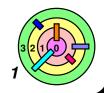
CS 350 PA3: Mutex & Condition Variable

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Based on slides created by Kivilcim Cumbul



PA3



Start with your PA2 code (no separate starter code will be provided)

```
cd
mkdir cs350/pa3
mkdir cs350/pa3/xv6-pa3-src
cd cs350/pa3/xv6-pa3-src
cp ../../pa2/xv6-pa2-src/* .
rm -f pa2-submit.tar.gz
make clean
```

- PA3 only makes sense if you have multithreading
 - if your PA2 code is not working, you will not be able to pass any of the PA3 tests and you will end up with a very low score
- Part 1 preparation
- reading code and documentation
- Part 2 implement mutex functions
 - kthread_mutex_alloc(), kthread_mutex_dealloc(),
 kthread_mutex_lock(), kthread_mutex_unlock()
 - no condition variables!



Submission



Which files do you need to modify?

open a terminal and type the following:

```
pwd
cd ~/cs350/pa3/xv6-pa3-src
make -n pa3-submit
```

you should see:

```
tar cvzf pa3-submit.tar.gz \
    Makefile \
    pa3-README.txt \
    proc.c \
    proc.h \
    syscall.c \
    sysproc.c \
    kthread.h \
    exec.c
```

- some files may be the same as in your PA2 submission
- these are the only files are are supposed to submit
 - if you submit additional files, the grader will have to delete them before grading
 - if you submit binary files, points will be deducted

Part 1: Preparation

Read POSIX Threads Programming tutorial by Blaise Barney from Lawrence Livermore National Laboratory

pthread_mutex_lock(), pthread_mutex_unlock()



Read the spinlock code in XV6



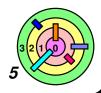
Part 2: Implement Mutex Functions



Implement mutex API for kernel

- wthread_mutex_alloc(), kthread_mutex_dealloc(),
 kthread_mutex_lock(), kthread_mutex_unlock()
- unless otherwise specified, we use the term *lock* and *mutex* interchangeably (although in general, a lock may allow multiple threads to have concurrent access to a resource)
- you will implement these functions in "proc.c" and add the following to "kthread.h"

```
#define MAX_MUTEXES 64
int kthread_mutex_alloc();
int kthread_mutex_dealloc(int mutex_id);
int kthread_mutex_lock(int mutex_id);
int kthread_mutex_unlock(int mutex_id);
```



Changes In "proc.h"



Mutex data structures

what are the possible mutex states?

```
enum mutexstate { MUNUSED, MLOCKED, MUNLOCKED };
```

- what should go into a mutex struct?
 - at a minimum:

```
int mid; // unique mutex ID ≥ 1
enum mutexstate;
```

Changes In "proc.c"



Similar to ptable, we need a mtable



Need a global variable to know what mutex ID to return next

```
int nextmid = 1;
```

must never reuse a mutex ID



kthread_mutex_alloc()



Allocates a mutex object and initializes it; the initial state should be unlocked

should return the ID of the initialized mutex, or -1 upon failure



Looping Through Mutex Table



Something like the following:

```
for (m = mtable.mutex; && m < &mtable.mutex[MAX_MUTEXES]; Mp++) {
    ...
}

Or:

for(i = 0; i < MAX_MUTEXES; i++) {
    m = &mtable.mutex[i];
    ...
}</pre>
```



kthread_mutex_alloc()

```
kthread_mutex_alloc():

    create a mutex pointer m (using struct in "proc.h")
Loop through mutex table
    if m is unused
        m->mutex_id = nextmid++;
        m->state = MUNLOCKED;
        initialize all other values if needed

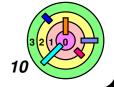
if m == &mtable.mutex[MAX_MUTEXES]
    return -1
else
    return m->mutex_id
```

- Note: the above is not the only way to implement mutex allocation
- also, this is not a complete pseudocode
 - you have to add locks if necessary



kthread_mutex_dealloc()

- De-allocates a mutex object which is no longer needed
- the function should return 0 upon success and -1 upon failure
 - if the given mutex is currently locked, this function should return -1



kthread_mutex_dealloc()

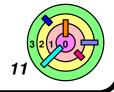
```
kthread_mutex_dealloc(int mutex_id):

create a mutex pointer m (using struct in "proc.h")
loop through mutex table to find given mutex_id
   if m is locked
      return -1

if not found
   return -1

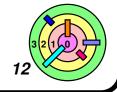
else
   m->mid = 0
   m->state = MUNUSED
   zero out all the other values if needed
   return 0
```

- Note: the above is not the only way to implement mutex deallocation
- also, this is not a complete pseudocode
 - you have to add locks if necessary



kthread_mutex_lock()

- This function is used by a thread to lock the mutex specified by the argument mutex_id
- if the mutex is already locked by another thread, this call will block the calling thread (change the thread state to TBLOCKED) until the mutex is unlocked
 - you may add a TBLOCKED state in "proc.h" if you'd like



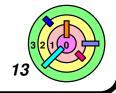
kthread_mutex_lock()

```
kthread_mutex_lock(int mutex_id):

create a mutex pointer m (using struct in "proc.h")
loop through mutex table to find the target mutex id (parameter)
   if m->mid == mutex_id
      break;
if not found
   return -1
while (m->state == MLOCKED)
   sleep // on m
if (m->state != MUNLOCKED)
   return -1

m->state = MLOCKED
return 0
```

- Note: the above is not the only way to implement mutex lock
- also, this is not a complete pseudocode
 - you have to add locks if necessary



kthread_mutex_unlock()



This function unlocks the mutex specified by the argument mutex_id if called by the owning thread, and if there are any blocked threads, one of the threads will acquire the mutex

- the mutex may be owned by one thread and unlocked by another
- an error will be returned if the mutex was already unlocked



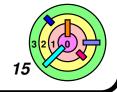
kthread_mutex_unlock()

```
kthread_mutex_unlock(int mutex_id):

create a mutex pointer m (using struct in "proc.h")
loop through mutex table to find the target mutex id (parameter)
   if m->mid == mutex_id
      break;
if not found
   return -1
while (m->state == MUNLOCKED)
   return -1

m->state = MUNLOCKED
call wakeup on m to wake up all threads waiting for this mutex
return 0
```

- Note: the above is not the only way to implement mutex lock
- also, this is not a complete pseudocode
 - you have to add locks if necessary



Common Errors



You are supposed to read the code of the test programs

- mutextest1.c
- mutextest2.c



You are supposed to be reading the XV6 book: xv6-rev11.pdf and the XV6 source code to understand how the spinlock (and locking in general) works in XV6



Common Errors

panic: sched locks



this means process holding multiple locks

before calling sched() make sure to release all other locks

```
ac (ptable)
ac (mtable)
rel (mtable)
rel (ptable)
```

