset 12

```
#include <stdio.h>
int X;
```

```
void subr(int i) {
    printf("i = %d\n", i);
}
```



subr.o

Data:

Size: 8

Contents: "i = %d\n"

bss:

Size: 4

Global: X, offset 0

Text:

Size: 24

Global: subr, offset 0

Undefined: printf

Relocation:

offset 7, size 4, value: addr of printfarg

offset 12, size 4, value: PC-relative addr of
printf

Contents: [machine instructions]

relocation is required for:

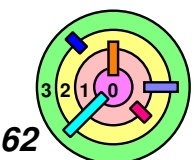
▢ printfarg at offset 7

▢ printf at offset 12

```
#include <stdio.h>
```

```
int X;
```

```
void subr(int i) {  
    printf("i = %d\n", i);  
}
```



subr.o

Data:

Size: 8

Contents: "i = %d\n"

bss:

Size: 4

→ Global: X, offset 0

Text:

Size: 24

→ Global: subr, offset 0

Undefined: printf

Relocation:

offset 7, size 4, value: addr of printfarg

offset 12, size 4, value: PC-relative addr of
printf

Contents: [machine instructions]

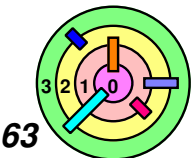
x and subr are *exported*

→ needed in main.o

```
#include <stdio.h>
```

```
int X;
```

```
void subr(int i) {  
    printf("i = %d\n", i);  
}
```



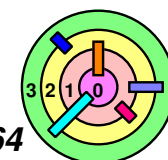
main.s

Offset	Op	Arg
0:	.data	; what follows is initialized data
0:	.globl	aX ; aX is global: ; by others
0:	aX:	
0:	.long	X
4:		
0:	.text	; offset restarts; w ; text (read-only co
0:	.globl	main
0:	main:	
0:	pushl	%ebp ; save the fram
1:	movl	%esp,%ebp ; point to
3:	subl	\$4,%esp ; make space
6:	movl	aX,%eax ; put conten
11:	movl	(%eax),%eax ; put *X
13:	movl	%eax,-4(%ebp) ; stor
16:	pushl	-4(%ebp) ; push y onto stack
19:	call	subr
24:	addl	\$4,%esp ; remove y from stack
27:	movl	\$0,%eax ; set return value to 0
31:	movl	%ebp, %esp ; restore stack pointer
33:	popl	%ebp ; pop frame pointer
35:	ret	

these 2 places remained
unresolved

```
extern int X;
int *aX = &X;

int main( ) {
    void subr(int);
    int y = *aX;
    subr(y);
    return(0);
}
```



main.o

Data:

Size: 4

Global: aX, offset 0

Undefined: X

Relocation: offset 0, size 4, value

Contents: 0x00000000

these 2 places remained

unresolved

→ they are noted in main.o

bss:

Size: 0

Text:

Size: 36

Global: main, offset 0

Undefined: subr

Relocation:

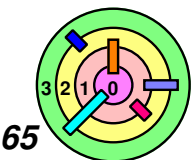
offset 7, size 4, value: addr of aX

offset 20, size 4, value: PC-relative
addr of subr

Contents: [machine instructions]

```
extern int X;
int *aX = &X;

int main( ) {
    void subr(int);
    int y = *aX;
    subr(y);
    return(0);
}
```



subr.o

Data:

Size: 8

Contents: "i = %d\n"

bss:

Size: 4

Global: X, offset 0

Text:

Size: 24

Global: subr, offset 0

Undefined: printf

Relocation:

offset 7, size 4, value: addr of printfarg

offset 12, size 4, value: PC-relative addr of
printf

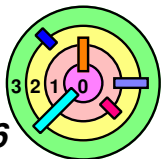
Contents: [machine instructions]

printf remained *unresolved*

```
#include <stdio.h>
```

```
int X;
```

```
void subr(int i) {  
    printf("i = %d\n", i);  
}
```



printf.o

Data:

Size: 1024

Global: StandardFiles

Contents: ...

assume that printf.o looks like this

→ write is *unresolved*

bss:

Size: 256

Text:

Size: 12000

Global: printf, offset 100

...



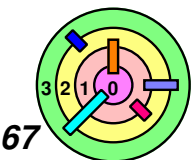
Undefined: write

Relocation:

offset 211, value: addr of StandardFiles

offset 723, value: PC-relative addr of write

Contents: [machine instructions]



write.o

Data:

Size: 0

bss:

Size: 4

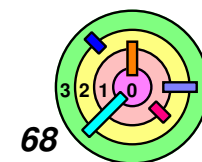
Global: errno, offset 0

Text:

Size: 16

Contents: [machine
instructions]

and write.o looks like this



startup function

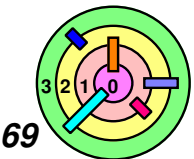
```
Data:
    Size:  0

bss:
    Size:  0

Text:
    Size:  36
    Undefined: main
    Relocation:
        offset 21, value: main
    Contents: [machine
               instructions]
```

every C program contains a *startup* routine that is called first

- ▢ it calls `main()`
- ▢ if `main()` returns, it calls `exit()`
- ▢ our example is incomplete



prog

Text

main	4096
subr	4132
printf	4156
write	16156
startup	16172

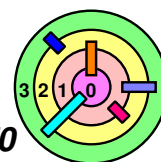
Data

aX	16384
printfargs	16388
StandardFiles	16396

BSS

X	17420
errno	17424

- this is how `ld` might set things up
 - `ld` *lays out* the *address space*
 - `ld` allocates memory in *pages* (typically 4KB each)
- `main` does not start at location 0
 - first "page" is typically made inaccessible so that references to null pointers will fail (get SIGSEG)



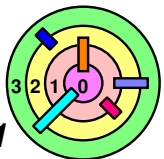
prog

Text			}	R/O
main	4096			
subr	4132			
printf	4156			
write	16156			
startup	16172			
Data			}	R/W
aX	16384			
printfargs	16388			
StandardFiles	16396			
BSS			}	R/W
X	17420			
errno	17424			

printfargs can be modified
by the application

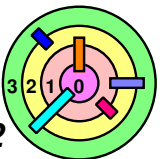
- ▢ is that okay?
- ▢ where else would you put it?
- ▢ who decides?

- ▢ due to the use of "*pages*", the data segment needs to start at a page boundary (i.e., multiple of page size)
 - this way, the *text segment* can be made *read-only* while the *data* and *bss segments* made *read-write*
 - here we assume 4KB pages (therefore, pages start at 4096, 8192, 12288, 16384, etc.)



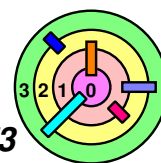
Page Protection & Virtual Memory Basics

- ➡ A process has, say, a 32-bit address space
- ➡ that's 4GB of memory
 - ➡ our prog process, when it starts, only needs about 16KB for text+data+bss (plus more for stack)
 - ➡ allocating 4GB of physical memory will be a *huge waste*



Page Protection & Virtual Memory Basics

- ➡ A process has, say, a 32-bit address space
 - ➡ that's 4GB of memory
 - ➡ our prog process, when it starts, only needs about 16KB for text+data+bss (plus more for stack)
 - ➡ allocating 4GB of physical memory will be a *huge waste*
- ➡ Solution: *indirection*
 - ➡ use *page table* to *map virtual to physical addresses*
 - a page table is a *kernel data structure*, used by the *CPU*
 - ➡ OS allocates *pages* of *physical memory* using the *Buddy System*
 - a page is 4KB in many systems
 - ➡ a page corresponds to physical memory that can be located (or "*mapped*") anywhere in virtual memory
 - "*address translation*" done in hardware
 - ➡ some advantages:
 - page protection
 - only allocate needed physical memory pages
 - ➡ a lot more to come in Ch 7



Page Protection & Virtual Memory Basics

Text

main 4096
 subr 4132
 printf 4156
 write 16156
 startup 16172

Data

aX 16384
 printfargs 16388
 StandardFiles 16396

BSS

X 17420
 errno 17424

Page Table

#	Start	Access	Physical Addr
0	0	-	-
1	4096	R	•
2	8192	R	•
3	12288	R	•
4	16384	R/W	•

ask buddy system to
allocate these pages

